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OBJECT-ORIENTED DESIGN
OF
TOKEN RING NETWORK

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BY

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CERTIFICATE

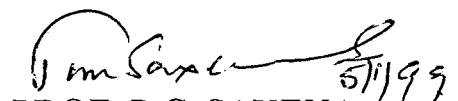
This is to certify that the dissertation entitled “ **OBJECT-ORIENTED DESIGN OF TOKEN RING NETWORK** “ which is being submitted by Mr. RAM BHAGAT to the School of Computer & Systems Sciences, JAWAHARLAL NEHRU UNIVERSITY for the award of Master of Technology in computer science is a bonafide work carried out by him under my supervision .

This work is original and has not been submitted in part or full to any university or institution for the award of any degree .


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RAM BHAGAT

ABSTRACT

The primary objective of this work is to design the Token Ring Network using object-oriented methodology based on specialization theory. In this work we first described object-oriented concepts which will be applied while designing the Token Ring Network. We further described Token Ring specification and commenced the designing process of token ring .In the next part we have given the full code of token ring in Abstract Syntax Notation.

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CHAPTER 1

INTRODUCTION

Networks are so complex that it is imperative to describe the network in terms of a model . The process of modelling takes the problem into the realm of the abstract where it is easier to find the solution. The solution of abstract problem then can be converted into concrete solutions. A good modelling technique provides a mechanism by which a complex problem can be decomposed into parts and those parts into further subparts. The overall model of the complex system describes not only components which results from such a decomposition but also the relationships among these components. The modelling technique must not only provide structure but also suggest function. A good modelling technique has to make balance between complexity and simplicity depending on the requirements of the application. Object-oriented modelling is a technique that allows classification based not only on structure but also on behavior. One of the reasons for the popularity of the object-oriented techniques is the intellectual appeal of their ability to capture both structure and function in an object. Some modelling techniques like entity-attribute modelling describe the problem domain only in structural terms whereas others like state transition analysis do so only in behavioral terms. The analysis of communication networks needs both aspects . A pure structural model of a communication network describes the nature of the elements that comprise it and their relationships with each other but would not capture behavioral aspects in terms of protocol message and responses. A pure behavioral model would describe their function but would not capture their structural information.

CHAPTER 2

OBJECT-ORIENTED METHODOLOGY

2.1 Object-Oriented themes

In the object-oriented model the themes of object-oriented analysis abstraction, encapsulation, inheritance all find expression. Abstraction is a mechanism for coping with complexity. We can concentrate only on essential details for solving a specific problem. Abstraction draws a boundary around an object inside which are the essential characteristics of the object from the point of view of the application domain. Essential characteristics isolated inside the abstraction barrier must capture the notion of both structure and function based distinctions. A correctly defined abstraction allows the same object model to be reused in various ways.

Encapsulation consists of identifying the internal implementation of the network element and separating those from its externally visible behaviour. Encapsulation preserves the integrity of the object - the underlying implementation may be changed as long as the interface visible to other objects remain consistent. Encapsulation and abstraction are complementary to each other. Just as abstraction separates the essential characteristics of an object from the non-essential ones, the encapsulation separates the externally visible characteristics from the hidden ones. The process of abstraction and encapsulation enables us to talk about network operations completely in the abstract.

An object class is a set of objects having common structural and behavioural properties. All objects in a class have the same purpose. The process of classification in combination with abstraction allows us to categorise the different elements that comprise a communication network as lines, circuits, LAN bridges, softwares, services switching equipments etc. In addition to this classes also serve as templates with which to create new objects. The attributes of an object describe those data values that the object possesses which could conceivably be different from the data

values possessed by other objects of the same class. When we assign an attribute to a class we say that the attribute is-a-property of that class . Conversely the class has-as-a-property each attribute assigned to it. The terms is-a-property-of and has-as-a-property are called property assignment associations. Each attribute has a data type which defines the nature and range of values the attribute can possess. The set of all values which an attribute may possibly possess is known the domain of the attribute. An attribute can be specified unique i.e. each object instance will have different values for that attribute.

The function describes the behavioral properties of an object. Functions are of two types procedural function and stream function. Procedural functions accept arguments and produce results which are well defined data types. Stream functions have at least one argument or one result that is not a well defined data type but is an unstructured continuous analog or digital stream.

Inheritance provides a way of classifying classes. If we find two similar object classes sharing a subset of their properties , we abstract their common properties into a superclass. We thus have a categorization of classes . The inheritance hierarchy provides a second order abstraction in which common properties and behavior can be isolated in yet other classes . This second order abstraction is the defining characteristic of object-oriented modelling paradigm. Inheritance is also a mechanism for implicit property assignment i.e. by assigning a property to an ancestral class , it becomes available in all its descendants. The properties available in an object class through inheritance from ancestral class are called inherited properties of that class. In addition to inheriting characteristics from its superclass, the subclass may choose to define additional characteristics of its own . This is called as extension . The properties defined by subclass on its own are called original properties. The class where a property is first defined is called originating class. Inheritance and extension are mechanisms to express the semantics of classification and evolution . They are a means of incorporating a formalized taxonomy within the object model itself

A paradigm that supports abstraction and encapsulation is said to be object based ,one that ,in addition , supports classification is said to be class based and further if it also

supports inheritance it is said to be object-oriented. An object instance is a concrete object which is an example of an object class. It has an identity which is different from other instances of the same class. The process of creating object of a class is called instantiation. A class which has direct instances is called a concrete class. Those classes which are not instantiated are called abstract classes. Classes, which not only themselves but whose descendants, are also abstract are called fully abstract classes.

Aggregation refers to the ability to construct complex objects by assembling simpler objects together in a meaningful configuration. Decomposition is a process of breaking down the complex system in a set of subproblems each of which we know how to solve independently. While decomposing we must take care that each part should preserve encapsulation and present a well defined set of interfaces at its boundary. An aggregation hierarchy identifies the part/whole relationships between objects.

The modelling methodology applied here is based on specialisation theory. It is used to organize unstructured knowledge to meaningfully represent a system. Here emphasis lies in understanding the networks in terms of its component objects, their attributes and the relationship they hold with each other. Object-oriented model serves as a basis for specification, design and documentation.

2.2 Object-oriented modelling v/s object-oriented programming

Object-oriented modelling is different from object-oriented programming in various aspects e.g. in object-oriented modelling inheritance does not necessarily mean reusability of implementations, it simply means reusability of specifications. Here inheritance is a mechanism of incremental improvement of the classes already defined. In specialisation theory which is stronger form of object-oriented modelling monotonic inheritance is applied i.e. already existing features of classes are not allowed to be cancelled. In case of object-oriented programming we can drop the properties(attributes or functions) of the superclass. Further multiple inheritance is

rarely used in object-oriented modelling . In fact it can be done away by using aggregation and hierarchy redesign . While in object-oriented programming it is encountered more frequently. In case of modelling overriding (a process by which subclass can redefine inherited attributes and functions) is not permitted but in object-oriented programming this feature is available and is often implemented by virtual functions. In object-oriented modelling polymorphism does not make much sense as here we are not concerned with implementation variations but in programming it plays an important role. Similarly the concept of dynamic binding is not meaningful in object-oriented modelling of networks. Again we do not use delegation which is used in certain object-based software environments. Delegation as the name suggests is a property by which an object can delegate its responsibility of performing certain operations to some other object which in turn can delegate this responsibility to some object.

Object-oriented programming provides a means of visibility control by the concept of class and superclass but there is no such need of visibility control in object-oriented modelling. In object-oriented programming declarative instantiation i.e. having defined a class in a formal syntax one can then instantiate variables of the class through declarations in the same syntax. However declarative instantiation is not meaningful in object-oriented modelling of communication networks . The issue of inter-object communication is very straightforward in object-oriented programming while it is much complicated in object-oriented modelling of communication networks. Referent and non-referent both types of classes are referred in object-oriented programming but not in object-oriented network model. A subclass can be arbitrarily written while overriding the original behavior of the class but in a subtype its original has to be maintained. It can be modified only in permissible ways. In object-oriented network modelling as we are applying specialisation theory the difference between subclasses and subtypes is irrelevant. This difference is significant in many object-oriented languages. Anyhow object-oriented modelling and object-oriented programming are complementary to each other. They both are of benefit in different phases of system development . The programmer need no longer expend

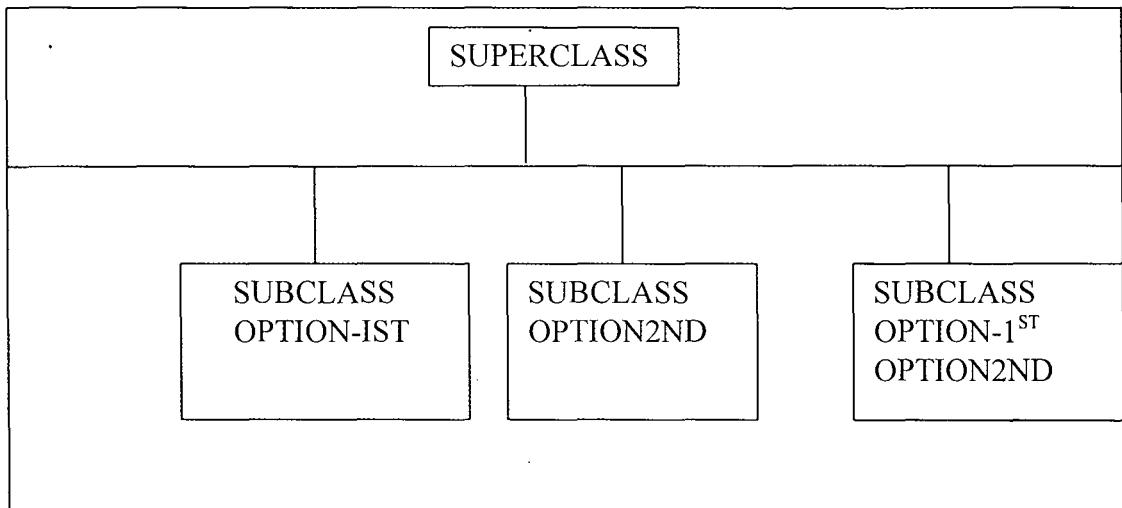
effort for designing software classes for network objects as the hierarchy obtained by object-oriented modelling can be converted into programming language.

2.3 Object Class

A meaningful object class in a network model must have a referent in the actual system i.e. it models some entity in the actual system which is relevant to our purposes. Object classes may be physical entities or logical entities e.g. protocol entities, network services and sites . Every object class is modelled by identifying attributes and functions of interest. The class is described as a modelling construct and its attributes and functions are described as independent modelling constructs which are associated with the class through property assignment.

There are two types of properties those which are mandatory and those which are optional. Mandatory properties are captured by using core specification while optional properties are captured using variant specification. Let us first describe the core properties .Here we identify the properties which must be present in every instance of the class i.e. if an object lacks any of the core properties then it is not an instance of that class but while identifying the attributes and functions we must be careful that they are orthogonal i.e. they are not functionally dependent on each other in any way. Thus we should not be in a position to compute any attribute or function in terms of other attributes e.g. suppose we are trying to find out three parts of a number so that their product is maximum. In this case if we know first two parts of the number then the third part can be computed just by subtracting the sum of the two parts from the number. Thus if we represent this third part as also an attribute then these attributes will not be orthogonal as a relationship exists between them. This may cause sometimes integrity violation . An object not having any knowledge of the dependency between the attributes may try to set them to values which are inconsistent with each other. Such an attribute can be specified by formally modelling their functional dependency on other attributes as an algorithm. While specifying functions care must be taken to specify function arguments such a way that the number of functions required can be minimized e.g. if there are two functions

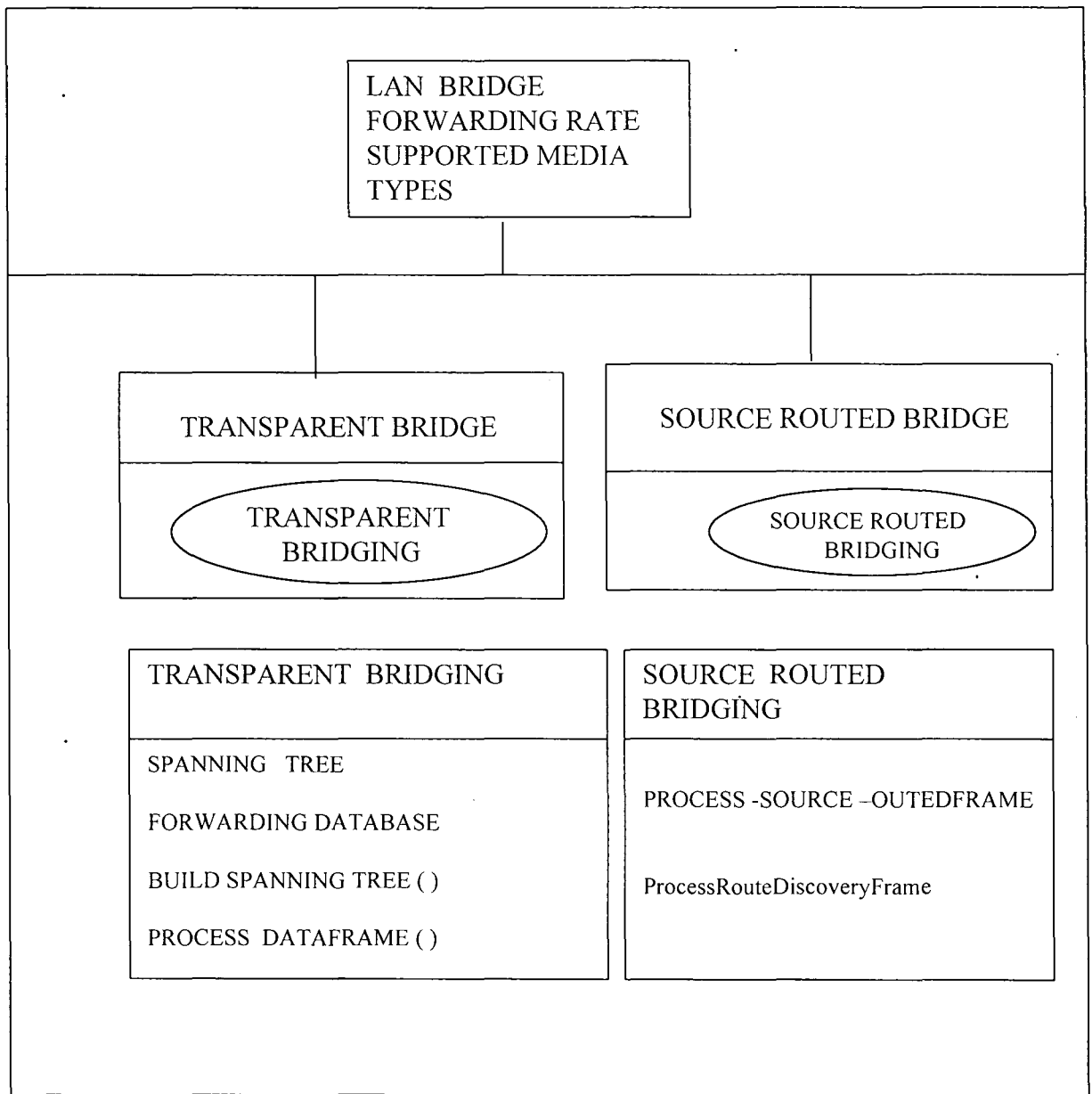
five minutes service and ten minutes service then they should not be modelled as two separate functions instead they should be modelled as a single function and an argument duration should be provided to that function. But this can be done only if the two functions have the same purpose. If their purpose is different then they must be modelled as separate functions. In communication networks there are variety of devices, protocols and services. There are different types of versions and options .. If the core property in a superclass has been defined then it can be extended into different subclasses using inheritance e.g. if there are two options option first and option second then the subclasses can be obtained as shown in the figure .



2.4 Capsules

In this way if we have n options then we will get $2^n - 1$ subclasses which gets out of hand for values like 20 or more. Therefore it complicates the whole design of the class hierarchy. Thus it is better to provide variant property specification. There certain properties which may be present or absent together. Such a collection of properties is modelled by the concept of capsule which may be named. A capsule is never instantiated independently . It is a mechanism to add flexible subsets to object classes . A capsule can be reproduced just by a reference to an already existing capsule . The properties of one capsule should not overlap those of others or with the core properties defined for a given class. They also reduce the need of modelling multiple

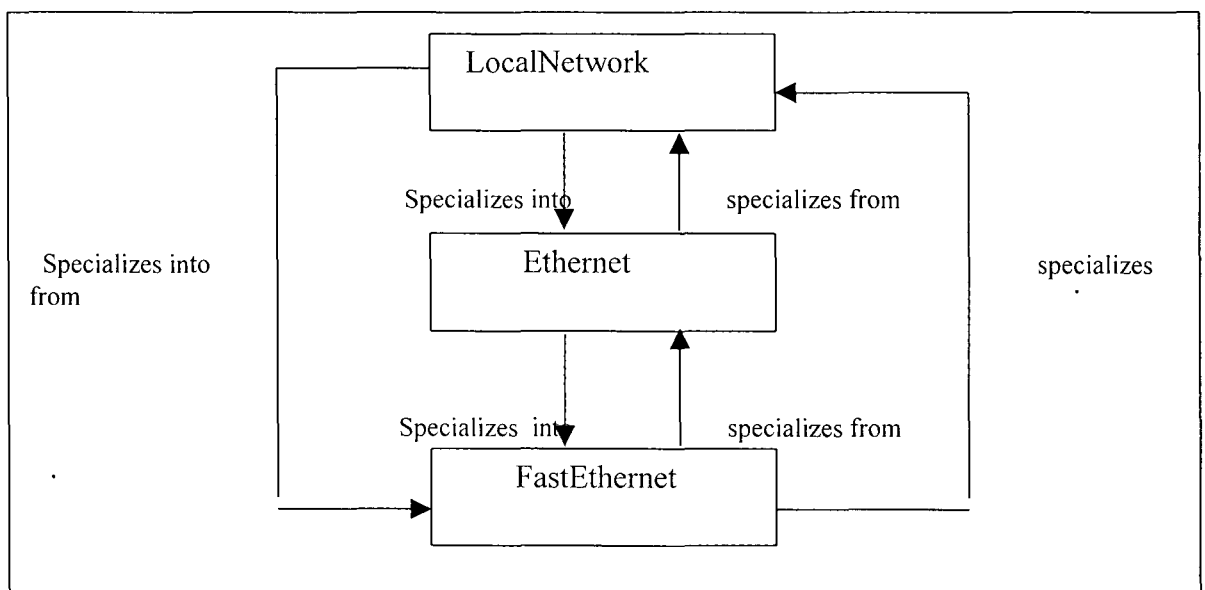
inheritance. A capsule provides a powerful abstraction to describe variant or optional behavior. A capsule may be mandatory or optional. In a capsule if we want to make an attribute to be optional then the entire capsule has to be made optional. It is in consistence with practice as almost all optional characteristics are collections of multiple related attributes and functions which occur as an assemblage. The presence or absence of an optional capsules is also specified using a flag attribute which specifies which optional capsules are present in that instance. In the process of specialisation an optional capsule can be fixed i.e. the capsule becomes mandatory in that class and all the descendant classes. This is termed as capsule fixing . Capsules can be nested with other capsules but not with themselves. Capsules are used as modelling abstractions to specify and import reusable assemblages of properties in network objects as shown in the figure. TransportBridging and sourceRoutedBridging are two capsules. Capsules also help in avoiding multiple inheritance. This accelerates processing in the model information base since multiple inheritance is costly to compute. Further it maintains exactness in the taxonomical structure for assigning instances to the class.



2.5 Generalisation & Specialisation

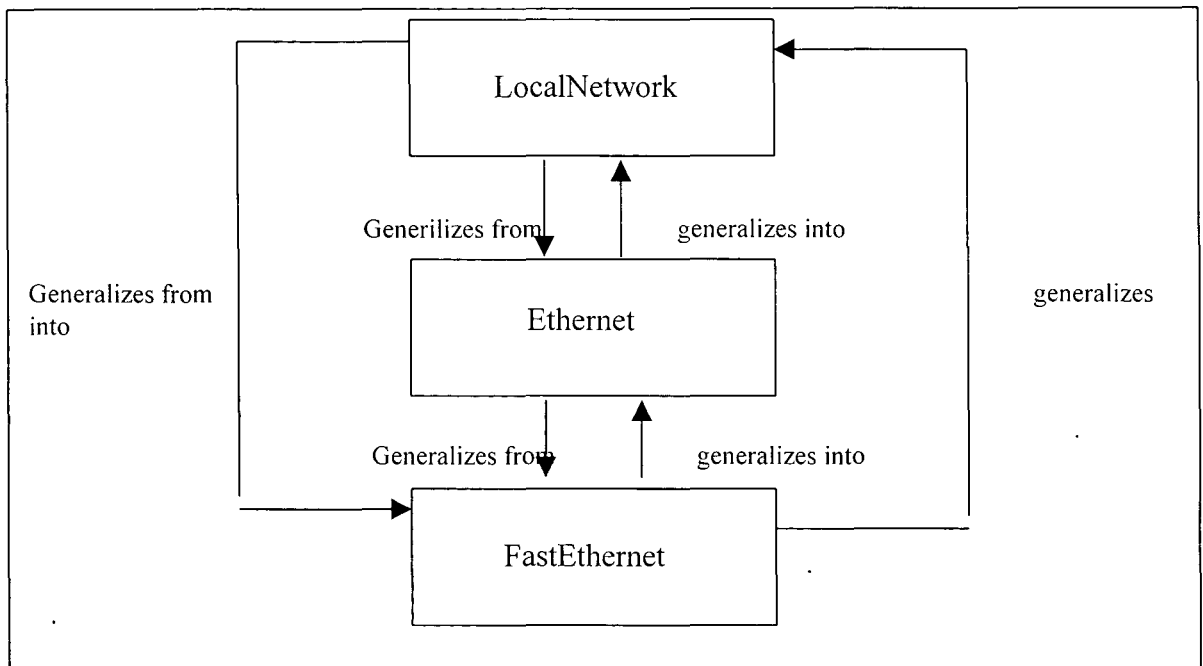
Generalisation and specialisation are the basic themes of inheritance hierarchy . As we go up in the inheritance hierarchy generalisation increases and as we go down the hierarchy specialisation increases . Generalisation is the process constructing an inheritance hierarchy in a bottom-up fashion . It is an exercise of identifying

similarities between object classes, abstracting them and raising them to the level of a superclass . The classes at the top of the hierarchy are simpler and the classes at the bottom of the hierarchy are more complicated in terms of the structure and advanced in terms of the functioning. The construction of the hierarchy in top-down fashion by providing more attributes and functions to the superclass is specialisation. Its purpose is to focus on more specific aspects of the problem domain. While creating the hierarchy we first apply the principle of generalisation and once a hierarchy is constructed specialisation is useful for assigning new object classes a place in the existing hierarchy. As we go up higher in the inheritance tree there comes a point from where no further abstraction is possible . This is known as the root of the inheritance hierarchy . We call it by the special name generic object. It captures the only common properties of every possible object module in the application domain . There is a specialises-into association between a superclass and its subclass. It is transitive from a class to all its descendant classes e.g. if the ethernet class derives from localnetwork class and fastethernet derives from ethernet class i.e. localnetwork class specialises into ethernet class which in turn specialises fastethernet class . Then localnetwork class also specialises into fastethernet class . The complement association is specialises from . That is we can say that ethernet class specialises from localnetwork class . These are shown in the figure as below .



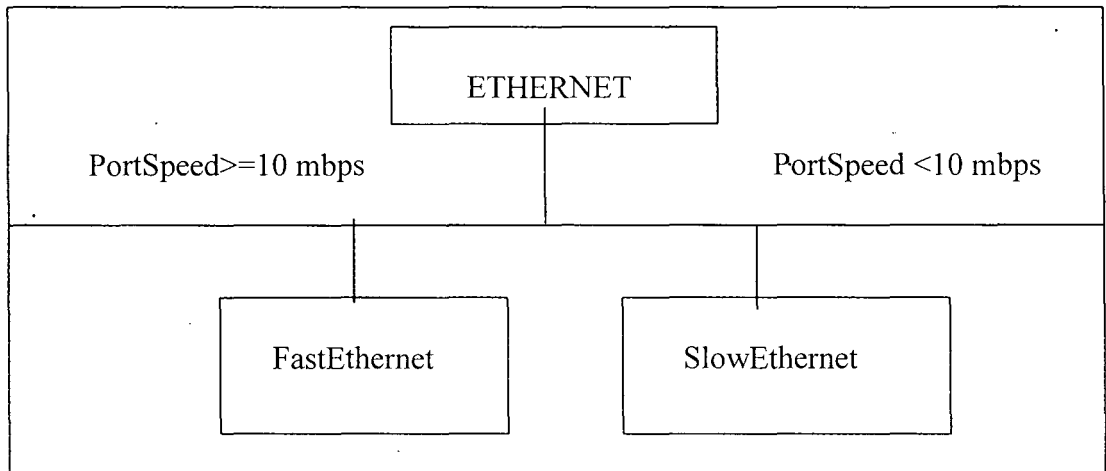
These associations are transitive but not symmetric e.g. ethernet class specialises from localnetwork class but localnetwork class does not specialises from ethernet class. All object classes can be enumerated from the root class by traversing the transitive closure of its specialises into association . It means determining all nodes on the graph which are reachable using the starting node. There are corresponding associations of generalisation as well . There are two associations generalises-from and generalises-into . A superclass generalises from a subclass and a subclass generalises into a

superclass . These associations are also transitive but not symmetric. They can be shown as below.

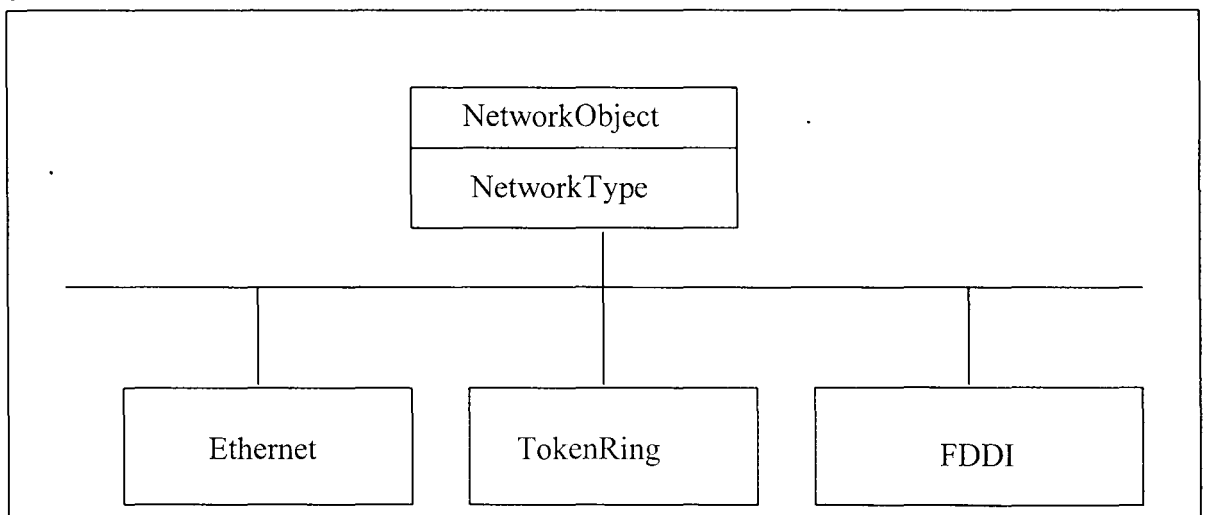


It is a good modelling practice to take care that concrete classes are leaf level classes in the hierarchy as much as possible. It is possible that some non-leaf classes also have instances . It is possible when a new subclass is added making the already existing class

a non-leaf class. In specialisation theory generalisation and specialisation are formally described in terms of a basis property i.e. when we have to specialise a new object class from an existing class we must choose a basis of specialisation. It acts as a distinguishing factor between different subclasses. Generally this basis property is an attribute of the superclass . However it can be some function based property also . As a subclass is more specialised form it has more capabilities than its superclass. Actually it operates in a narrower range of the overall problem space . The basis property which has more domain in the superclass is restricted in domain in the subclass . Then there are two kinds of basis of specialisation Quantitative and Qualitative .If the basis property is a numerical attribute of the superclass specialisation is done by restricting the domain of the values that the attribute is permitted to take in the subclass . This is known as quantitative basis of specialisation . If the basis of property is an enumerable attribute of the superclass, specialisation is done by restricting the set of enumerated values this attribute can have in the subclass. This is called qualitative basis of specialisation

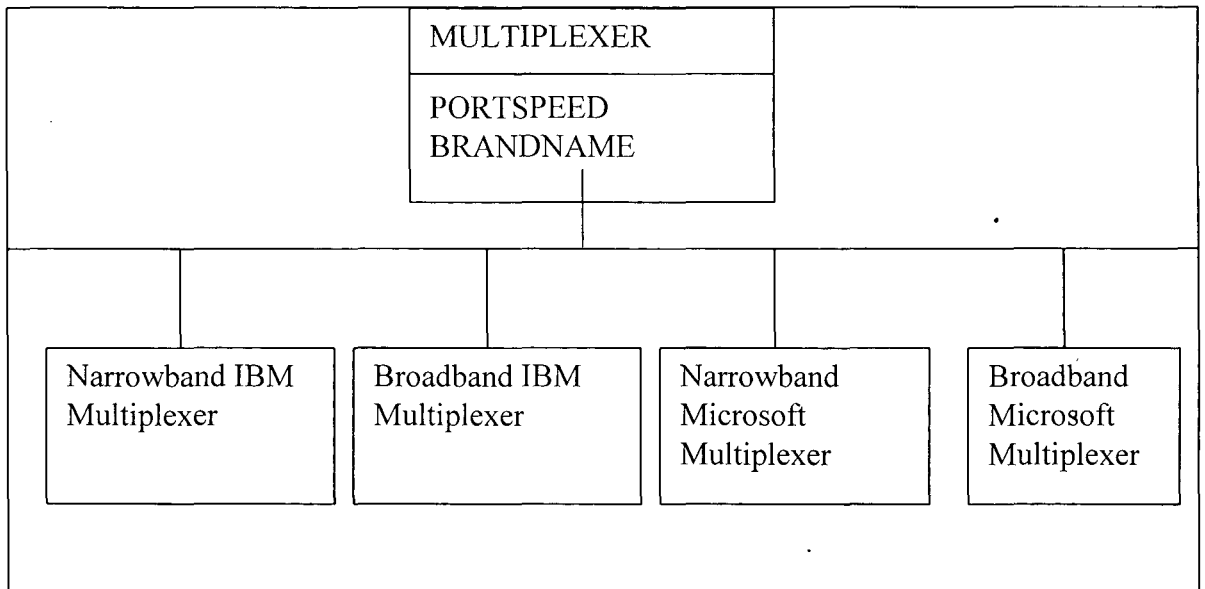


Consider a superclass ethernet which is specialised into two subclasses fastethernet and slowethernet depending on the portspeed. Thus it is an example of quantitative basis of specialisation . Here portspeed attribute has been restricted in domain for subclass.

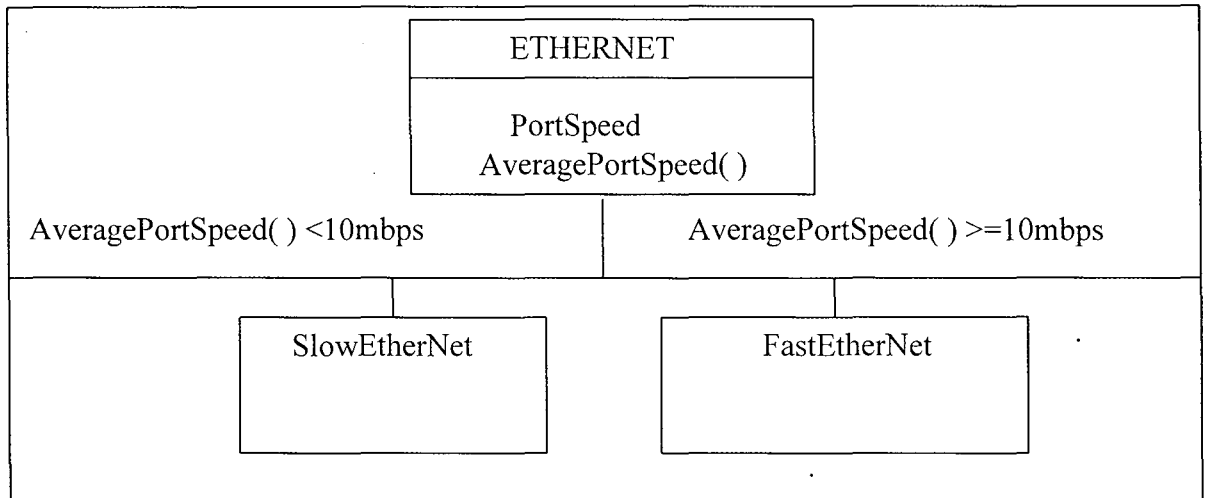


As shown in the figure there are three types of network object ethernet, tokenRing and fddi depending on the value of the attribute network type. This networktype is here basis of specialisation . When specialisation is done by restricting the domain of an attribute of the superclass the mode of specialisation is known as specialisation by Attribute domain restriction There are two notions associated with specialisation disjointness and completeness. A disjoint specialisation means that all the subclasses of a superclass must form non-intersecting sets of objects i.e. an object which is a member of a superclass can be a member a member of only one subclass. An object can not be a member of more than one class i.e. two classes can not have any object in common e.g. consider the example in which we considered ethernet superclass and specialised it on the basis of portspeed . If we take portspeed ≥ 10 mbps for fastethernet and portspeed ≤ 10 mbps for slowethernet then we have overlapping i.e. these two subclasses are not disjoint . For portspeed = 10 mbps objects belong to both

subclasses .A complete specialisation implies that every object of the superclass has to be a member of some subclass . Considering the same example an object of ethernet class has to be a member of fastethernet or a member of slowethernet subclass as by the definition its portspeed has to be either ≥ 10 mbps or < 10 mbps . The notions of disjointness and completeness are also useful when basis of specialisation is qualitative . In network modelling it is better to have subclassing disjoint and complete as far as possible . It improves the richness of the taxonomical structure by removing any ambiguity and confusion with regard to where instances can be assigned in the hierarchy . Furthermore completeness assures us that each instance can be assigned to at least one class. Sometimes we need to choose more than one basis property for the same specialisation . Then it is called compound specialisation . For example if two different basis attributes undergo domain restriction simultaneously as part of the same specialisation as shown in the figure .



Compound specialisation can be replaced by simple specialisation at more than one levels. There is no harm in using compound specialisation if it makes sense for the application domain and number of subclasses at each level are manageable. Another mode of specialisation is known as specialisation by argument contravariance. Instead of restricting the domain here we relax the domain of the arguments of the function. Thus the function possesses now more versatility. They can handle more situations. This is known as contravariance. It is named so because as subclasses become specialised and restricted in their basis properties the arguments that their function can handle actually become weaker and relaxed. As all the subclasses are actually subtypes here therefore contravariance works. We may specialise by tightening function results. In fact finding an attribute's value may be internally implemented as the result returned by a function computation. So as we can restrict domain of an attribute to specialise. Similarly we can do the same thing with result of a function. Consider the figure



Thus if the return value of the function averageportspeed () is less than 10 mbps then it is specialised into slowethernet otherwise fastethernet .Here the basis of specialisation is the result of the function averageportspeed(). Function results can be restricted for specialisation like attribute domain . Thus function results are said to be covariant with specialisation . We can also do specialisation by the capsule fixing i.e. if we create a subclass from an existing object class in which a capsule is optional we make that capsule mandatory . All instances of the subclass and the classes further derived will now possess this capsule . Now we describe certain specialisation principles. A specialisation principle is a rule which states how a subclass can be legally derived from a superclass .

2.6 Specialisation Principles

Specialisation principle of attribute addition:- A descendant class may add an attribute as new property .

Specialisation principle of Function Addition:- A subclass may add a function as a new property .

Specialisation principle of capsule addition :- A subclass may add a capsule as a new property

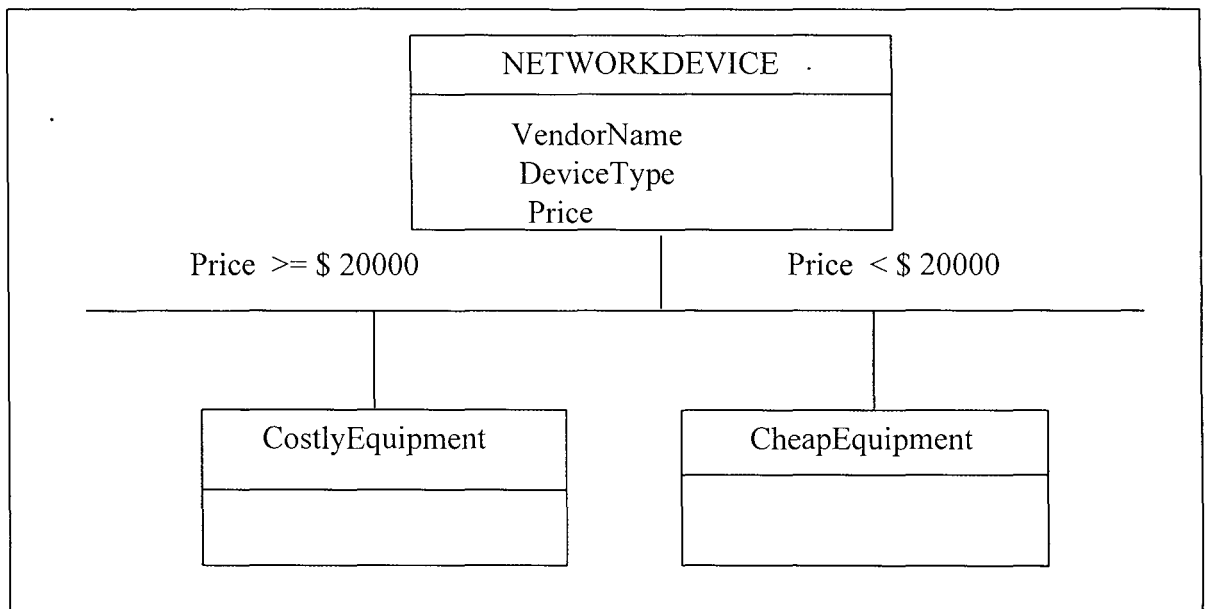
Specialisation principle of attribute domain restriction :- A subclass may restrict the domain of an attribute inherited from an ancestral class.

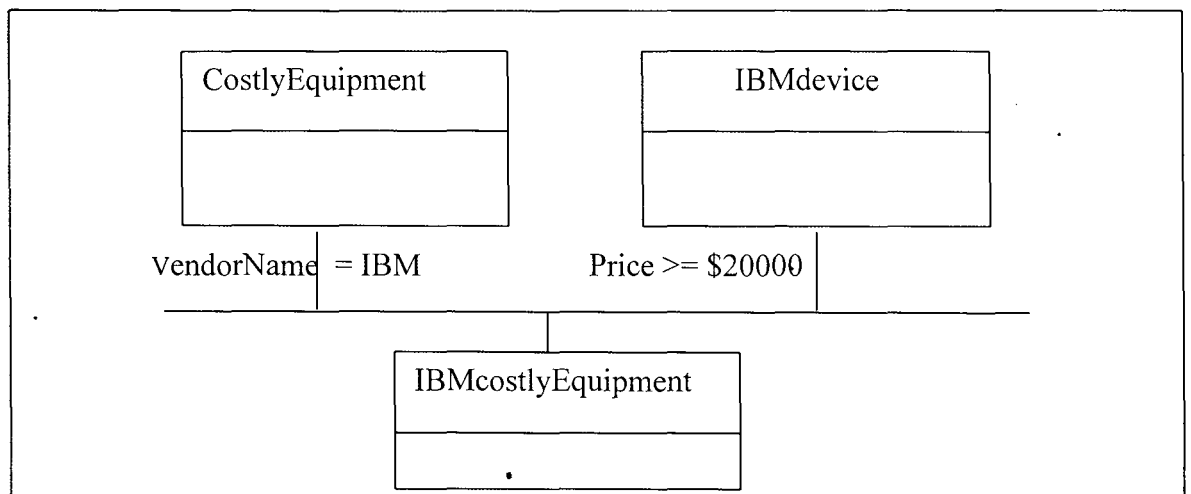
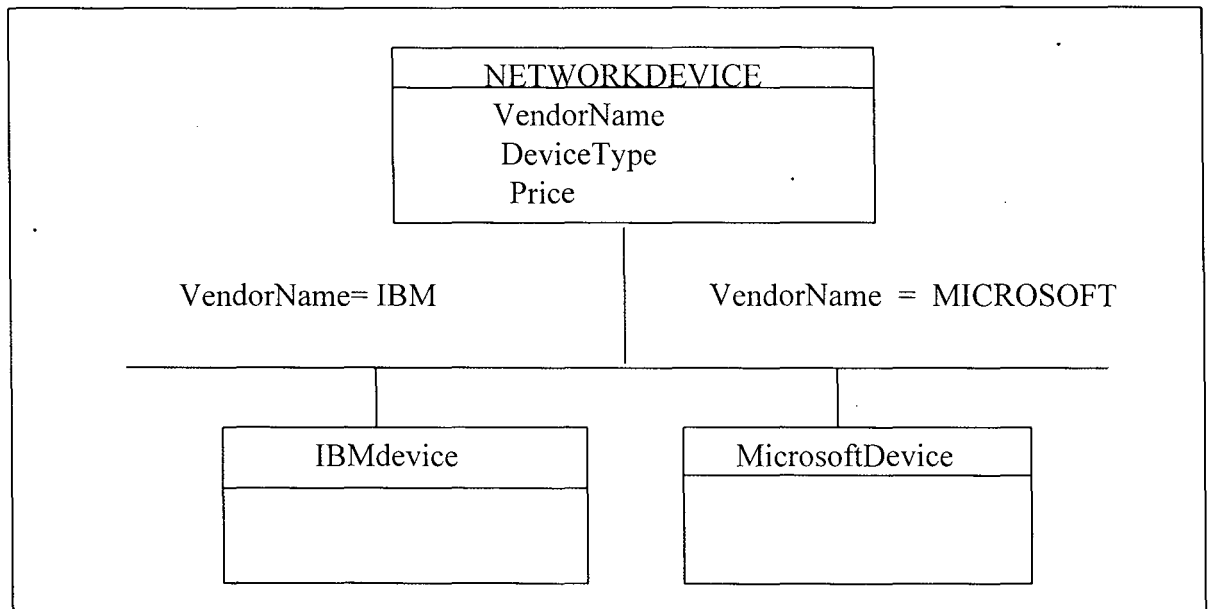
Specialisation principle of argument contravariance :- A subclass may expand the domain of an argument of a function inherited from a superclass .

Specialisation principle of result covariance :-A subclass may restrict the domain of the result of a function inherited from a superclass .

Specialisation principle of capsule fixing :- A subclass may mandate the use of an optional capsule inherited from a superclass .

After describing specialisation principles lets refer to multiple inheritance . It is the capability of an object class to inherit properties from more than one superclass . It is generally suitable when an ancestral class has been specialised at different times using different bases of specialisation . In multiple inheritance specialisation proceeding from the same superclass at two different times we have two different bases. of specialisation . Consider the example as shown in the figure





In this multiple derivation that basis of specialisation which is not used earlier is applied from each superclass. From the costlyequipment class the vendorName basis is being used while from the IBMdevice class the price basis is being used . In specialisation theory we permit the use of multiple inheritance when required . However as far as possible when subclassing only one basis property should be selected .

2.7 Normalization

Normalization is the technique used to avoid selective inheritance . It helps us find out the most appropriate class in which a particular property must be placed . The

normalization of a class hierarchy may occur either at architecture time or at operations time . However at operations time unrestricted redefinition of classes may cause instances to lose their instantiation associations with existing classes . So there are some restrictions with respect to normalizing a class hierarchy e.g. we can not change name of any concrete class . The extent of a concrete class can not change . The name of any ancestral class of a concrete class can not change . The membership of any ancestral class of a concrete class can not change and the set of properties accumulating in every concrete class through inheritance can not decrease . An advantage of normalization is that concrete classes tend to move to the bottom of the hierarchy and intermediate classes tend to stay abstract classes . It helps us in designing better hierarchies . The following are certain guidelines to design well structured object class hierarchies .

1. Use capsules to model minor variations of a class and specialisation to model major variations .
2. Avoid free specialisation . A formal basis of specialisation should be chosen at each level of subclassing .
3. Select only one basis of specialisation as far as possible
4. Ensure that each specialisation is disjoint.
5. Ensure that each specialisation is complete
6. push all concrete classes as far down in the hierarchy as possible.
7. Ensure that the hierarchy is normalized . By this it can be ensured that there is no need to apply selective inheritance.
8. Avoid multiple inheritance as far as possible . Try to find other ways to deal with the situation.

2.8 Aggregation

Aggregation is an important mechanism in object-oriented designing . It is a way of breaking down information in order to create complex objects by making an assembly of smaller objects . In object-oriented modelling this breakdown process is called decomposition . The reverse process i.e. assembling objects to create complex objects is called composition. The object formed on such a way is called aggregate object . The simpler objects which together form the aggregate are known as component objects. The aggregate is taken as a single object. Aggregation provides a powerful abstraction that facilitates the description and manipulation of all component objects . Aggregation specifies part-whole relationship between object classes The components are considered to be part and aggregate is whole . Aggregation relationships can be represented by aggregation hierarchies . An aggregation hierarchy is rooted at the aggregate object and proceeds downward to show its first order components . There may need of more than one aggregation hierarchy to describe the complete model of a single network.

Sometimes confusion occurs between inheritance hierarchy and aggregate hierarchy. The inheritance hierarchy describes the specializes from association between classes . Its goal is to describe taxonomies so that we can categorize objects . The categorisation of an object does not affect the operation of the object itself. But aggregation hierarchy defines the compositions so that we can assemble the objects to propagate the effects of operations on one object to another. It actually affects the operations of the objects. The same object class may appear twice in the aggregation hierarchy but every object class can occur exactly once in the inheritance hierarchy . These two hierarchies are complementary to each other . None of them is sufficient . They are both required together . An aggregation hierarchy is like an assembly blueprint which tells us how to put them together .The inheritance hierarchy helps us in finding the parts we need .The definition of aggregation is not concerned with individual properties of the components.

There are two types of aggregation exclusion aggregation and inclusive aggregation . In exclusion aggregation each specific component instance can belong to only one

aggregation instance while in inclusive aggregation some component may belong to more than one aggregate instance e.g. when we say that a hubcard component is a part of a wiringhub aggregate we imply that no other wiringhub object may possess the same hubcard component . This is an example of an exclusive aggregation . On the other hand suppose that a digitalvideoclip component object is a part of multimediadocument aggregate object . This does not prevent the same digitalvideoclip object from also being a component of a different multimediadocument object . This is an example of inclusive aggregation . There exists some confusion aggregation and capsules . It should be kept in mind that capsules are never instantiated whereas component objects are . Further a class may contain more than one instance of the same component whereas it can contain each capsule only once.

In order to describe how many times a component is required to form an aggregate , the concept of component multiplicity is introduced . It allows us to specify the number of components which may be contained in the aggregate . Component multiplicity is shown in the form of an interval $[N,M]$ where N and M are non-negative integers . This interval infers that there are at least N and at most M instances of the component to form the aggregate . Component multiplicity may also be specified by a discrete set of enumerated integers . It is also possible to define component multiplicity in terms of attribute-defined expressions . In fact attribute-defined expression is actually the most general and most satisfactory way of specifying the component multiplicity . If the aggregate class has been specified correctly , it is usually possible to select one its capacity related attributes to specify component multiplicity . In every multiplicity specification where the network model seems to require a non-trivial integer constant for a limit it is better to look closely for any capacity related attribute of the aggregate class and examine it for possible use in an attribute defined expression for the component multiplicity.

An aggregate class and its component classes both have a place in the inheritance hierarchy . Therefore each can specialise into a number of descendants classes . Further the descendants of the component object classes have inherited its

composition relationship with the aggregate object class and the descendants of the aggregate object class have inherited its decomposition relationship with the component object class . There are certain aggregation principles .

Aggregation principle of composition inheritance :- A descendant of a component class inherits the composition relationship of an ancestral with the aggregate class .

Aggregation principle of decomposition inheritance :- A descendant of an aggregate class inherits the decomposition relationship of an ancestral class with the component class .

Aggregation principle of monotonic aggregation inheritance :- A descendant class may not cancel any composition or decomposition relationship defined for an ancestral class .

Aggregation principle of component multiplicity inheritance :- A descendant of an aggregate class inherits by default the component multiplicity of each decomposition relationship inherited from an ancestral class that it may modify in permissible ways .

Specialisation principle of component addition :- A descendant class may add a decomposition relationship with a new component class .

Specialisation principle of aggregate addition :- a descendant class may add a composition relationship with a new aggregate class .

Specialisation principle of component multiplicity restriction :- A descendant class may restrict the multiplicity domain of a decomposition relationship inherited from an ancestral class .

Aggregation principle of component multiplicity inheritance :- A descendant of an aggregate class inherits by default the component multiplicity of each decomposition relationship inherited from an ancestral class, which it may restrict .

Specialisation principle of component specialisation :- A descendant of an aggregate class may restrict an inherited decomposition relationship to selected descendants of the component class .

Specialisation principle of aggregate specialisation :- A descendant of a component class may restrict an inherited composition relationship to selected descendants of the aggregate class .

Aggregation principle of specialized component inheritance :- A descendant of an aggregate class has-as-a part at least one descendant of a mandatory component class , if any specialisation of the component class is complete .

Aggregation principle of specialised aggregate inheritance :- A descendant of a component class is-a-part-of at least one descendant of a mandatory aggregate class if any specialisation of the aggregate class is complete .



TH-7889



CHAPTER 3

TOKEN RING SPECIFICATION & DESIGN

3.1 Definitions

Abort sequence:- A sequence that terminates the transmission of a frame prematurely.

fill. A bit sequence which may be either 0 bits, 1 bits, or any combination thereof.

frame. A transmission unit that carries a protocol data unit (PDU) on the ring.

logical link control (LLC). That part of the data link layer that supports media independent data link functions, and uses the services of the medium access control sublayer to provide services to the network layer.

medium. The material on which the data may be represented. Twisted pairs, coaxial cables, and optical fibers are examples of media.

medium access control (MAC). The portion of the IEEE 802 data station that controls and mediates the access to the ring.

medium interface connector (MIC). The connector between the station and trunk coupling unit (TCU) at which all transmitted and received signals are specified.

monitor. The monitor is that function that recovers from various error situations. It is contained in each ring station; however, only the monitor in one of the stations on a ring is the active monitor at any point in time. The monitor function in all other stations on the ring is in standby mode.

multiple frame transmission. A transmission where more than one frame is transmitted when a token is captured.

network management (NMT). The conceptual control element of a station which interfaces with all of the layers of the station and is responsible for the setting and resetting of control parameters, obtaining reports of error conditions and determining if the station should be connected to or disconnected from the medium.

repeater. A device used to extend the length, topology, or interconnectivity of the transmission medium beyond that imposed by a single transmission segment.

ring latency. In a token ring medium access control system, the time required for a

signal to propagate once around the ring. The ring latency time includes the signal propagation delay through the ring medium plus the sum of the propagation delays through each station connected to the token ring.

service data unit (SDU). Information delivered as a unit between adjacent entities which may also contain a PDU of the upper layer.

Station (or data station). A physical device that may be attached to a shared medium local area network for the purpose of transmitting and receiving information on that shared medium. A data station is identified by a destination address.

Token. The symbol of authority that is passed between stations using a token access method to indicate which station is currently in control of the medium.

Transmit. The action of a station generating a frame, token, abort sequence, or fill and placing it on the medium to the next station. In use, this term contrasts with repeat.

Trunk cable. The transmission cable that interconnects two trunk coupling units.

trunk coupling unit (TCU). A physical device that enables a station to connect to a trunk cable. The trunk coupling unit contains the means for inserting the station into the ring or. Conversely bypassing the station.

3.2 General Description

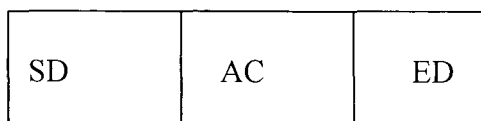
A token ring consists of a set of stations serially connected by a transmission medium. Information is transferred sequentially, bit by bit, from one active station to the next. Each station generally regenerates and repeats each bit and serves as the means for attaching one or more devices (terminals, work-stations) to the ring for the purpose of communicating with other devices on the network. A given station (the one that has access to the medium) transfers information onto the ring, where the information circulates from one station to the next. The addressed destination station(s) copies the information as it passes. Finally, the station that transmitted the information effectively removes the information from the ring. A station gains the right to transmit its information onto the medium when it detects a token passing on the medium. The token is a control signal comprised of a unique signaling sequence

that circulates on the medium following each information transfer. Any station, upon detection of an appropriate token, may capture the token by modifying it to a start-of-frame sequence and appending appropriate control and status fields, address fields, information field, frame-check sequence and the end-of-frame sequence. At the completion of its information transfer and after appropriate checking for proper operation, the station initiates a new token, which provides other stations the opportunity to gain access to the ring. A token holding timer controls the maximum period of time a station shall use (occupy) the medium before passing the token. Multiple levels of priority are available for independent and dynamic assignment depending upon the relative class of service required for any given message. For example, synchronous real-time voice), asynchronous (interactive), immediate (network recovery). The allocation of priorities shall be by mutual agreement among users of the network. Error detection and recovery mechanisms are provided to restore network operation in the event that transmission errors or medium transients (for example, those resulting from station insertion or removal) cause the access method to deviate from normal operation. Detection and recovery for these cases utilize a network monitoring function that is performed in a specific station with back-up capability in all other stations that are attached to the ring.

3.3 Formats and Facilities

Formats. There are two basic formats used in token rings: tokens and frames. In the following discussion, the figures depict the formats of the fields in the sequence they are transmitted on the medium, with the left-most bit or symbol transmitted first.

Token Format

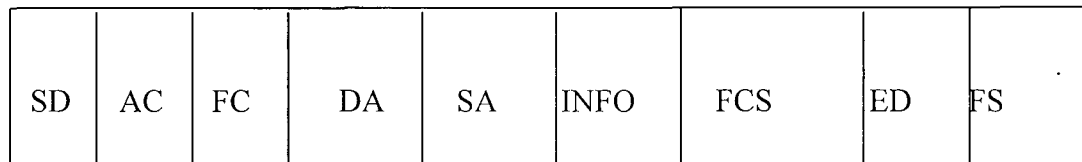


SD = Starting Delimiter (1 octet)

AC = Access Control (1 octet)

ED = Ending Delimiter (1 octet)

Frame Format



SD : Starting Delimiter (1 octet)

AC : Access Control (1 octet)

FC : Frame Control (1 octet)

DA : Destination Address (2 or 6 octets)

INFO Information (0 or more octets)

FCS :Frame Check Sequence (4 octets)

EFS End-of-Frame Sequence

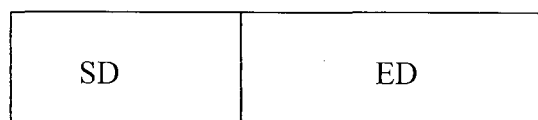
ED : Ending Delimiter (1 octet)

SA = Source Address (2 or 6 octets)

FS : Frame Status (1 octet)

The frame format shall be used for transmitting both medium access control (MAC) and logical link control (LLC) messages to the destination station. It may or may not have an information INFO field.

ABORT -SEQUENCE



This sequence shall be used for the purpose of terminating the transmission of a frame prematurely. The abort sequence may occur any-where in the bit stream; that is, receiving stations shall be able to detect an abort sequence even if it does not occur on octet boundaries.

Fill. When a station is transmitting (as opposed to repeating), it shall transmit fill preceding or following frames, tokens, or abort sequences to avoid what would otherwise be an inactive or indeterminate transmitter state.

Fill may be either 0 or 1 bits or any combination thereof and may be *any number* of bits in length, within the constraints of the token holding timer.

Field Descriptions. The following is a detailed description of the individual fields in the tokens and frames.

Starting Delimiter (SD)

J	K	0	J	K	0	0	0
---	---	---	---	---	---	---	---

J = non-data-J

K = non-data-K

0 = binary zero

Access Control (AC)

P	P	P	T	M	R	R	R
---	---	---	---	---	---	---	---

PPP = priority bits

T = token hit

M = monitor bit

RRR = reservation bits

Priority Bits. The priority bits shall indicate the priority of a token and, therefore, which stations are allowed to use the token. In a multiple-priority system, stations use different priorities depending on the priority of the PDU to be transmitted. The eight levels of priority increase from the lowest (000) to the highest (111) priority. For purposes of comparing priority values, the priority shall be transmitted most significant bit first; for example, 110 has higher priority than 011 (left-most bit transmitted first).

Token Bit. The token bit is a 0 in a token and a 1 in a frame. When a station with a PDU to transmit detects a token which has a priority equal to or less than the PDU to be transmitted, it may change the token to a start-of-frame sequence and transmit the PDU.

Monitor Bit. The monitor bit is used to prevent a token whose priority is greater than 0 or any frame from continuously circulating on the ring. If an active monitor detects a frame or a high priority token with the monitor bit equal to 1, the frame or token is aborted. This bit shall be transmitted as 0 in all frames and tokens. The active monitor inspects and modifies this bit. All other stations shall repeat this bit as received.

Reservation Bits. The reservation bits allow stations with high priority PDUs to request (in frames or tokens as they are repeated) that the next token be issued at the requested priority. The eight levels of reservation increase from 000 to 111. For purposes of comparing reservation values, the reservation shall be transmitted most significant bit first; for example, 110 has higher priority than 011 (left-most bit transmitted first).

Frame Control (FC)

F	F	Z	Z	Z	Z	Z	Z
---	---	---	---	---	---	---	---

The FC field defines the type of the frame and certain MAC and information frame functions.

Frame-Type Bits. The frame-type bits shall indicate the type of the frame as follows:

00 = MAC frame (contains an MAC PDU)

01 = LLC frame (contains an LLC PDU)

1x = undefined format (reserved for future use)

Medium Access Control (MAC) Frames. If the frame-type bits indicate a MAC frame, all stations on the ring shall interpret and, based on the finite state of the station, act on the *ZZZZZZ* control bits.

Logical Link Control (LLC) Frames. If the frame-type bits indicate an LLC frame, the *ZZZZZZ* bits are designated as *rrYYYY*. The *rrr* bits are reserved and shall be transmitted as 0's in all transmitted frames and ignored upon reception. The *YYY* bits may be used to carry the priority (*Pm*) of the PDU from the source LLC entity to the target LLC entity or entities. Note that *P* (the priority in the access control [AC] field of a frame) is less than or equal to *Pm* when the frame is transmitted onto the ring.

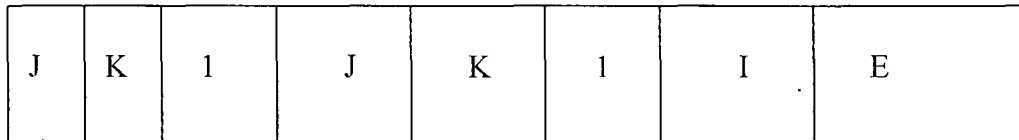
Destination and Source Address (DA and SA) Fields- Each frame shall contain two address fields: the destination (station) address and the source (station) address, in that order. Addresses may be either 2 or 6 octets in length; however, all stations of a specific LAN shall have addresses of equal length.

Information (INFO) Field. The information field contains 0, 1, or more octets that are intended for MAC, NMT, or LLC. Although there is no maximum length specified for the information field, the time required to transmit a frame may be no greater than the token holding period that has been established for the station.

The format of the information field is indicated in the frame-type bits of the FC field. The frame types defined are MAC frame and LLC frame.

Frame-Check Sequence (FCS). The FCS shall be a 32-bit sequence. The FCS shall be transmitted commencing with the coefficient of the highest term

Ending Delimiter (ED)



J = non-data-J

K = non-data-K

1 = binary one

I = intermediate frame bit

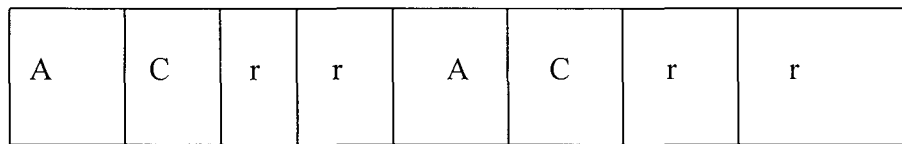
E = error-detected bit

The transmitting station shall transmit the delimiter as shown. Receiving stations shall consider the ending delimiter (ED) valid if the first six symbols J K 1 J K 1 are received correctly.

Intermediate Frame Bit (I Bit). To indicate that this is an intermediate (or first) frame of a multiple frame transmission, the I bit shall be transmitted as 1. An I bit of 0 indicates the last or only frame of the transmission.

Error-Detected Bit (E Bit). The error-detected bit (E) shall be transmitted as 0 by the station that originates the token, abort sequence, or frame. All stations on the ring check tokens and frames for errors (for example, FCS error, non-data symbols: see 4.2.1). The E bit of tokens and frames that are repeated shall be set to 1 when a frame with error is detected; otherwise the E bit is repeated as received.

Frame Status (FS).



A = address-recognized bits

C = frame-copied bits

r = reserved bits

reserved bits. These bits are reserved for future standardization. They shall be transmitted as 0's; however, their value shall be ignored by the receivers.

Address-Recognized (A) Bits and Frame-Copied (C) Bits. The A and C bits shall be transmitted as 0 by the station originating the frame. If another station recognizes the destination address as its own address or relevant group address, it shall set the A bits to 1. If it copies the frame (into its receive buffer), it shall also set the C bits to 1. This allows the originating station to differentiate among three conditions:

1. Station non-existent/non-active on this ring
2. Station exists but frame not copied
3. Frame copied

Claim Token MAC Frame (CL__TK). When a station determines that there is no active monitor operating on the ring, it shall send claim token frames and inspect the source address of the claim token MAC frames it receives. If the SA matches its own (MA) address it has claimed the token and shall enter active monitor mode and generate a new token.

Duplicate Address Test MAC Frame (DAT). This frame is transmitted with DA = MA as part of the initialization process. If the frame returns with the A bits set to 1, it indicates that there is another station on the ring with the same address. If such an

event occurs, the station's network manager is notified and the station returns to bypass state. A station that copies a DAT frame will ignore it.

Active Monitor Present MAC Frame (AMP). This frame is transmitted by the active monitor. It shall be queued for transmission following the successful purging of the ring or following the expiration of the TAM. Any station that receives this frame shall reset its TSM whose A and C bits equal 0, the TQP is reset. When timer TQP expires, an SMP PDU shall be queued for transmission.

Beacon MAC Frame (BCN). This frame shall be sent as a result of serious ring failure (for example, broken cable, jabbering station, etc). It is useful in localizing the fault. The immediate upstream station is part of the failure domain about which the beacon is reporting. Therefore the address of the upstream station that was previously recorded is included in the MAC INFO field.

Purge MAC Frame (PRG). This frame is transmitted by the active monitor. It shall be transmitted following claiming the token or to perform reinitialization of the ring following the detection of an M bit set to I or the expiration of timer TVX.

Timer, Return to Repeat (TRR). Each station shall have a timer TRR to ensure that the station shall return to Repeat State. TRR shall have a value greater than the maximum ring latency. The maximum ring latency consists of the signal propagation delay around a maximum-length ring plus the sum of all station latencies.

Timer, Holding Token (THT). Each station shall have a timer THT to control the maximum period of time the station may transmit frames after capturing a token. A station may initiate transmission of a frame if such transmission can be completed before timer THT expires.

Timer, Queue PDU(TQP):- Each station shall have a timer TQP for the purpose of timing the enqueueing of an SMP PDU after reception of an AMP or SMP frame in which the A and C bits were equal to 0. The default time-out value of TQP is 10 ins.

Timer, Valid Transmission (TVX). Each station shall have a timer TVX ,which is used by the active monitor to detect the absence of valid transmissions. The time-out value of TVX shall be the sum of the time-out value of THT plus the time-out value of TRR.

Timer, No Token (TNT). Each station shall have a timer TNT to recover from various token-related error situations. TNT shall have a time-out value equal to TRR plus n times THT (where n is the maximum number of stations on the ring).

Timer, Active Monitor (TAM). Each station shall have a timer TAM which is used by the active monitor to stimulate the enqueueing of an AMP PDU for transmission. The default time-out value of timer TAM shall be 3 s.

Timer, Standby Monitor (TSM). Each station shall have a timer TSM which is used by the stand-by monitor(s) to assure that there is an active monitor on the ring and to detect a continuous stream of tokens. The default time-out value of timer TSM shall be 7 s.

Flags. Flags are used to *remember* the occurrence of a particular event. They shall be set when the event occurs. The flags used are:

I Flag: A flag which is set upon receiving an ED with the I bit equal to 0.

SFS FLAG :A flag which is set upon receiving an SFS sequence.

MA Flag: A flag which is set upon receiving an SA which is equal to the station's address.

Latency Buffer. The latency buffer serves two purposes. The first is to ensure that there are at least 24 bits of latency in the ring. The second is to provide phase jitter compensation. The token management is structured so that only one latency buffer

shall be active in a normally functioning ring and is provided by the active monitor in the ring.

3.4 Token Ring Protocols

Frame Transmission. Access to the physical medium (the ring) is controlled by passing a token around the ring. The token gives the downstream (receiving) station (relative to the station passing the token) the opportunity to transmit a frame or a sequence of frames. Upon request for transmission of an LLC PDU or NMT PDU, MAC prefixes the PDU with the appropriate FC, DA, and SA fields and enqueues it to await the reception of a token that may be used for transmission.

Such a token has a priority less than or equal to the priority of the PDU(s) that is to be sent. Upon queuing the PDU for transmission and prior to receiving a usable token, if a frame or an unusable token is repeated on the ring, the station requests a token of appropriate priority in the RRR bits of the repeated AC field. Upon receipt of a usable token, it is changed to a start-of-frame sequence by setting the token bit. At this time, the station stops repeating the incoming signal and begins transmitting a frame. During transmission, the FCS for the frame is accumulated and appended to the end of the information field.

Token Transmission. After transmission of the frames has been completed, the station checks to see if the station's address has returned in the SA field, as indicated by the MA__FLAG. If it has not been seen, the station transmits fill until the MA_ FLAG is set, at which time the station transmits a token.

Stripping. After transmission of the token, the station will remain in transmit state until all of the frames that the station originated are removed from the ring. This is done to avoid unnecessary recovery action that would be caused if a frame were allowed to continuously circulate on the ring.

Frame Reception. Stations, while repeating the incoming signal stream, check it for frames they should copy or act upon. If the frame-type bits indicate a MAC frame, the control bits are interpreted by all stations on the ring. In addition, if the frame's DA field matches the station's individual address, relevant group address, or broadcast address, the FC, DA, SA, INFO, and FS fields are copied into a receive buffer and subsequently forwarded to the appropriate sublayer.

Priority Operation. The priority bits (PPP) and the reservation bits (RRR) contained in the access control (AC) field work together in an attempt to match the service priority of the ring to the highest priority PDU that is ready for transmission on the ring. These values are stored in registers as Pr and Rr. The current ring service priority is indicated by the priority bits in the AC field, which is circulated on the ring. The priority mechanism operates in such a way that fairness (equal access to the ring) is maintained for all stations within a priority level. This is accomplished by having the same station that raised the service priority level of the ring (the stacking station) return the ring to the original service priority. The priority operation is explained as follows: When a station has a priority (a value greater than zero) PDU (or EDU's) ready to transmit, it requests a priority token. This is done by changing the reservation bits (RRR) as the station repeats the AC field. If the priority level (Pm) of the PDU that is ready for transmission is greater than the RRR bits, the station increases the value of RRR field to the value Pm. If the value of the RRR bits is equal to or greater than Pm, the reservation bits (RRR) are repeated unchanged. After a station has claimed the token, the station transmits PDUs that are at or above the present ring service priority level until it has completed transmission of those PDUs or until the transmission of another frame could not be completed before timer TNT expires). The priority of all of the PDUs that are transmitted should be at the present ring service priority value. The station will then generate a new token for transmission on the ring. If the station does not have additional PDUs to transmit that have a priority (Pm) or does not have a reservation request (as contained in register Rr) neither of which is greater than the present ring service priority (as contained in register Pr), the token is transmitted with its priority at the present ring service priority and the reservation bits RRR) at the greater of Rr

or P_m and no further action taken. However, if the station has a PDI ready for transmission or a reservation request (E_r), either of which is greater than the present ring service priority, the token is generated with its priority at the greater of P_m or R_r and its reservation bits (RRR) as 0. Since the station has raised the service priority level of the ring, the station becomes a stacking station and, as such, stores the value of the old ring service priority as S_r and the new ring service priority as S_x . (These values will be used later to lower the service priority of the ring when there are no PDUs ready to transmit on the ring whose P_m is equal to or greater than the stacked S_x .) Having become a stacking station, the station claims every token that it receives that has a priority (PPP) equal to its highest stacked transmitted priority (S_x) in order to examine the RRR bits of the AC field for the purpose of raising, maintaining, or lowering the service priority of the ring. The new token is transmitted with its PPP bits equal to the value of the reservation bits (RRR) but no lower than the value of the highest stacked received priority (S_r), which was the original ring priority service level. If the value of the new ring service priority (PPP equal to R_r) is greater than S_r the HER bits are transmitted as 0, the old ring service priority contained in S_x is replaced with a new value S_x equal to R_r , and the station continues its role as a stacking station. However, if the R_r value is equal to or less than the value of the highest stacked received priority (S_r) the new token is transmitted at a priority value of the S_r , both S_x and S_r are removed (popped) from the stack, and if no other values of S_x and S_r are stacked, the station discontinues its role as a stacking station. The frames that are transmitted to initialize the ring have a PPP field that is equal to 0. The receipt of a PPP field whose value is less than a stacked S_x will cause any S_x or S_r values that may be stacked to be cleared in all stations on the ring.

Beaconing and Neighbor Notification. When a hard failure is detected in a token ring, its cause must be isolated to the proper failure domain so that recovery actions can take place. The failure domain consists of

1. the station reporting the failure (the beaconing station)
2. the station upstream of the beaconing station
3. the ring medium between them

To do accurate problem determination, all elements of the failure domain must be known at the time that the failure is detected. This implies that at any given time, each station should know the identity of its upstream neighbor station. A process for obtaining this identity known as Neighbor Notification is described below. Neighbor Notification has its basis in the address-recognized and frame-copied bits (the A and C bits) of the FS field. These bits are transmitted as 0's. If a station recognizes the destination address of the frame as one of its own, the station sets the A bits to 1 in the passing frame. If a station also copies the *frame*, then the C bits are also set to a 1. When a frame is broadcast to all stations on a ring, the first station downstream of the broadcaster will see that the A and C bits are all 0's. Since a broadcast frame will have its destination address recognized by all of the stations on the ring, the first station downstream will, in particular, set the A bits to 1. All stations further downstream will, therefore, not see the A and C bits as all 0's. This process continues in a circular, daisy-chained fashion to let every station know the identity of its upstream neighbor. The monitor begins Neighbor Notification by broadcasting the active monitor present (AMP) MAC frame. The station immediately downstream from it takes the following actions:

1. resets its timer TSM, based on seeing the AMP value in the FC field;
2. if possible, copies the broadcast AMP MAC frame and stores the upstream station's identity in an upstream neighbor's address (UNA) memory location;
3. sets the A bits (and C bits if the frame was copied) of the passing frame to 1's
4. at a suitable transmit opportunity, broadcasts a similar standby monitor present (SMP) MAC frame.

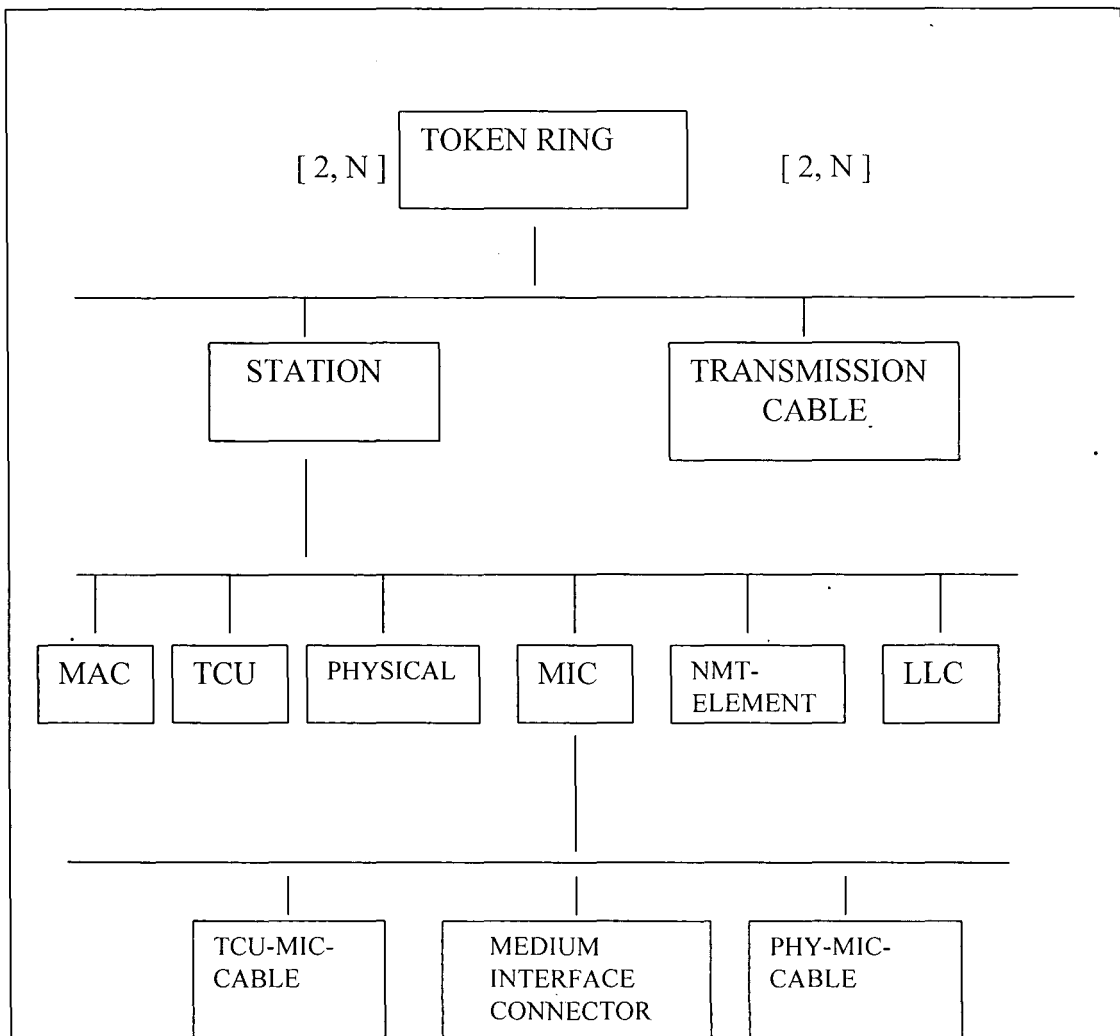
One by one, each station receives an SMP frame with the A and C bits set to 0's, stores its UNA, and continues the process by broadcasting such a frame itself. Since the AMP frame must pass each station on a regular basis (the active monitor present MAC frame sent by the monitor), the continuous transmission of tokens onto a ring can be detected. In addition to the timer TAM in the active monitor, each standby

station has a timer TSM that is reset each time an AMP MAC frame passes. If timer TSM expires, that standby monitor station begins transmitting claim token frames. From the received symbols MAC detects various types of input data, such as tokens, MAC frames, and LLC information frames. In turn, MAC stores values, sets flags, and performs certain internal actions as well as generating tokens, frames, or fill, or flipping bits and delivering them to the PHY layer in the form of a serial stream of the 0, 1, J, and K symbols. For the purpose of accumulating the FCS and storing the contents of a frame, J and K symbols that are not part of the SD or ED shall be interpreted as 1 and 0 bits, respectively.

3.5 Carving out the design of token ring network

The token ring Network consists of the following distinguished components. Token Ring as a whole, stations, transmission cable, transmission control unit , medium interface cable , medium interface connector , network management and network administrator .

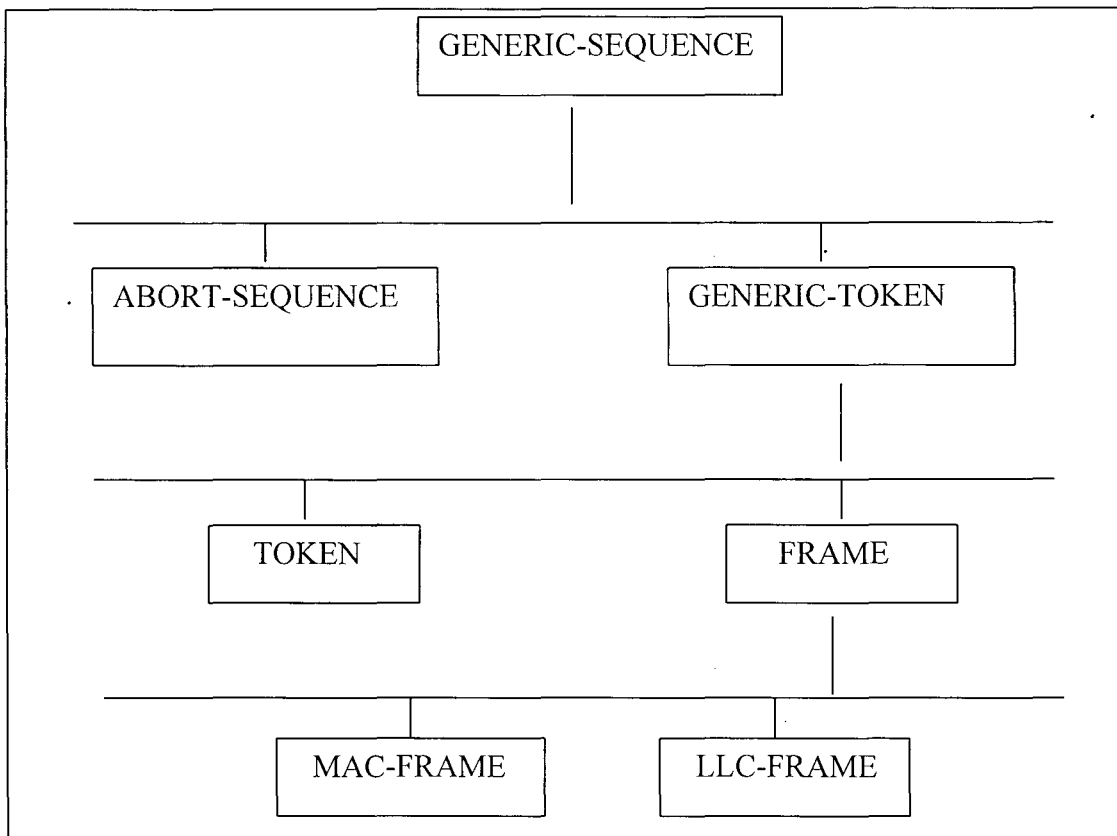
AGGREGATION:- Considering Token Ring Network as a whole , the stations are its main components . For a network the number of stations may be at least two or more but less than a fixed number . Corresponding to a station there is a medium interface cable and TCU (Transmission control unit). Therefore they will be shown as part of station class. The medium interface cable consists of medium interface connector, tcu-mic-cable and physical-layer-mic-cable. The working of station is layered in physical-layer, macSublayer and nmtElement. The aggregation hierarchy can be shown as follows



AGGREGATION HIERARCHY OF TOKEN RING

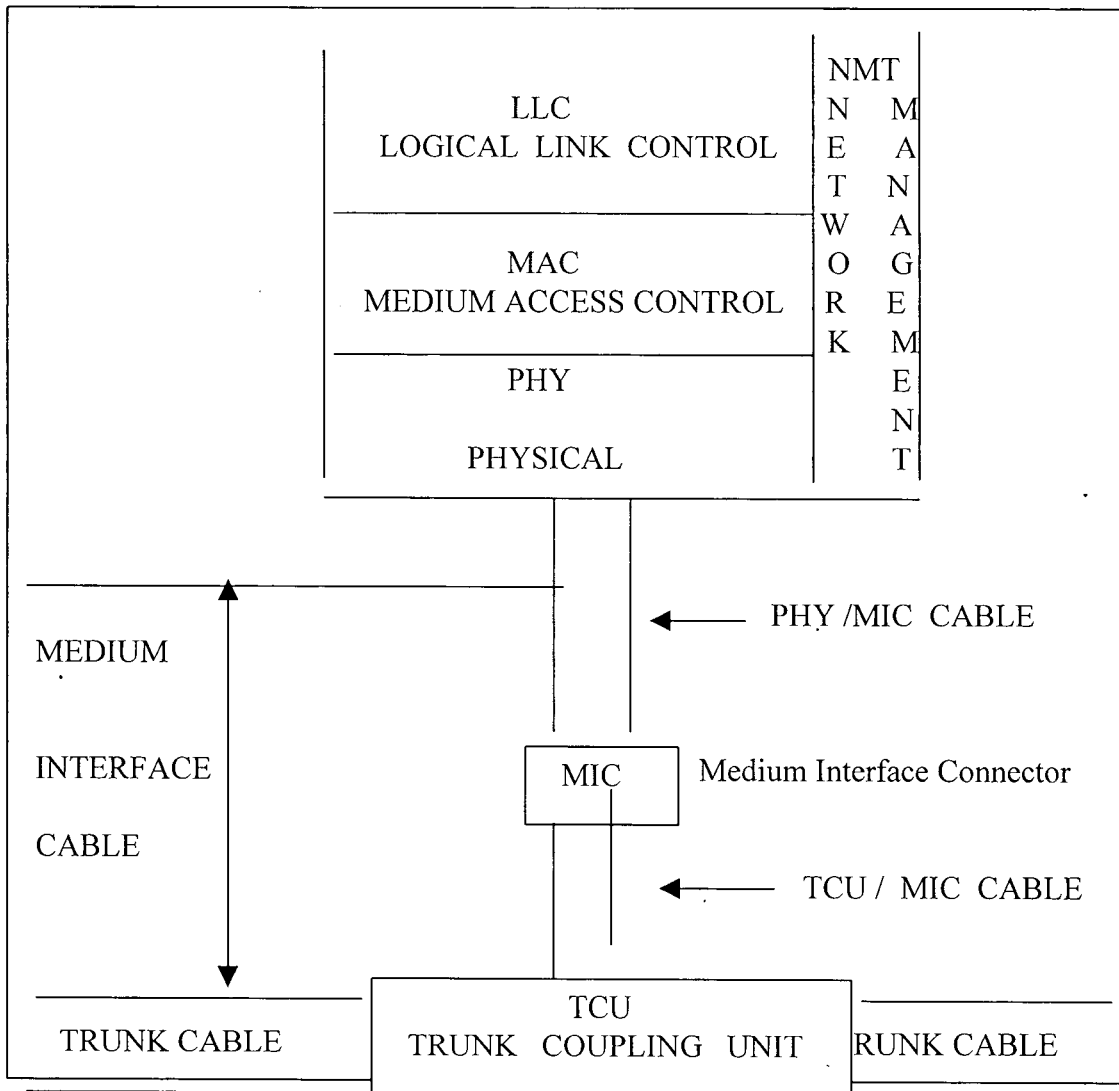
INHERITANCE :- Token and frame are two important components which are transmitted over the network. In the access-control if the token-bit is zero it means it

is token otherwise it is a frame. So the token bit is chosen as a basis of specialization. A generic-token is the abstract class from which token and frame classes are derived using token bit as a basis of specialization . There is another sequence termed as abort-sequence consisting of start-delimiter and ending-delimiter only . An abstract class called as generic-sequence is created for the purpose of specification reuse. There are different types of frames known as mac-frame and llc-frame These are differentiated on the basis of frame-type-bits. So it is used as a basis of specialization. The specialization principle of attribute domain restriction is applied here .The inheritance tree created for the token and frames is as shown below.



Identifying the attributes and behavior:- The abort-sequence , token and frame all contain the start-delimiter and ending-delimiter . The start-delimiter and ending-delimiter are sequence of bits . They have no actual instance and they are not independently instantiated . So they are treated as capsules . Further access-control is also not instantiated independently so it is also treated as capsule . Generic-sequence is developed as an abstract class for the purpose of specification reuse. This class will have the functionality of checking start and ending-delimiter Another abstract class called generic-token class is required to be created as the capsule access-control is added providing the access and set functions .The frame class needs to have source and destination address, frame-check-sequence and info part. In addition it will have

frame-control and frame-status. The class abort-sequence is meant to stop the transmission so it will have a function abort-transmission.



Layering:- The functioning of the station class is decomposed into layers viz. Physical-layer, mac-sublayer, llc-sublayer and nmtElement. Access and transmission is controlled by the mac-sublayer. These layers are arranged on the basis of layer principles i.e. lower layers provide services to upper layers through service access points. The arrangement of layers is as shown in the figure.

The logical link control layer accesses services from medium access control sublayer in order to services to higher layers which are not shown in the figure . In this dissertation layering is kept limited upto the llc sublayer . Medium access control sublayer accesses services provided by physical layer and interacts with nmtElement class for proper network management . It provides all information like status changes to nmtElement . Physical layer also interacts with nmtElement so that functions like insertion or removal of a station can be controlled automatically by network management without human intervention.

SERVICES :- The following are the interactions between logical link control layer and medium access control sublayer

MA-DATA-request.

MA-DATA-indication

MA-DATA-confirmation

MA-DATA-request:- This primitive is used by llc-sublayer to request macSublayer to send a service data unit to another llc-sublayer.

MA-DATA-indication :-It will be used by macSublayer to send a macSublayer frame to llc-sublayer to indicate the arrival of an llc frame at the local macSublayer entity .

MA-DATA-confirmation :- This primitive is used to provide an appropriate response to the llc-sublayer signifying the success or failure of the request

PHY to MAC Service. The services provided by the PHY layer allow the local MAC sublayer entity to exchange MAC data units with peer MAC sub-layer entities.

Interactions. The following primitives are defined for the MAC sub-layer to request service from the PHY layer:

PH-DATA-request

PH-DATA-indication

PH-DATA-confirmation

PH-DATA-request. This primitive defines the transfer of data from a local MAC sublayer entity to the station's PHY layer. The MAC sublayer shall send the PHY layer a PH-DATA-request every time the MAC sublayer has a symbol to output.

PH-DATA-indication. This primitive defines the transfer of data from the PHY layer to the MAC sublayer entity. The PHY layer shall use this primitive to send the MAC sublayer a PH-DATA-indication every time the PHY layer decodes a symbol.

PH-DATA.confirmation. This primitive will be used by physical layer to provide an appropriate response to the MAC sublayer PH-DATA-request signifying the acceptance of a symbol specified by the PH-DATA-request and willingness to accept another symbol.

MAC-NMT Services:- This interface is used by NMT to monitor and control operations of the MAC sublayer. The following primitives are defined for the NMT to request service from the MACsublayer.

MA-INITIALIZE-PROTOCOL-request

MA-INITIALIZE-PROTOCOL-confirmation

MA-CONTROL-request

MA-STATUS-indication

MA-NMT-DATA-request

MA-NMT-DATA-indication

MA-NMT-DATA-confirmation.

MA-INITIALIZE-PROTOCOL-request:-This primitive is used by NMT to reset the MAC sublayer and optionally to change the operational parameters of the MAC sublayer.

MA-INITIALIZE-PROTOCOL-confirmation:-This primitive is used by the MAC sublayer to inform NMT that the MA-INITIALIZE-PROTOCOL-request primitive is complete.

MA-STATUS-indication:-This primitive is used by the MAC sublayer to inform NMT of errors and significant status changes.

MA-NMT-DATA-request:-This primitive shall be generated by the NMT entity whenever data must be transferred to one or more NMT entities.

MA-NMT-DATA-indication:-This primitive defines the transfer of data from the MAC sublayer entity. The MA-NMT-DATA-indication primitive shall be generated by the MAC sublayer entity to the NMT entity to indicate the arrival of a MAC frame at the local MAC sublayer entity.

MA-NMT-DATA-confirmation:-This primitive shall provide an appropriate response to the NMT's MA-NMT-DATA-request primitive signifying the success or failure of the request.

PHY-NMT SERVICES :-The following primitives are defined for the NMT to request services from the physical layer.

PH-CONTROL-request

PH-STATUS-indication

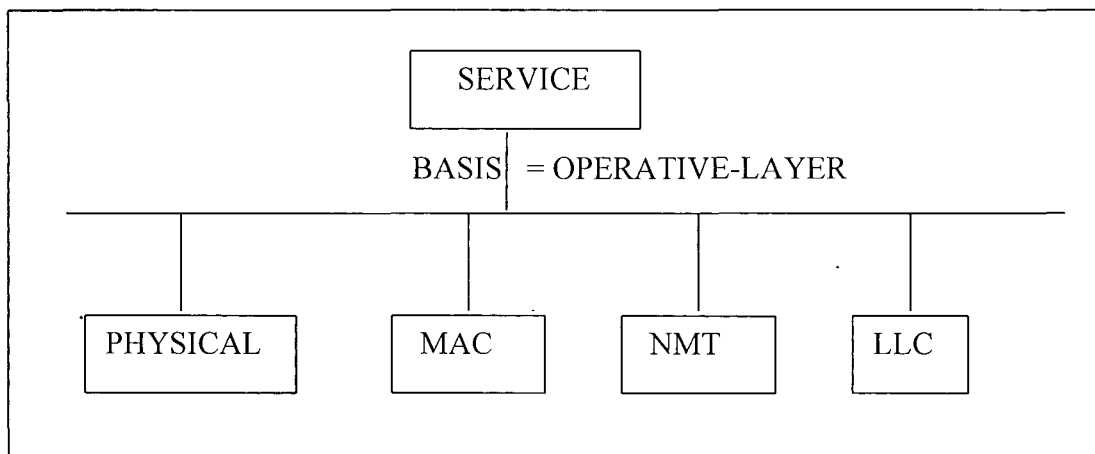
PH-CONTROL-request :- This primitive shall be generated by NMT to request the physical layer to insert or remove itself to/from the ring.

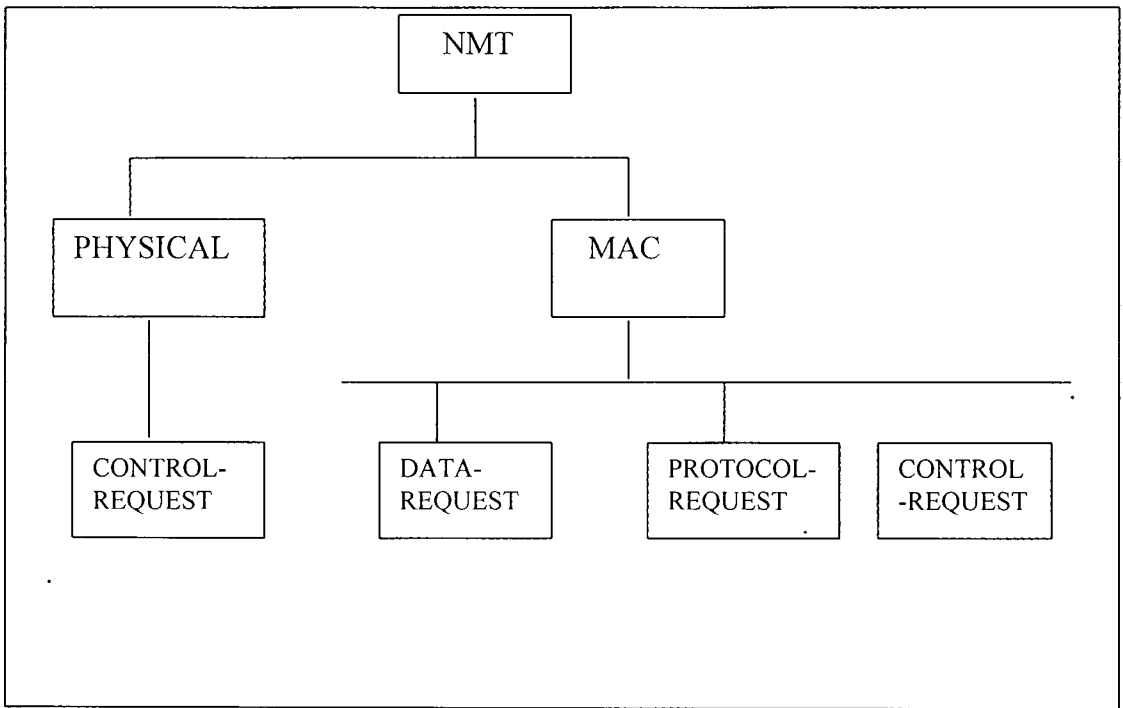
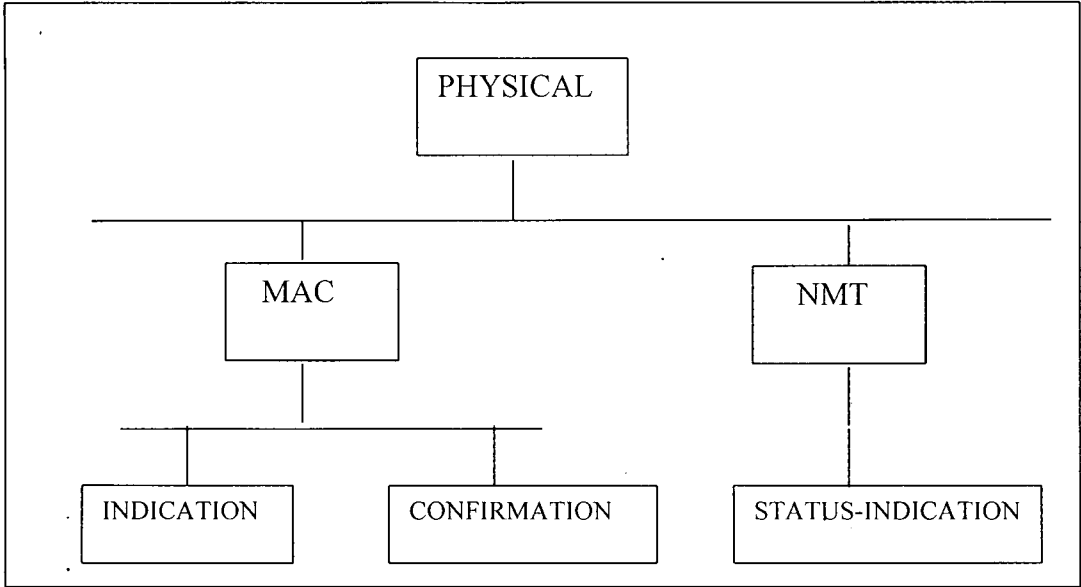
PH-STATUS-indication:-This primitive is used by the physical layer to inform NMT of errors and significant status changes.

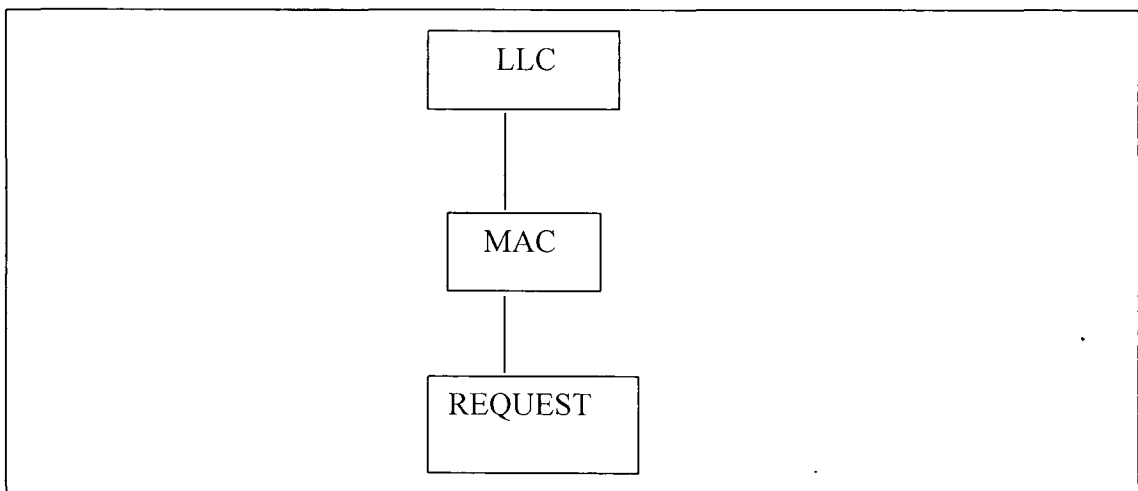
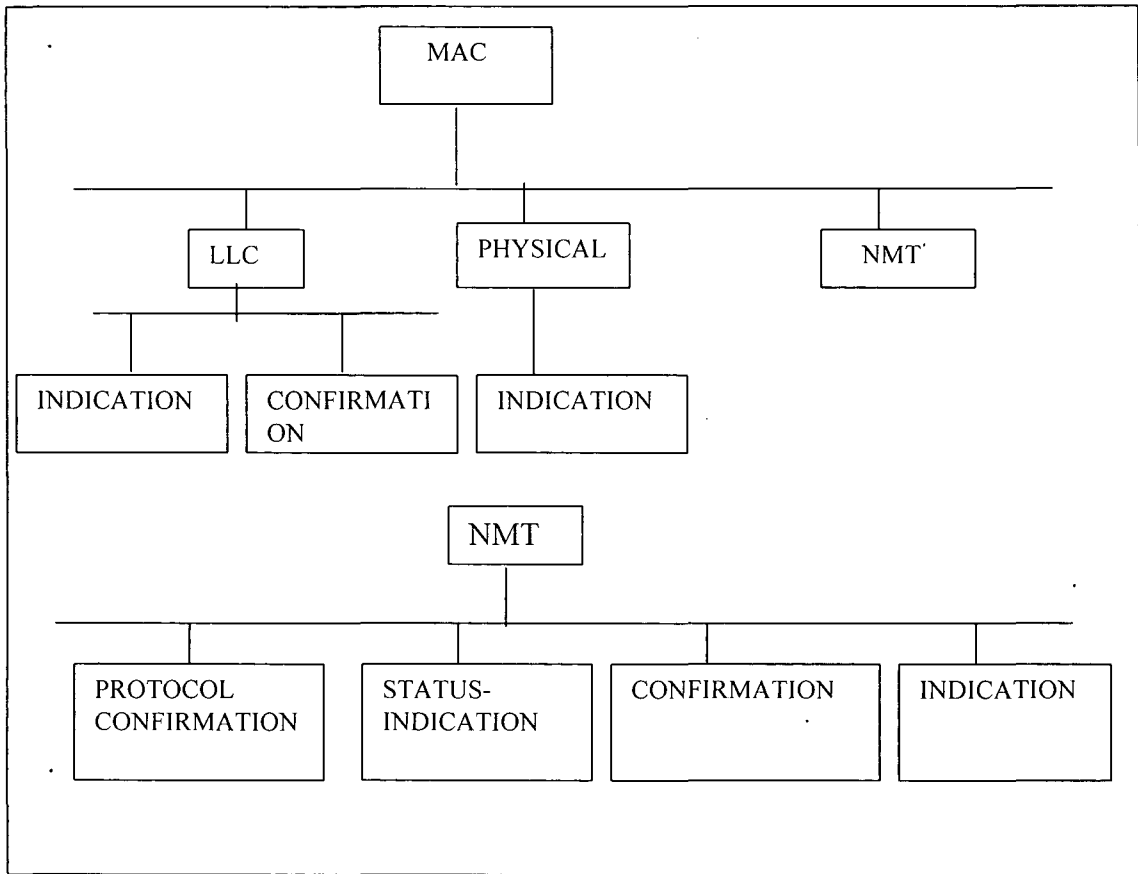
TRUNK-CABLE:-The function of the trunk cable medium is to transport data signals between successive stations of a baseband ring local area network. This communication medium consists of a set of TCU's interconnected sequentially by

trunk cable links. Each TCU is connected to a TCU/MIC cable to which a station may be connected . The relationship between these embodiments and the LAN model is as shown in the figure. Repeaters may be used where required to extend the length of a trunk link beyond limits imposed by normal signal degradation due to link impairments. These repeaters serve to restore the amplitude,shape and timing of signals passing through them. The repeater's regenerative functions have the same characteristics as a repeating station on the ring and must be included in the count of the number of stations supported by the ring.

Ring Access Control:- Station insertion into the ring is controlled by the station. The mechanism for effecting the insertion or bypass of the station resides in the TCU. The station exercises control of the mechanism via the media interface cable .







The services are classified on the basis of operative layer , receiving layer and the type of services . The inheritance tree of service class is shown. A service abstract class is created which has the common attributes and functions of all services . Choosing operative layer as the basis of specialisation. It is derived into four subclasses viz. Physical ,mac,nmtElement and llcSublayer . Specialisation is done using the

specialisation principle of attribute domain restriction . as in case of each of these sublayers domain of the attribute operative layer has been restricted . These service classes are further derived into other services based on the value of the attribute receiving layer . This attribute specifies the layer using that service . In the inheritance tree downward derivation is again done on the basis of service type .

Identifying attributes and functions of the layers :- The mac sublayer is responsible for controlling access to the medium . so it needs certain timer by virtue of which it may control the access. These timers are timer-holding-token ,timer-return-to-repeat, timer-valid-transmission, timer-no-token, timer-queue-PDU, timer-standby-monitor and timer-active-monitor. These attributes have certain fixed values when the timers are reset. The mac-sublayer informs the nmtElement success or failure of the requests made by it so it needs to have an attribute status .The layer has to check frame condition and its report has to be sent to nmt so another attribute required is status-report. This layer needs the functions to check frame condition , reset timers TNT and TSM , enqueue the SMP-PDU, check for presence of active monitor , duplicate address , transmit sdu, append various fields to a frame , hold and regenerate the token on the ring . The physical layer carries out its operation with the aid of trunk-coupling-unit and medium-interface-cable . The trunk-coupling-unit has two state repeat and transmit so it has an attribute state which can have the values repeat and transmit and corresponding functions. The medium-interface-cable consists of physical-mic-cable, tcu-mic-cable and medium interface connector. The nmtElement class has to take actions in error conditions so it has parameters for different error conditions like beacon-state, active-monitor-not-present ,duplication-of-address ,insertion and removal of station from the ring and corresponding functions .

CHAPTER 4

DESIGN IN ASN.1

-- The following is the definition of the aggregation modeling construct

```
AGGREGATION ::= CLASS
{
    &aggregateClass      OBJECT-CLASS
    &componentClass      OBJECT-CLASS
    &componentMultiplicity Multiplicity
    &aggregationLabel    OBJECT IDENTIFIER UNIQUE
}
```

WITH SYNTAX

```
{
    AGGREGATE      &aggreagteClass
    [ WITH MULTIPLICITY &aggreagteMultiplicity ]
    HAS-AS-A-PART
    COMPONENT      &componentClass
    [ WITH MULTIPLICITY &componentMultiplicity ]
    IDENTIFIED BY  &aggreagtionLabel
}
```

```
stationPhysical_layer AGGREGATION ::=
{
    AGGREGATE      station
    HAS-AS-A-PART
    COMPONENT      physical_layer
    IDENTIFIED BY  stationPhysical_layerLbl
}
```

-- Physical Layer is a component of station class

```
stationLlc_layer AGGREGATION ::=
{
    AGGREGATE      station
    HAS-AS-A-PART
    COMPONENT      llc_layer
    IDENTIFIED BY  stationLlc_layerLbl
}
```

-- Llc is a component of station class

```
stationMac_sublayer AGGREGATION ::=
{
    AGGREGATE      station
    HAS-AS-A-PART
```

```

        COMPONENT                mac_sublayer
        IDENTIFIED BY            stationMac_sublayerLbl
    }
-- Mac Layer is a component class of station class

stationNmt_element              AGGREGATION ::=
{
    AGGREGATE                    station
    HAS-AS-A-PART
    COMPONENT                    nmt_element
    IDENTIFIED BY                stationNmt_elementLbl
}
-- Nmt_element is the network manager and controls the overall operation

stationTcu                      AGGREGATION ::=
{
    AGGREGATE                    station
    HAS-AS-A-PART
    COMPONENT                    tcu
    IDENTIFIED BY                stationTcuLbl
}
-- Tcu stands for Transmission control unit

stationMic                      AGGREGATION ::=
{
    AGGREGATE                    station
    HAS-AS-A-PART
    COMPONENT                    mic
    IDENTIFIED BY                stationMicLbl
}
-- Mic stands for medium interface cable

tokenringStation              AGGREGATION ::=
{
    AGGREGATE                    tokenring
    HAS-AS-A-PART
    COMPONENT                    station
    WITH MULTIPLICITY            {
                                    theRangeWith :
                                    { lowerBound { constant : 2};
                                      upperBound {ade : {
                                        numberOfStationsLbl} }
                                    }
    }
    IDENTIFIED BY                tokenringStationLbl
}
-- Token Ring contains at least two stations but the maximum number of stations
-- depends on the attribute numberOfStations.

```

```

micMediumInterfaceConnector      AGGREGATION ::=
{
    AGGREGATE                    mic
    HAS-AS-A-PART
    COMPONENT                    mediumInterfaceConnector
    IDENTIFIED BY                micMediumInterfaceConnectorLbl
}

```

```

micPhy/micCable                 AGGREGATION ::=
{
    AGGREGATE                    mic
    HAS-AS-A-PART
    COMPONENT                    phy/micCable
    IDENTIFIED BY                micPhy/micCableLbl
}

```

```

micTcu/micCable                 AGGREGATION ::=
{
    AGGREGATE                    mic
    HAS-AS-A-PART
    COMPONENT                    tcu/micCable
    IDENTIFIED BY                micTcu/micCableLbl
}

```

-- The following is the definition of the capsule modeling construct

```

CAPSULE ::= CLASS
{
    &Attributes                    ATTRIBUTE OPTIONAL
    &Functions                      FUNCTION OPTIONAL
    &Capsules                       CAPSULE OPTIONAL

    &capsuleLabel                    OBJECT IDENTIFIER UNIQUE
}

```

```

WITH SYNTAX
{
    [ ATTRIBUTES                    &Attributes          ]
    [ FUNCTIONS                      &Functions          ]
    [ CAPSULES                       &Capsules           ]
    IDENTIFIED BY                    &capsuleLabel
}

```

```

start-delimiter                 CAPSULE ::=
{
    ATTRIBUTES
    {
        symbol_J_one,

```



```

symbol_K_one,
symbol_zero_one,
symbol_J_two,
symbol_K_two,
symbol_zero_two,
symbol_zero_three,
symbol_zero_four
    }
IDENTIFIED BY start-delimiterLbl
}

access-control CAPSULE ::=
{
    ATTRIBUTES
        {
            priority,
            token-bit,
            monitor-bit,
            reserved-priority
        }
    IDENTIFIED BY access-controlLbl
}

ending-delimiter CAPSULE ::=
{
    ATTRIBUTES
        {
            symbol-J-one,
            symbol-K-one,
            symbol-one-one,
            symbol-J-two,
            symbol-K-two,
            symbol-one-two,
            symbol-I,
            symbol-E
        }
    IDENTIFIED BY ending-delimiterLbl
}
-- generic sequence is the abstract class from which abort-sequence class is derived

generic-sequence OBJECT-CLASS ::=
{
    MANDATORY CAPSULES
        {
            start-delimiter,
            ending-delimiter,
        }
    FUNCTIONS
        {

```

```

        check-start-delimiter,-- checks starting delimiter
        check-ending-delimiter,
    }
    IDENTIFIED BY generic-sequenceLbl
}

abort-sequence OBJECT-CLASS ::=
{
    SPECIALIZES-FROM
        {{
            superclass generic-sequenceLbl,
            basisOfSpecialization
            {simplePredicate :
                {functionAdded : abort-
                    transmissionLbl }
            }
        }}
    IDENTIFIED BY abort-sequenceLbl }

```

-- abort-sequence is the class used for aborting the transmission

```

generic-token OBJECT-CLASS ::=
{
    SPECIALIZES-FROM
        {{
            superclass generic-sequenceLbl,
            basisOfSpecialization
            { simplePredicate :
                {capsuleAdded :
                    access-controlLbl}
            }
        }}
    FUNCTIONS
        {
            get-priority,
            get-token-bit,
            get-monitor-bit,
            get-reserved-priority,
            set-priority,
            set-token-bit,
            set-monitor-bit,
            set-reserved-priority
        }
    IDENTIFIED BY generic-tokenLbl
}

```

-- generic token is also an abstract class

-- get function returns the value of that attribute

-- set function modifies the value of that attribute

```
token          OBJECT-CLASS
{
    SPECIALIZES-FROM
        {{
            superclass    generic-tokenLbl,
            basisOfSpecialization
            { simplePredicate :
              {attributeDomainRestricted :
                {
                    newValueOf :
                    {attribute    token-bit,
                      is          equalTo,
                      value       0
                    }
                }
            }
        }}—token is derived from the generic-token class on
```

the

```
        }—token-bit attribute with 0 value
    IDENTIFIED BY    tokenLbl
}
```

```
frame          OBJECT-CLASS ::=
{
    SPECIALIZES-FROM
        {{
            superclass    generic-tokenLbl,
            basisOfSpecialization
            {simplePredicate :
              {attributeDomainRestricted :
                { newValueOf:
                  { attribute    token-bit,
                    is          equalTo,
                    value       1
                  }
                }
            }
        }}
    ATTRIBUTES
        {
            destination-address,
            source-address,
            frame-check-sequence,
            info              OPTIONAL
        }
}
```

```

CAPSULES
    {
        frame-control,
        frame-status
    }
IDENTIFIED BY    frameLbl
}
-- frame class is also derived from generic-token class with token-bit =1

frameLbl ::=frame.&objectClassLabel

```

```

frame-control    CAPSULE ::=
{
    ATTRIBUTES
        {
            frame-type-bits,
            control-type-bits,
        }
    IDENTIFIED BY    frame-controlLbl
}

```

```

macframe        OBJECT-CLASS ::=
{
    SPECIALIZES-FROM
        {{
            superclass    frameLbl,
            basisOfSpecialization
            {simplePredicate :
            {attributeDomainRestricted:
            {newValueOf :
            {attribute    frame-type-bits,
            is            equalTo,
            value        0
            }
            }}}
        }}
}
-- macframe is derived from the frame class with frame-type-bit =0

```

```

llcframe        OBJECT-CLASS ::=
{
    SPECIALIZES-FROM
        {{
            superclass    frameLbl,
            basisOfSpecialization
            {simplePredicate :
            {attributeDomainRestricted:
            {newValueOf :

```

```

        {attribute      frame-type-bits,
          is            equalTo,
          value        1
        }
      }}}}
    }}
  }
-- llcframe is derived from the frame class with frame-type-bit = 1

```

```

OBJECT-CLASS ::= CLASS
{
  &SpecializesFrom Specialization OPTIONAL,
  &Attributes       ATTRIBUTE OPTIONAL,
  &Functions        FUNCTION  OPTIONAL,
  &MandatoryCapsules CAPSULE  OPTIONAL,
  &OptionalCapsules CAPSULE  OPTIONAL,
  &ObjectClassLabel OBJECT IDENTIFIER UNIQUE
}

```

WITH SYNTAX

```

{
  [SPECIALIZES-FROM      &Specialization      ]
  [ATTRIBUTES           &Attributes           ]
  [FUNCTIONS             &Functions           ]
  [MANDATORY CAPSULES   &MandatoryCapsules   ]
  [OPTIONAL CAPSULES    &OptionalCapsules    ]
  IDENTIFIED BY         &objectClassLabel
}

```

-- The following is the definition of the attribute modelling construct

```

ATTRIBUTE ::= CLASS
{
  &AttributeType,
  &attributeDomain Domain{&AttributeType}OPTIONAL,
  &attributeLabel OBJECT IDENTIFIER UNIQUE
}

```

WITH SYNTAX

```

{
  ATTRIBUTE-TYPE      &AttributeType
  [ATTRIBUTE-DOMAIN   &attributeDomain  ]
  IDENTIFIED BY      &attributeLabel
}

```

-- The following is the definition of the function modelling construct

```

FUNCTION ::= CLASS
{
  &Arguments ARGUMENT OPTIONAL,
  &Results   RESULT   OPTIONAL,
  &Exceptions EXCEPTION OPTIONAL,
}

```

```

        &Specification                                SET OF SEQUENCE OF
                                                    formalSpecification OPTIONAL,
        &functionLabel                               OBJECT IDENTIFIER UNIQUE
    }
WITH SYNTAX
{
    [ARGUMENTS                                &Arguments ]
    [RESULTS                                  &Results ]
    [EXCEPTIONS                              &Exceptions ]
    [SPECIFICATION                            &Specification ]
    IDENTIFIED BY                             &functionLabel
}

```

-- The following is the definition of the argument modelling construct

```

ARGUMENT ::= CLASS
{
    &ArgumentType,
    &argumentDomain           Domain{&ArgumentType}
    &argumentLabel           OBJECT IDENTIFIER UNIQUE
}

```

WITH SYNTAX

```

{
    ARGUMENT-TYPE           &ArgumentType
    [ARGUMENT-DOMAIN       &argumentDomain ]
    IDENTIFIED BY         argumentLabel
}

```

-- The following is the definition of the result modelling construct

```

RESULT ::= CLASS
{
    &ResultType,
    &ResultDomain           Domain{&ResultType}
    &resultLabel           OBJECT IDENTIFIER UNIQUE
}

```

WITH SYNTAX

```

{
    RESULT-TYPE           &ResultType
    [RESULT-DOMAIN       &ResultDomain ]
    IDENTIFIED BY         &resultLabel
}

```

-- The following is the definition of the exception modelling construct

```

EXCEPTION ::= CLASS
{

```

```

        &ExceptionParamType          OPTIONAL
        &exceptionParamDomain        Domain{&ExceptionParamType}
        &exceptionLabel              OBJECT IDENTIFIER UNIQUE
    }
    WITH SYNTAX
    {
    [EXCEPTION-PARAMETER-TYPE &ExceptionParamType
    [EXCEPTION-PARAMETER-DOMAIN &exceptionParamDomain ]
    IDENTIFIED BY &exceptionLabel.
    }

```

```

Bit ::= INTEGER(0..1)

```

```

symbol-J-one ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE Bit,
    IDENTIFIED BY symbol-J-oneLbl
}

```

```

symbol-J-oneLbl ::= symbol-J-one.&objectClassLabel

```

```

symbol-K-one ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE Bit,
    IDENTIFIED BY symbol-K-oneLbl
}

```

```

symbol-K-oneLbl ::= symbol-K-one.&objectClassLabel

```

```

symbol-zero-one ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE Bit,
    IDENTIFIED BY symbol-zero-oneLbl
}

```

```

symbol-zero-oneLbl ::= symbol-zero-one.&objectClassLabel

```

```

symbol-J-two ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE Bit,
    IDENTIFIED BY symbol-J-twoLbl
}

```

```

symbol-J-twoLbl ::= symbol-J-two.&objectClassLabel

```

```

symbol-K-two ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE Bit,

```

IDENTIFIED BY symbol-K-twoLbl
}

symbol-K-twoLbl ::= symbol-K-two.&objectClassLabel

symbol-zero-two ATTRIBUTE ::=
{
ATTRIBUTE-TYPE Bit,
IDENTIFIED BY symbol-zero-twoLbl
}

symbol-zero-twoLbl ::= symbol-zero-two.&objectClassLabel

symbol-zero-three ATTRIBUTE ::=
{
ATTRIBUTE-TYPE Bit,
IDENTIFIED BY symbol-zero-threeLbl
}

symbol-zero-threeLbl ::= symbol-zero-three.&objectClassLabel

symbol-zero-four ATTRIBUTE ::=
{
ATTRIBUTE-TYPE Bit,
IDENTIFIED BY symbol-zero-fourLbl
}

symbol-zero-fourLbl ::= symbol-zero-four.&objectClassLabel

Priority ::= ENUMERATED{0,1,2,3,4,5,6,7}

priority ATTRIBUTE ::=
{
ATTRIBUTE-TYPE Priority,
IDENTIFIED BY priorityLbl
}

priorityLbl ::=priority.&objectClassLabel

token-bit ATTRIBUTE ::=
{
ATTRIBUTE-TYPE Bit,
IDENTIFIED BY token-bitLbl
}token-bitLbl ::= token-bit.&objectClassLabel

monitor-bit ATTRIBUTE ::=
{
ATTRIBUTE-TYPE Bit,


```

        IDENTIFIED BY      monitor-bitLbl
    }
monitor-bitLbl ::=      monitor-bit.&objectClassLabel

reserved-priority      ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      Priority,
    IDENTIFIED BY      reserved-priorityLbl
}
reserved-priorityLbl ::=reserved-priority.&objectClassLabel

symbol-one-one          ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      symbol-one-oneLbl
}
symbol-one-oneLbl ::= symbol-one-one.&objectClassLabel

symbol-one-two          ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      symbol-one-twoLbl
}
symbol-one-twoLbl ::= symbol-one-two.&objectClassLabel

symbol-I                ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      symbol-ILbl
}
symbol-ILbl ::= symbol-I.&objectClassLabel

symbol-E                ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      symbol-ELbl
}
symbol-ELbl ::= symbol-E.&objectClassLabel

source-address          ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      Address,
    IDENTIFIED BY      source-addressLbl
}
source-addressLbl ::= source-address.&objectClassLabel

destination-address     ATTRIBUTE ::=

```

```

{
    ATTRIBUTE-TYPE      Address,
    IDENTIFIED BY      destination-addressLbl
}
destination-addressLbl ::= destination-address.&objectClassLabel

info
{
    ATTRIBUTE-TYPE      INFO,
    IDENTIFIED BY      infoLbl
}
infoLbl ::= info.&objectClassLabel

INFO ::= SEQUENCE OF OCTET
STRING(SIZE(NUMBER))
NUMBER ::= ENUMERATED {2,6}

Frame-check-sequence
{
    ATTRIBUTE-TYPE      INTEGER32,
    IDENTIFIED BY      frame-check-sequenceLbl
}
frame-check-sequenceLbl ::= frame-check-sequence.&objectClassLabel

frame-type-bits
{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      frame-type-bitsLbl
}
frame-type-bitsLbl ::= frame-type-bits.&objectClassLabel

control-type-bits
{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      control-type-bitsLbl
}
control-type-bitsLbl ::= control-type-bits.&objectClassLabel

check-start-delimiter
{
    ARGUMENTS          start-delimiter
    RESULTS            result
    IDENTIFIED BY      check-start-delimiterLbl
}
check-start-delimiterLbl ::= check-start-delimiter.&objectClassLabel

start-delimiter
ARGUMENT ::=

```

```

{
    ARGUMENT-TYPE          CAPSULE,
    IDENTIFIED BY          start-delimiterLbl
}
start-delimiterLbl      ::=  start-delimiter.&objectClassLabel

result
{
    RESULT-TYPE           Bit,
    IDENTIFIED BY         resultLbl
}
resultLbl               ::=  result.&objectClassLabel

check-ending-delimiter  FUNCTION ::=
{
    ARGUMENTS             ending-delimiter
    RESULTS               result
    IDENTIFIED BY         check-ending-delimiterLbl
}
check-ending-delimiterLbl ::=  check-ending-delimiter.&objectClassLabel

ending-delimiter
{
    ARGUMENT-TYPE          CAPSULE,
    IDENTIFIED BY          ending-delimiterLbl
}
ending-delimiterLbl    ::=  ending-delimiter.&objectClassLabel

get-priority
{
    RESULTS               priority,
    IDENTIFIED BY         get-priorityLbl
}
get-priorityLbl        ::=  getpriority.&objectClassLabel

get-token-bit
{
    RESULTS               token-bit,
    IDENTIFIED BY         get-token-bitLbl
}
get-token-bitLbl       ::=  get-token-bit.&objectClassLabel

token-bit
{
    RESULT-TYPE           Bit,
    IDENTIFIED BY         token-bitLbl
}
token-bitLbl           ::=  token-bit.&objectClassLabel

```

```

get-monitor-bit                FUNCTION ::=
{
    RESULTS                    monitor-bit,
    IDENTIFIED BY              get-monitor-bitLbl
}
get-monitor-bitLbl             ::=  get-monitor-bit.&objectClassLabel

monitor-bit                    RESULT  ::=
{
    RESULT-TYPE                Bit,
    IDENTIFIED BY              monitor-bitLbl
}
monitor-bitLbl                 ::=  monitor-bit.&objectClassLabel

get-reserved-priority         FUNCTION ::=
{
    RESULTS                    reserved-priority,
    IDENTIFIED BY              get-reserved-priorityLbl
}
get-reserved-priorityLbl      ::=  get-reserved-priority.&objectClassLabel

reserved-priority            RESULT  ::=
{
    RESULT-TYPE                Priority,
    IDENTIFIED BY              reserved-priorityLbl
}
reserved-priorityLbl         ::=  reserved-priority.&objectClassLabel

set-priority                  FUNCTION ::=
{
    ARGUMENTS                  priority,
    RESULTS                    priority,
    IDENTIFIED BY              set-priorityLbl
}
set-priorityLbl              ::=  setpriority.&objectClassLabel

priority                      ARGUMENT ::=
{
    ARGUMENT-TYPE              Priority,
    IDENTIFIED BY              priorityLbl
}
priorityLbl                   ::=  priority.&objectClassLabel

priority                      RESULT  ::=
{
    RESULT-TYPE                Priority,
    IDENTIFIED BY              priorityLbl
}

```

```

}
priorityLbl ::= priority.&objectClassLabel

set-monitor-bit FUNCTION ::=
{
    ARGUMENTS monitor-bit,
    RESULTS monitor-bit,
    IDENTIFIED BY set-monitor-bitLbl
}
set-monitor-bitLbl ::= set-monitor-bit.&objectClassLabel

monitor-bit ARGUMENT ::=
{
    ARGUMENT-TYPE Bit,
    IDENTIFIED BY monitor-bitLbl
}
monitor-bitLbl ::= monitor-bit.&objectClassLabel

set-reserved-priority FUNCTION ::=
{
    ARGUMENT reserved-priority,
    RESULTS reserved-priority,
    IDENTIFIED BY set-reserved-priorityLbl
}
set-reserved-priorityLbl ::= set-reserved-priority.&objectClassLabel

reserved-priority ARGUMENT ::=
{
    ARGUMENT-TYPE Priority,
    IDENTIFIED BY reserved-priorityLbl
}
reserved-priorityLbl ::= reserved-priority.&objectClassLabel

frame-status CAPSULE ::=
{
    ATTRIBUTES {
        first-A-bit,
        first-C-bit,
        first-r-bit,
        second-r-bit,
        second-A-bit,
        second-C-bit,
        third-r-bit,
        fourth-r-bit
    }
    IDENTIFIED BY frane-statusLbl
}

```

frame-statusLbl	::=	frame-status.&objectClassLabel
first-A-bit	ATTRIBUTE ::=	
{		
ATTRIBUTE-TYPE	Bit,	
IDENTIFIED BY	first-A-bitLbl	
}		
first-A-bitLbl	::=	first-A-bit.&objectClassLabel
first-C-bit	ATTRIBUTE ::=	
{		
ATTRIBUTE-TYPE	Bit,	
IDENTIFIED BY	first-C-bitLbl	
}		
first-C-bitLbl	::=	first-C-bit.&objectClassLabel
first-r-bit	ATTRIBUTE ::=	
{		
ATTRIBUTE-TYPE	Bit,	
IDENTIFIED BY	first-r-bitLbl	
}		
first-r-bitLbl	::=	first-r-bit.&objectClassLabel
second-A-bit	ATTRIBUTE ::=	
{		
ATTRIBUTE-TYPE	Bit,	
IDENTIFIED BY	second-A-bitLbl	
}		
second-A-bitLbl	::=	second-A-bit.&objectClassLabel
second-C-bit	ATTRIBUTE ::=	
{		
ATTRIBUTE-TYPE	Bit,	
IDENTIFIED BY	second-C-bitLbl	
}		
second-C-bitLbl	::=	second-C-bit.&objectClassLabel
second-r-bit	ATTRIBUTE ::=	
{		
ATTRIBUTE-TYPE	Bit,	
IDENTIFIED BY	second-r-bitLbl	
}		
second-r-bitLbl	::=	second-r-bit.&objectClassLabel
third-r-bit	ATTRIBUTE ::=	

```

{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      third-r-bitLbl
}
third-r-bitLbl        ::=      third-r-bit.&objectClassLabel

forth-r-bit
{
    ATTRIBUTE-TYPE      Bit,
    IDENTIFIED BY      forth-r-bitLbl
}
forth-r-bitLbl        ::=      forth-r-bit.&objectClassLabel

Specialization        ::=      SEQUENCE
{
    superclass          ObjectClassLabel,
    basisOfSpecialization LogicalPredicate
}

abort-transmission    FUNCTION ::=
{
    RESULTS             status,
    IDENTIFIED BY      abort-tarnsmissionLbl
}
abort-transmissionLbl ::=      abort-transmission.&objectClassLabel

status
{
    RESULT-TYPE         Bit,
    IDENTIFIED BY      statusLbl
}
statusLbl             ::=      status.&objectClassLabel

service
{
    ATTRIBUTES
    {
        serviceName,
        serviceType,
        serviceEntity,
        operativeLayer,
        receivingLayer
    }
    IDENTIFIED BY      serviceLbl
}
serviceLbl            ::=      service.&objectClassLabel

physical-LayerService OBJECT-CLASS ::=

```

```

{
    SPECIALIZES-FROM {{
        Superclass    serviceLbl,
        BasisOfSpecialization
        {simplePredicate:
        {attributeDomainRestricted:
        {newValueOf:
        {attribute    operativeLayerLbl,
          is          equalTo,
          value       physical-Layer,
        }}}}}
    }}

ph-data-indication
{
    SPECIALIZES-FROM
        {{
        superclass    physiacal-LayerLbl,
        basisOfSpecialization {compoundPredicate:
            {connectBy and,
              predicates
            {simplePredicate:
            {attributeDomainRestricted:
            {newValueOf:
            {attribute    serviceType,
              is          equalTo,
              value       indication
            }}}}}
            {simplePredicate:
            {attributeDomainRestricted:
            {newValueOf:
            {attribute    receivingLayer,
              is          equalTo,
              value       macSublayer
            }}}}}
        }}
    IDENTIFIED BY
}
ph-data-indicationLbl ::= ph-data-indication.&objectClassLabel

ph-data-confirmation
{
    SPECIALIZES-FROM
        {{
        superclass    physiacal-LayerLbl,
        basisOfSpecialization {compoundPredicate:
            {connectBy and,
              predicates

```



```

    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          serviceType,
    is                  equalTo,
    value               confirmation
    }}}}}
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          receivingLayer,
    is                  equalTo,
    value               macSublayer
    }}}}}
    }}
    IDENTIFIED BY ph-data-confirmationLbl
  }
  ph-data-confirmationLbl ::= ph-data-confirmation.&objectClassLabel

  ph-status-indication OBJECT-CLASS ::=
  {
    SPECIALIZES-FROM
    superclass
    {{
    macSublayerLbl,
    basisOfSpecialization {compoundPredicate:
    {connectBy and,
    predicates
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          serviceType,
    is                  equalTo,
    value               status-indication
    }}}}}
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          receivingLayer,
    is                  equalTo,
    value               nmtElement
    }}}}}
    }}
    IDENTIFIED BY ph- status-indication Lbl
  }
  ph- status-indication Lbl ::= ph- status-indication .&objectClassLabel

  ma-status-indication OBJECT-CLASS ::=
  {

```

```

SPECIALIZES-FROM      {{
                        superclass    macSublayerLbl,
                        basisOfSpecialization {compoundPredicate:
                            {connectBy and,
                             predicates
                              {simplePredicate:
                               {attributeDomainRestricted:
                                {newValueOf:
                                 {attribute      serviceType,
                                  is            equalTo,
                                   value        status-indication
                                }}}}}
                              {simplePredicate:
                               {attributeDomainRestricted:
                                {newValueOf:
                                 {attribute      receivingLayer,
                                  is            equalTo,
                                   value        nmtElement
                                }}}}}
                              }}
                        IDENTIFIED BY ma- status-indication Lbl
}
ma- status-indication Lbl ::= ma- status-indication .&objectClassLabel

macSublayerService    OBJECT-CLASS ::=
{
    SPECIALIZES-FROM {{
                        Superclass    serviceLbl,
                        BasisOfSpecialization
                        {simplePredicate:
                         {attributeDomainRestricted:
                          {newValueOf:
                           {attribute    operativeLayerLbl,
                            is          equalTo,
                             value      macSublayer,
                           }}}}}
                        }}
}

ph-data-request       OBJECT-CLASS ::=
{
    SPECIALIZES-FROM      {{
                        superclass    macSublayerLbl,
                        basisOfSpecialization {compoundPredicate:
                            {connectBy and,
                             predicates
                              {simplePredicate:

```

```

{attributeDomainRestricted:
{newValueOf:
{attribute          serviceType,
is                  equalTo,
value              request
}}}}
{simplePredicate:
{attributeDomainRestricted:
{newValueOf:
{attribute          receivingLayer,
is                  equalTo,
value              physical-Layer
}}}}
}}}}

ma-nmt-data-indication OBJECT-CLASS ::=
{
    SPECIALIZES-FROM      {{
superclass      macSublayerLbl,
basisOfSpecialization {compoundPredicate:
    {connectBy and,
predicates
{simplePredicate:
{attributeDomainRestricted:
{newValueOf:
{attribute          serviceType,
is                  equalTo,
value              indication
}}}}
{simplePredicate:
{attributeDomainRestricted:
{newValueOf:
{attribute          receivingLayer,
is                  equalTo,
value              nmtElement
}}}}
}}}}

ma-data-indication OBJECT-CLASS ::=
{
    SPECIALIZES-FROM      {{
superclass      macSublayerLbl,
basisOfSpecialization {compoundPredicate:
    {connectBy and,
predicates
{simplePredicate:
{attributeDomainRestricted:
{newValueOf:
{attribute          serviceType,
is                  equalTo,

```

```

value          indication
}}}}
{simplePredicate:
{attributeDomainRestricted:
{newValueOf:
{attribute      receivingLayer,
  is            equalTo,
  value        llcSublayer
}}}}
ma-nmt-data-confirmation OBJECT-CLASS ::=
{
  SPECIALIZES-FROM      {{
                        superclass    macSublayerLbl,
                        basisOfSpecialization {compoundPredicate:
                          {connectBy and,
                        predicates
                        {simplePredicate:
                        {attributeDomainRestricted:
                        {newValueOf:
                        {attribute      serviceType,
                          is            equalTo,
                          value        confirmation
                        }}}}}
                        {simplePredicate:
                        {attributeDomainRestricted:
                        {newValueOf:
                        {attribute      receivingLayer,
                          is            equalTo,
                          value        nmtElement
                        }}}}}
  }}
  IDENTIFIED BY        ma-nmt-data-confirmationLbl
}
ma-nmt-data-confirmationLbl ::= ma-nmt-data-confirmation.&objectClassLabel

ma-data-confirmation OBJECT-CLASS ::=
{
  SPECIALIZES-FROM      {{
                        superclass    macSublayerLbl,
                        basisOfSpecialization {compoundPredicate:
                          {connectBy and,
                        predicates
                        {simplePredicate:
                        {attributeDomainRestricted:
                        {newValueOf:
                        {attribute      serviceType,

```

```

        is          equalTo,
        value       confirmation
    }}}}
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute      receivingLayer,
    is             equalTo,
    value          llcSublayer
    }}}}
}}
    IDENTIFIED BY      ma-data-confirmationLbl
}
ma-data-confirmationLbl ::= ma-data-confirmation.&objectClassLabel

ma-initialize-protocol-confirmation OBJECT-CLASS ::=
{
    SPECIALIZES-FROM    {{
        superclass      macSublayerLbl,
        basisOfSpecialization {compoundPredicate:
            {connectBy and,
            predicates
            {simplePredicate:
            {attributeDomainRestricted:
            {newValueOf:
            {attribute      serviceType,
            is             equalTo,
            value          procol-confirmation
            }}}}
            {simplePredicate:
            {attributeDomainRestricted:
            {newValueOf:
            {attribute      receivingLayer,
            is             equalTo,
            value          nmtElement
            }}}}
    }}
}
    IDENTIFIED BY      ma-initialize-protocol-confirmationLbl
}
ma-initialize-protocol-confirmationLbl ::= ma-initialize-protocol-confirmation.&objectClassLabel

ma-data-request OBJECT-CLASS ::=
{
    SPECIALIZES-FROM    {{
        superclass      llcSublayerLbl,
        basisOfSpecialization {compoundPredicate:
            {connectBy and,

```

```

    predicates
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          serviceType,
      is                equalTo,
      value             request
    }}}}}
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          receivingLayer,
      is                equalTo,
      value             macSublayer
    }}}}}

ma-nmt-data-request OBJECT-CLASS ::=
{
    SPECIALIZES-FROM
    {
    superclass      nmtElementLbl,
    basisOfSpecialization {compoundPredicate:
        {connectBy and,
    predicates
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          serviceType,
      is                equalTo,
      value             request
    }}}}}
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute          receivingLayer,
      is                equalTo,
      value             macSublayer
    }}}}}

ma-initialize-protocol-request OBJECT-CLASS ::=
{
    SPECIALIZES-FROM
    {
    superclass      nmtElementLbl,
    basisOfSpecialization {compoundPredicate:
        {connectBy and,
    predicates
    {simplePredicate:
    {attributeDomainRestricted:

```

```

    {newValueOf:
    {attribute      serviceType,
      is           equalTo,
      value        protocol-request
    }}}
    {simplePredicate:
    {attributeDomainRestricted:
    {newValueOf:
    {attribute      receivingLayer,
      is           equalTo,
      value        macSublayer
    }}}
  }
}

ma-control-request      OBJECT-CLASS ::=
{
    SPECIALIZES-FROM    {{
                        superclass      nmtElementLbl,
                        basisOfSpecialization {compoundPredicate:
                        {connectBy and,
                        predicates
                        {simplePredicate:
                        {attributeDomainRestricted:
                        {newValueOf:
                        {attribute      serviceType,
                          is           equalTo,
                          value        contol-request
                        }}}
                        {simplePredicate:
                        {attributeDomainRestricted:
                        {newValueOf:
                        {attribute      receivingLayer,
                          is           equalTo,
                          value        macSublayer
                        }}}
                    }
}

PROTOCOL-LAYERING      ::= CLASS
{
    &upperProtocol      ProtocolLabel,
    &lowerProtocol      ProtocolLabel,
    &demultiplexPoint   INTEGER OPTIONAL,
    &layeringLabel      OBJECT IDENTIFIER UNIQUE
}

WITH SYNTAX
{
    PROTOCOL            &upperProtocol
}

```

<pre> LAYERS-ABOVE PROTOCOL [DEMULTIPLEX POINT IDENTIFIED BY } </pre>	<pre> &lowerProtocol &demultiplexPoint] &layerLabel </pre>
<pre> llcSublayer-macSublayer { PROTOCOL LAYERS-ABOVE PROTOCOL IDENTIFIED BY } llcSublayer-macSublayerLbl ::= macSublayer.&objectClassLabel </pre>	<pre> PROTOCOL-LAYERING : ::= llcSublayer macSublayer llcSublayer-macSublayerLbl llcSublayer- </pre>
<pre> macSublayer-physicalLayer { PROTOCOL LAYERS-ABOVE PROTOCOL IDENTIFIED BY } macSublayer-physicalLayerLbl ::= physicalLayer.&objectClassLabel </pre>	<pre> PROTOCOL-LAYERING ::= macSublayer physicalLayer macSublayer-physicalLayerLbl macSublayer- </pre>
<pre> nmtElement-physicalLayer { PROTOCOL LAYERS-ABOVE PROTOCOL IDENTIFIED BY } nmtElement-physicalLayerLbl ::= physicalLayer.&objectClassLabel </pre>	<pre> PROTOCOL-LAYERING ::= nmtElement physicalLayer nmtElement-physicalLayerLbl nmtElement- </pre>
<pre> nmtElement-macSublayer { PROTOCOL LAYERS-ABOVE PROTOCOL IDENTIFIED BY } nmtElement-macSublayerLbl ::= </pre>	<pre> PROTOCOL-LAYERING : ::= nmtElement macSublayer nmtElement-macSublayerLbl nmtElement-macSublayer.&objectClassLabel </pre>


```

serviceName          ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE    PrintableString
    IDENTIFIED BY    serviceNameLbl
}
serviceNameLbl      ::=  serviceName.&objectClassLabel

PrintableString      ::=  SEQUENCE OF OCTET STRING

ServiceType          ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE    PrintableString,
    IDENTIFIED BY    serviceTypeLbl
}
serviceTypeLbl      ::=  serviceType.&objectClassLabel

operativeLayer       ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE    Layer
    IDENTIFIED BY    operativeLayerLbl
}
operativeLayerLbl   ::=  operativeLayer.&objectClassLabel

LAYER                ::=  ENUMERATED
{
    physicalLayer(0),
    macSublayer(1),
    llcSublayer(2),
    nmtElement(3)
}

receivingLayer       ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE    Layer
    IDENTIFIED BY    receivingLayerLbl
}
receivingLayerLbl   ::=  receivingLayer.&objectClassLabel

macSublayer          OBJECT-CLASS ::=
{
    ATTRIBUTES
    {
        individual-mac-address,
        group-mac-address,
        tht-value,
        trr-value,
        tvx-value,
        tnt-value,
    }
}

```

<pre> } FUNCTIONS { </pre>	<pre> tqp-value, tsm-value, tam-value, priority-Amp-Data-Unit, status, m-sdu, status-report, frame-control, reception-status, transmission-status, provided-service-class, symbol </pre>
<pre> } IDENTIFIED BY } macSublayerLbl ::= nmtElement { ATTRIBUTES { </pre>	<pre> report-transmission-status, receive-frame-control, receive-m-sdu-identification, receive-requested-service-class, master-reset, insert, check-frame-condition, check-for-active-monitor, set-parameters, reset-TNT, reset-TSM, enqueue-SMP-PDU, set-flag, append-DA-to-msdu, transmit-msdu-macSublayer, macSublayerLbl macSublayer.&objectClassLabel OBJECT-CLASS ::= control-action, frame-control, destination-address, m-sdu, requested-service-class, status, </pre>
<pre> } FUNCTIONS { </pre>	

		reporting-frame-condition, report-transmission-status, report-activemonitor-not-present, report-beacon-state, report-provided-service-class, report-frame-control, report-destination-address, report-source-address, report-m-sdu, report-reception-status,
	}	
	IDENTIFIED BY	nmtElementLbl
}		
nmtElementLbl	::=	nmtElement.&objectClassLabel
status-report		ATTRIBUTE ::=
{		
	ATTRIBUTE-TYPE	PrintableString,
	IDENTIFIED BY	status-reportLbl
}		
status-reportLbl	::=	status-report.&objectClassLabel
transmission-status		ATTRIBUTE ::=
{		
	ATTRIBUTE-TYPE	Bit,
	IDENTIFIED BY	transmission-statusLbl
}		
transmission-statusLbl	::=	transmission-status.&objectClassLabel
get-symbol		FUNCTION ::=
{		
	RESULTS	symbol,
	IDENTIFIED BY	get-symbolLbl
}		
get-symbolLbl	::=	get-symbol.&objectClassLabel
symbol		RESULT ::=
{		
	RESULT-TYPE	Bit,
	IDENTIFIED BY	symbolLbl
}		
symbolLbl	::=	symbol.&objectClassLabel
symbol-decoded		FUNCTION ::=
{		
	ARGUMENTS	stream,
	RESULTS	symbol,

<pre> IDENTIFIED BY } symbol-decodedLbl ::= symbol-decoded.&objectClassLabel </pre>	<pre> symbol-decodedLbl </pre>
<pre> stream { ARGUMENT-TYPE IDENTIFIED BY } streamLbl ::= stream.&objectClassLabel </pre>	<pre> ARGUMENT ::= Stream, streamLbl </pre>
<pre> burst-correction-start { RESULTS IDENTIFIED BY } burst-correction-startLbl ::= burst-correction-start.&objectClassLabel </pre>	<pre> FUNCTION ::= symbol burst-correction-startLbl </pre>
<pre> burst-correction-end { ARGUMENTS IDENTIFIED BY } burst-correction-endLbl ::= burst-correction-end.&objectClassLabel </pre>	<pre> FUNCTION ::= transition, burst-correction-endLbl </pre>
<pre> latency-buffer-overflow { RESULTS IDENTIFIED BY } latency-buffer-overflowLbl ::= latency-buffer-overflow.&objectClassLabel </pre>	<pre> FUNCTION ::= status-report, latency-buffer-overflowLbl </pre>
<pre> status-report { RESULT-TYPE IDENTIFIED BY } </pre>	<pre> RESULT ::= PrintableString, status-report.&objectClassLabel </pre>
<pre> latency-buffer-underflow { RESULTS IDENTIFIED BY } latency-buffer-underflowLbl ::= latency-buffer-underflow.&objectClassLabel </pre>	<pre> FUNCTION ::= status-report, latency-buffer-underflowLbl </pre>
<pre> tht-value </pre>	<pre> ATTRIBUTE ::= </pre>

```

{
    ATTRIBUTE-TYPE      TimeStamp,
    IDENTIFIED BY      tht-valueLbl
}
tht-valueLbl          ::= tht-value.&objectClassLabel

trr-value            ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      TimeStamp,
    IDENTIFIED BY      trr-valueLbl
}
trr-valueLbl         ::= trr-value.&objectClassLabel

tvx-value            ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      TimeStamp,
    IDENTIFIED BY      tvx-valueLbl
}
tvx-valueLbl         ::= tvx-value.&objectClassLabel

tnt-value            ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      TimeStamp,
    IDENTIFIED BY      tnt-valueLbl
}
tnt-valueLbl         ::= tnt-value.&objectClassLabel

tqp-value            ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      TimeStamp,
    IDENTIFIED BY      tqp-valueLbl
}
tqp-valueLbl         ::= tqp-value.&objectClassLabel

tsm-value            ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      TimeStamp,
    IDENTIFIED BY      tsm-valueLbl
}
tsm-valueLbl         ::= tsm-value.&objectClassLabel

tam-value            ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE      TimeStamp,
    IDENTIFIED BY      tam-valueLbl
}
tam-valueLbl         ::= tam-value.&objectClassLabel

```

individual-mac-address	ATTRIBUTE ::=
{	
ATTRIBUTE-TYPE	Address,
IDENTIFIED BY	individual-mac-addressLbl
}	
individual-mac-addressLbl	::= individual-mac-address.&objectClassLabel
group-mac-address	ATTRIBUTE ::=
{	
ATTRIBUTE-TYPE	AddressSet,
IDENTIFIED BY	group-mac-addressLbl
}	
group-mac-addressLbl	::= group-mac-address.&objectClassLabel
AddressSet	::= SET OF Address
priority-Amp-Data-Unit	ATTRIBUTE ::=
{	
ATTRIBUTE-TYPE	Priority
IDENTIFIED BY	priority-Amp-Data-UnitLbl
}	
priority-Amp-Data-UnitLbl	::= priority-Amp-Data-Unit.&objectClassLabel
status	ATTRIBUTE ::=
{	
ATTRIBUTE-TYPE	Bit,
IDENTIFIED BY	statusLbl
}	
statusLbl	::= status.&objectClassLabel
reception-status	ATTRIBUTE ::=
{	
ATTRIBUTE-TYPE	Bit,
IDENTIFIED BY	reception-statusLbl
}	
reception-statusLbl	::= reception-status.&objectClassLabel
m-sdu	ATTRIBUTE ::=
{	
ATTRIBUTE-TYPE	INFO,
IDENTIFIED BY	m-sduLbl
}	
m-sduLbl	::= m-sdu.&objectClassLabel
transmission-status	ATTRIBUTE ::=
{	
ATTRIBUTE-TYPE	Bit,
IDENTIFIED BY	transmission-statusLbl

```

}
transmission-statusLbl ::= transmission-status.&objectClassLabel

provided-service-class ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE Priority,
    IDENTIFIED BY provide-service-classLbl
}
provided-service-classLbl ::= provided-service-class.&objectClassLabel

check-frame-conditio FUNCTION ::=
{
    ARGUMENT input-frame,
    RESULTS status-report,
    IDENTIFIED BY check-frame-conditionLbl
}
check-frame-conditionLbl ::= check-frame-condition.&objectClassLabel

input-frame ARGUMENT ::=
{
    ARGUMENT-TYPE Frame,
    IDENTIFIED BY input-frameLbl
}
input-frameLbl ::= input-frame.&objectClassLabel

check-for-active-monitor FUNCTION ::=
{
    RESULTS value,
    IDENTIFIED BY check-for-active-monitorLbl
}
check-for-active-monitorLbl ::= check-for-active-monitor.&objectClassLabel

value RESULT ::=
{
    RESULT-TYPE Bit,
    IDENTIFIED BY valueLbl,
}
valueLbl ::= value.&objectClassLabel

set-parameters FUNCTION ::=
{
    RESULTS result-status,
    IDENTIFIED BY set-parameterLbl
}
set-parameterLbl ::= set-parameter.&objectClassLabel

result-status RESULT ::=
{

```

<pre> RESULT-TYPE IDENTIFIED BY } result-statusLbl </pre>	<pre> Bit, result-statusLbl ::= result-status.&objectClassLabel </pre>
<pre> reset-TNT { ARGUMENTS RESULTS IDENTIFIED BY } reset-TNTLbl </pre>	<pre> FUNCTION ::= tnt-value, tnt-value, reset-TNTLbl reset-TNT.&objectClassLabel </pre>
<pre> tnt-value { ARGUMENT-TYPE IDENTIFIED BY } tnt-valueLbl </pre>	<pre> ARGUMENT ::= TimeStamp, tnt-valueLbl tnt-value.&objectClassLabel </pre>
<pre> tnt-value { RESULT-TYPE IDENTIFIED BY } tnt-valueLbl </pre>	<pre> RESULT ::= TimeStamp, tnt-valueLbl tnt-value.&objectClassLabel </pre>
<pre> reset-TSM { ARGUMENTS RESULTS IDENTIFIED BY } reset-TSMLbl </pre>	<pre> FUNCTION ::= tsm-value, tsm-value, reset-TSMLbl reset-TSM.&objectClassLabel </pre>
<pre> tsm-value { ARGUMENT-TYPE IDENTIFIED BY } tsm-valueLbl </pre>	<pre> ARGUMENT ::= TimeStamp, tsm-valueLbl tsm-value.&objectClassLabel </pre>
<pre> tsm-value { RESULT-TYPE IDENTIFIED BY } </pre>	<pre> RESULT ::= TimeStamp, tsm-valueLbl </pre>

tsm-valueLbl	::=	tsm-value.&objectClassLabel
enqueue-SMP-PDU	FUNCTION ::=	
{		
ARGUMENTS		pdu,
RESULTS		position-in-queue,
IDENTIFIED BY		enqueue-SMP-PDULbl
}		
enqueue-SMP-PDULbl	::=	enqueue-SMP-PDU.&objectClassLabel
pdu	ARGUMENT ::=	
{		
ARGUMENT-TYPE		INFO,
IDENTIFIED BY		pduLbl,
}		
pduLbl	::=	pdu.&objectClassLabel
position-in-queue	RESULT ::=	
{		
RESULT-TYPE		INTEGER,
IDENTIFIED BY		position-in-queueLbl,
}		
position-in-queueLbl	::=	position-in-queue.&objectClassLabel
set-flag	FUNCTION ::=	
{		
ARGUMENTS		flag,
RESULTS		flag,
IDENTIFIED BY		set-flagLbl
}		
set-flagLbl	::=	set-flag.&objectClassLabel
flag	ARGUMENT ::=	
{		
ARGUMENT-TYPE		Bit,
IDENTIFIED BY		flagLbl
}		
flagLbl	::=	flag.&objectClassLabel
flag	RESULT ::=	
{		
RESULT-TYPE		Bit,
IDENTIFIED BY		flagLbl
}		
flagLbl	::=	flag.&objectClassLabel
append-source-address-to-m-sdu	FUNCTION ::=	
{		

```

        ARGUMENTS          m-sdu,
        RESULTS            m-sdu,
        IDENTIFIED BY     append-source-address-to-m-sduLbl
    }
append-source-address-to-m-sduLbl ::=append-source-address-to-m-
sdu.&objectClassLabel

m-sdu          ARGUMENT ::=
{
    ARGUMENT-TYPE      INFO,
    IDENTIFIED BY     m-sduLbl
}
m-sduLbl      ::= m-sdu.&objectClassLabel

m-sdu          RESULT ::=
{
    RESULT-TYPE        INFO,
    IDENTIFIED BY     m-sduLbl
}

append-DA-to-m-sdu      FUNCTION ::=
{
    ARGUMENTS          m-sdu,
    RESULTS            m-sdu,
    IDENTIFIED BY     append-DA-to-m-sduLbl
}
append-DA-to-m-sduLbl  ::= append-DA-to-m-sdu.&objectClassLabel

transmit-sdu-macSublayer      FUNCTION ::=
{
    ARGUMENTS          frame,
    RESULTS            transmission-status,
    IDENTIFIED BY     transmit-sdu-macSublayerLbl
}
transmit-sdu-macSublayerLbl::= transmit-sdu-macSublayer.&objectClassLabel

transmission-status          RESULT ::=
{
    RESULT-TYPE          Bit,
    IDENTIFIED BY       transmission-statusLbl
}
transmission-statusLbl      ::= transmission-status.&objectClassLabel

report-transmission-status    FUNCTION ::=
{
    RESULTS            transmission-status
    IDENTIFIED BY     report-transmission-statusLbl
}

```

report-transmission-statusLbl ::=	report-transmission-status.&objecClassLabel
receive-frame-control { ARGUMENTS IDENTIFIED BY }	FUNCTION ::= frame-control, receive-frame-controlLbl
receive-frame-controlLbl ::=	receive-frame-control.&objecClassLabel
frame-control { ARGUMENT-TYPE IDENTIFIED BY }	ARGUMENT ::= Frame-control, frame-controlLbl
frame-controlLbl ::=	frame-control.&objectClassLabel
receive-destination-address { ARGUMENTS IDENTIFIED BY }	FUNCTION ::= destination-address, receive-destination-addressLbl
receive-destination-addressLbl address.&objectClassLabel	::= receive-destination-
destination-address { ARGUMENT-TYPE IDENTIFIED BY }	ARGUMENT ::= Address, destination-addressLbl
destination-addressLbl	::= destination-address.&objectClassLabel
receive-m-sdu-identification { ARGUMENTS IDENTIFIED BY }	FUNCTION ::= m-sdu, receive-m-sdu-identificationLbl
receive-m-sdu-identificationLbl identification.&objectClassLabel	::= receive-m-sdu-
m-sdu { ARGUMENT-TYPE IDENTIFIED BY }	ARGUMENT ::= INFO, m-sduLbl
receive-requested-service-class { ARGUMENTS	FUNCTION ::= requested-service-class,

<pre> IDENTIFIED BY } receive-requested-service-classLbl ::= receive-requested-service- class.&objectClassLabel </pre>	<pre> receive-requested-service-classLbl </pre>
<pre> requested-service-class { ARGUMENT-TYPE IDENTIFIED BY } </pre>	<pre> ARGUMENT ::= Priority, requested-service-classLbl </pre>
<pre> master-reset { ARGUMENTS RESULTS IDENTIFIED BY } master-resetLbl ::= master-reset.&objectClassLabel </pre>	<pre> FUNCTION ::= parameters, parameters, master-resetLbl </pre>
<pre> parameters { ARGUMENT-TYPE IDENTIFIED BY } parametersLbl ::= parameters.&objectClassLabel </pre>	<pre> ARGUMENT ::= Parameters, parametersLbl </pre>
<pre> parameters { RESULT-TYPE IDENTIFIED BY } </pre>	<pre> RESULT ::= Parameters, parametersLbl </pre>
<pre> insert { RESULTS IDENTIFIED BY } insertLbl ::= insert.&objectClassLabel </pre>	<pre> FUNCTION ::= status, insertLbl </pre>
<pre> control-action { ATTRIBUTE-TYPE IDENTIFIED BY } control-actionLbl ::= control-action.&objectClassLabel </pre>	<pre> ATTRIBUTE ::= Bit, control-actionLbl </pre>
<pre> destination-address { </pre>	<pre> ATTRIBUTE ::= </pre>

ATTRIBUTE-TYPE IDENTIFIED BY } destination-addressLbl	Address, destination-address Lbl ::= destination-address.&objectClassLabel
reporting-frame-condition { ARGUMENTS RESULTS IDENTIFIED BY } reporting-frame-conditionLbl ::=	FUNCTION ::= input-frame, status-report, reporting-frame-conditionLbl reporting-frame-condition.&objectClassLabel
report-active-monitor-not-present { RESULTS IDENTIFIED BY } report-active-monitor-not-presentLbl ::=	FUNCTION ::= status, report-active-monitor-not-presentLbl report-active-monitor-not-present.&objectClassLabel
report-beacon-state { ARGUMENTS RESULTS IDENTIFIED BY } report-beacon-stateLbl ::=	FUNCTION ::= tnt-value, result-status, report-beacon-stateLbl report-beacon-state.&objectClassLabel
report-provided-service-class { RESULTS IDENTIFIED BY } report-provided-service-classLbl ::=	FUNCTION ::= provided-service-class, report-provided-service-classLbl report-provided-service-class.&objectClassLabel
provided-service-class { RESULT-TYPE IDENTIFIED BY }	RESULT ::= Priority, provided-service-classLbl,
report-frame-control { RESULTS IDENTIFIED BY }	FUNCTION ::= frame-control, report-frame-controlLbl,

report-frame-controlLbl	::=	report-frame-control.&objectClassLabel
frame-control	RESULT ::=	
{		
RESULT-TYPE		Frame-control,
IDENTIFIED BY		frame-controlLbl
}		
report-destination-address	FUNCTION ::=	
{		
ARGUMENTS		frame,
RESULTS		destination-address,
IDENTIFIED BY		report-destination-addresssLbl
}		
report-destination-addressLbl ::=		report-destination-addresss.&objectClassLabel
destination-address	RESULT ::=	
{		
RESULT-TYPE		Address,
IDENTIFIED BY		destination-addressLbl
}		
report-source-address	FUNCTION ::=	
{		
ARGUMENTS		frame,
RESULTS		source-address,
IDENTIFIED BY		report-source-addressLbl
}		
report-source-addressLbl ::=		report-source-address.&objectClassLabel
source-address	RESULT ::=	
{		
RESULT-TYPE		Address,
IDENTIFIED BY		source-addressLbl
}		
report-m-sdu	FUNCTION ::=	
{		
ARGUMENTS		frame,
RESULTS		m-sdu,
IDENTIFIED BY		report-m-sduLbl
}		
report-m-sduLbl ::=		report-m-sdu.&objectClassLabel
m-sdu	RESULT ::=	
{		
RESULT-TYPE		INFO,
IDENTIFIED BY		m-sduLbl
}		

```

report-reception-status          FUNCTION ::=
{
    ARGUMENTS                    input-frame,
    RESULTS                       reception-status,
    IDENTIFIED BY                 report-reception-statusLbl
}
report-reception-statusLbl ::= report-reception-status.&objectClassLabel

reception-status                RESULT ::=
{
    RESULT-TYPE                   Bit,
    IDENTIFIED BY                 reception-statusLbl
}

tcu                              OBJECT-CLASS ::=
{
    ATTRIBUTES                    { state,
                                }
    FUNCTIONS                      {
                                insert,
                                repeat,
                                transmit,
                                loopbacktest,
                                }
    IDENTIFIED BY                 tcuLbl
}
tcuLbl                          ::= tcu.&objectClassLabel

state                            ATTRIBUTE ::=
{
    ATTRIBUTE-TYPE                State,
    IDENTIFIED BY                 stateLbl
}
stateLbl                        ::= state.&objectClasslabel

State                            ::= ENUMERATED
{
    repeat(0),
    transmit(1)
}

insert                           FUNCTION ::=
{
    RESULTS                       output,
    IDENTIFIED BY                 insertLbl
}

```

```

}
insertLbl ::= insert.&objectClassLabel

output RESULT ::=
{
    RESULT-TYPE Bit,
    IDENTIFIED BY outputLbl
}
outputLbl ::= output.&objectClassLabel

remove FUNCTION ::=
{
    RESULTS output,
    IDENTIFIED BY removeLbl
}
removeLbl ::= remove.&objectClassLabel

repeat FUNCTION ::=
{
    ARGUMENTS input-frame,
    RESULTS output,
    IDENTIFIED BY repeatLbl
}

```


5. CONCLUSION

Object-oriented methodology described in chapter 2 is a powerful modelling technique to design communication networks . It is applied for designing Token Ring Network . It can be equally applied for designing other communication networks like Token Bus and FDDI. The work here included the design upto the Medium Access control Layer. It can be further enhanced by including the full design of Logical Link Control Sublayer.

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