

EXTFOR PREPROCESSOR

A Thesis submitted to the Jawaharlal Nehru University
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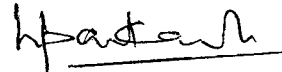
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The work embodied in this dissertation has been carried out at the School of Computer and Systems Sciences, Jawaharlal Nehru University, New Delhi-110 067.

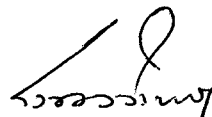
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A C K N O W L E D G E M E N T S

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HARI PARKASH

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INTRODUCTION

A new trend in program writing was introduced by Edsger W. Dijkstra in his monograph Notes on Structured Programming (1972). It takes a 'constructive' approach to the problem of program correctness by controlling the process of program generation to produce a correct program. Initially the term structured programming (sp) was misunderstood by many programmers (e.g. sp is programming without GOTOs, one must use only the conditional and while statements). Although people have different interpretations of sp, it is now widely accepted as an approach to understand the complete programming process so as to carefully control the structure of one's code - its control flow and data organization. The overall efficiency in program development process found to be considerably improved with sp.

Sp methodology tells us not to program in a programming language, but into it. Express the algorithm in a notation that best fits the problem, then refine this algorithm into the programming language. Use any notation that fits the problem and follows clean, understandable approach. If the notation is not with available resources, simulate it in the programming language.

The algorithm expressed in programming language should be easier to read, easier to get right, easier to change, and easier to keep right as they are changed. This is often harder to do with FORTRAN but there is

no harm in simulating the features required to make program writing easy through structured approach.

RATFOR was the first attempt by Dr Kernighan in 1975, which was called a structured language "Rational FORTRAN". Some new features were introduced to FORTRAN and a tool was presented called the preprocessor to translate RATFOR into FORTRAN. The preprocessor proved as a powerful tool to provide with missing features in any programming language. A number of attempts were made by different people on development of an efficient preprocessor and for FORTRAN with better features.

We believe that building on the work of others is the only way to make substantial progress in any field. Therefore we work on the similar guidelines for FORTRAN with the idea of structured programming and use of preprocessor as a tool. Considering FORTRAN as a poor language for structured programming we have extended FORTRAN with better features which we call as EXTFOR (short for Extended FORTRAN). EXTFOR provides modern control flow statements like those in PL/I, ALGOL, or PASCAL, retaining the merits of FORTRAN - a language that is universal and portable while at the same time concealing its worst drawbacks. We are not going to touch the data types of FORTRAN. The new features are easy to read, write, and understand and readily transformable into FORTRAN.

Chapter 1 contains a brief history of development in way of FORTRAN usage. The structured programming is adopted as a trend in programming in FORTRAN.

Chapter 2 introduces the EXTFOR language features with examples and concise summary.

Preprocessing the EXTFOR into FORTRAN is a convenient way to do the job and FORTRAN compiler can take the burden of further processing. Chapter 3 presents the EXTFOR preprocessor.

Finally we discuss the efficiency and overhead in preprocessor, conclusion and experience in Chapters 4 and 5.

In appendices the procedure for interfacing EXTFOR with local OS environment and the EXTFOR operating procedure with examples, error messages list are described.

CONCEPTS AND PROBLEM1.1 HISTORY OF FORTRAN

Hardly anything associated with computer that is as much as twenty five years old is still in use, unless it is the original Von Neumann "Stored Program" concept itself. Yet in spite of vast changes in hardware architecture and in spite of the reproduction of programming languages since 1957, we are still using FORTRAN in a form - not withstanding a major revision with the introduction of FORTRAN-IV, not much different from the original version announced by John W. Backus of IBM in 1957.

FORTRAN is mostly used for scientific and engineering applications because it is easy to learn, availability of many readymade programs written in FORTRAN and very efficient and optimized compilers, attribute to the popularity of the language.

Trends in computers have resulted in some changes in the FORTRAN language and the way it is used. The changes have occurred in two main areas. First, the FORTRAN users have consolidated to a considerable extent around the ANSI standard version of the language, and the rather arbitrary variations (which were wide spread as late as 1966) have declined in number. Most current compilers contain actual conflicts with or violations of the standard and as was intended, the acceptance of the standard has placed a certain onus on variant versions. This has resulted, of course, in increased interchangeability of programs among various users and is no doubt partly responsible for the combined popularity of the language.

1.2 STRUCTURED PROGRAMMING

Reliability is the sine qua non of software systems. There is no point in engaging in efficiency consideration for unreliable systems. And to achieve reliability we should make our code correct and clear before we try to make it faster or more compact. Structured programming is a technical elaboration of this basic point of view.

The term structured programming has been used with many different meanings since Edsger W. Dijkstra first coined the term. The term is generally misunderstood and lack of a precise definition has allowed, even encouraged, people to use it as they wished, for the sake of fashion, to attribute to structured programming (sp) what they themselves learned from monograph notes on structured programming (Dijkstra). Taken out of context or viewed in the wrong light, some of the resulting definition of sp that have appeared seem stupid (e.g. sp is programming without GOTOs), and it is quite understandable that programmers have looked askance when asked to learn and practice it.

Dijkstra proposed that we take a 'constructive' approach to the problem of program correctness by controlling "the process of generation such as to produce a priori correct program". To facilitate this process we limit the rules for program composition to those which are well understood. We would like to informally establish that our program is correct as we are designing and coding it.

The structured programming provides a systematic approach to the programming. A program can be disintegrated into small components so as to make them smaller programs. These component programs can, in turn, be regarded as systems composed of still smaller programs, and so on, down to some reasonable level where the components level to be consisting of a few statements each. The programs written with this approach are structured.

The important thing here is the interface between the components of the program which are systems themselves. Less is the interaction more it attributes to be called as structured. In the concept of structured programming both a design method and a coding technique are implied. The design aspects often embrace the phases "top-down expansion" and stepwise refinement. The coding features include restriction to a set of single entry/single exit control structures, often with explicit formatting conventions to promote readability. And top-down coding is often considered to be a part of structured programming. The top-down programming is step-by-step process of going from a program functional specifications to the coded program.

1.3 STRUCTURED PROGRAMMING - ADVANTAGES

The general view of the structured programming is that it is the general method by which leading programmers program.

It is the process of minimizing the number of interactions between the components of a program.

It helps with converting arbitrarily large and complete

flow chart: into a program represented by inter-relating and nesting a small number of basic and standard control logic structures.

It is a manner of organizing and coding programs that makes the programs easily understood and modified.

The main advantages of structured programming include ease of writing, understanding and debugging. Writing a program becomes easier because of the top-down approach. The program is written as an inter-connection of a few routines with the few interactions between components, they can be designed and tested in relative isolation.

Another important advantage of sp is that the programs so written are easy to understand by a person other than the program writer. Application programs are often modified and run for years by different programmers after the original programmer has gone some where else.

The debugging is the main phase of program development and the time involved in debugging a program is a function of the way (methodology) used in writing the program. Sp facilitates to debug small components in isolation using test data appropriate to the program. When small components have been debugged they can be put together to form the next level of component, which can be debugged in turn and any errors introduced at this level can be attributed with some possibility to the interface between components. The source of the error is narrowed down in this way

and debugging of the large components is simplified. We often use subroutines as component and they are tested independently with variety of test data. The components are interfaced through the parameters.

As far as efficiency is concerned, overall efficiency in a program development is increased considerably.

The current software/hardware ratio of a system is greater than in the early age of computers. Programmers in large software projects typically spend their time as follows:

40-50% is program check out, and one third in program design and less than 20% is actual coding. So the efficiency can be increased if we give more emphasis on design and coding so that the debugging time is minimized. This makes quite clear that sp helps in increasing the overall efficiency.

1.4 STATEMENT OF THE PROBLEM

Some of the main features of a programming language are control structure, data type and input/output. In view of modern trend of programming, FORTRAN has quite weak control structures. In spite of the inadequacy of some features in the language, not many extensions have been made of the FORTRAN and whatsoever were made, have continued to improve the usefulness of the language, both for certain advanced application and for the novice users. Some considerable extensions include WATFOR-WATFIV with its excellent pedagogical features - format free input-output, expressions in output lists, CHARACTER data type, expressions of mixed type, multiple assignments, etc.

As we know that the worst deficiency in using FORTRAN is in the control flow statements - conditional branches and loops which express the logic of the problem. The following examples illustrate:

FORTRAN DO restricts the user to going forward in an arithmetic progression. One cannot use DO with index going from N to M where $M < N$:

```

DO 10 I=10,5,-1
  .
  .
10 CONTINUE

DO 20 R=1.0,3.0,0.5
  .
  .
20 CONTINUE

```

In first case the index is negative arithmetic progression. Second case, the index is real. Both the cases violate the FORTRAN loop index rule. Anyway these deficiencies do not come in the way of developing structured programs using FORTRAN, but the users do feel difficulties with such restrictions.

A statement label is written just after 'DO', which is required to terminate the scope of the DO loop. This label is more serious as far as sp concept is concerned.

The conditional statements in FORTRAN are primitive. The arithmetic IF statement forces the user into at least two statement numbers which in turn give rise to two (implied) GOTOs leading to unintelligible code.

For example:

```

      IF (arithmetic expression) 10,100,200
      .
      .
10   CONTINUE
      .
100  I=1
      .
200  TEMP=ARRY(I)

```

The above IF statement is a three-way branch. Execution of this statement causes evaluation of the expression following which the statement identified by the label 10, 100 or 200 is executed next as the value of expression is less than zero, zero or greater than zero, respectively.

The logical IF is better:

```
IF (logical expression) "statement"
```

Upon the execution of this statement, the logical expression is evaluated; if the value is false statement which following IF statement (next statement) is executed, if value is true the "statement" is executed. Unfortunately this statement is restricted from the use of another "IF statement", DO loop as "statement" in above format. Secondly there can be no ELSE part to the IF statement: There is no way to specify an alternative action if the logical expression is false.

The result of these feelings is that FORTRAN programs must be written with numerous labels and branches. The resulting code is particularly difficult to read and understand, and thus hard to debug and modify.

In view of the above discussion we feel that the FORTRAN language requires some modifications or extensions so as to improve the correctness, readability and efficiency in program writing, in other words to write structured programs. The modification of the existing control structures leads to writing a new compiler, which means too much of a burden for a simple problem.

The approach we choose here is to make FORTRAN a better programming language by extension of the control structures and retaining the merits of the FORTRAN.

The primary purpose of this work is to define a new language which is an extension of the FORTRAN. The extensions are only of control structures; the format of the FORTRAN is not changed and the data structures are not extended. This new language we call as "Extended FORTRAN" (EXTFOR). Further, a system program called preprocessor is written to translate the EXTFOR program into the unpleasant FORTRAN.

EXTFOR - EXTENDED FORTRAN2.1 INTRODUCTION

A program is a static representation of a computation. The execution of a program is a dynamically changing activity. When we use program text to reason about a program execution, we are crucially concerned with the program's flow of control which is determined by the type of constituent control structures. In structured programming we limit the use of three forms of control structures. These forms are simple sequencing or concatenation of statements, selecting the next statement(s) to be executed by the result of a test (often with a true or false answer) and conditional iteration. The restriction to these three means of program compositions is far from arbitrary, rather, it is based directly upon consideration of program correctness, and readability of the program and efficiency.

Each of the forms has the property that it is single entry/single exit. As such they provide alternative ways to structure or to decompose a one-in/one-out box.

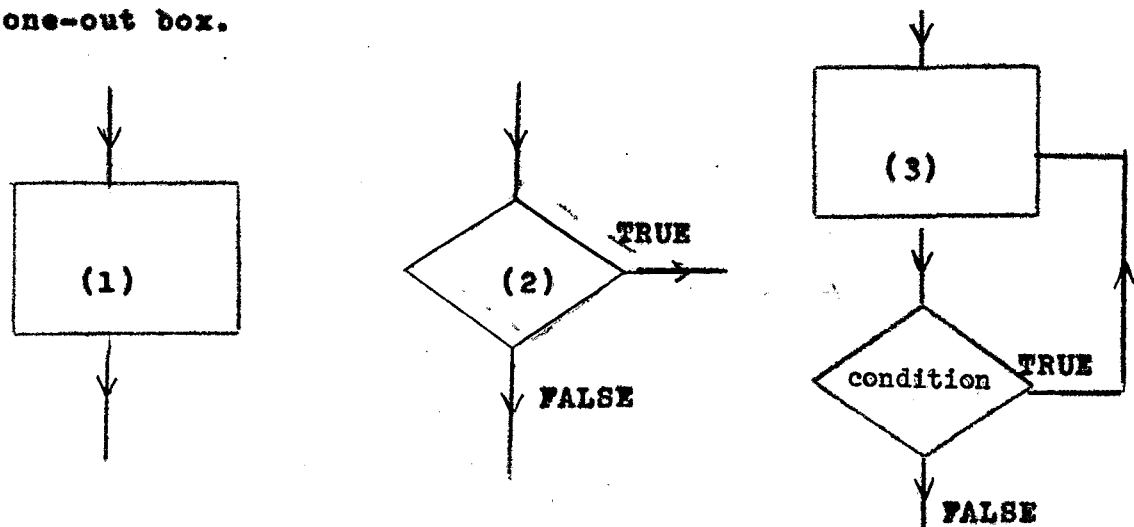


FIG. 2.1

When expanded into code, the forms can be read from top to bottom with no jumping around in the program text, minimizing the interaction among program components.

The control flow structures discussed in this chapter are the extensions to the FORTRAN and selected to provide, keeping in mind the above discussed three forms of control structures for writing structured programs.

2.2 STATEMENT GROUPING

FORTRAN provides no way to group statements together, short of making them into a unit of program. The construction "if a condition is true, do this group of statements", for example:

```
IF (X.GT.100) Y=X+10
              CALL ERROR (X)
              WRITE (6, 10) Y
              X=0
```

cannot be written directly in FORTRAN. The only way to write is:

```
IF (X.LE.100) GO TO 100
Y=X+10
CALL ERROR (X)
WRITE (6, 10) Y
10  FORMAT ("ERROR --", I4)
    X = 0
100 CONTINUE
```

2.2.1 THE BEGIN STATEMENT

In our EXTFOR a group of statements can be treated as a unit by enclosing them between BEGIN and END, of course this END has different meaning than FORTRAN "END" which terminates the scope of a program segment. The format is:

```

BEGIN
  statement-1
  statement-2
  .
  .
  statement-n
END

```

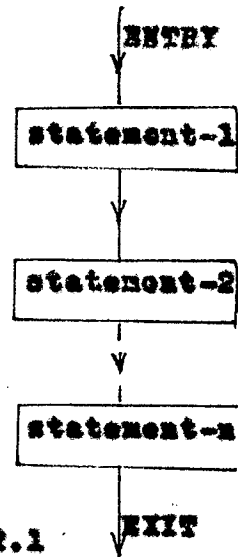


FIG. 2.2.1

BEGIN-END enables us to execute the statements sequentially provided the execution of any of the statements does not transfer control outside the block of statements. This structure is of first type of program components. Using this structure the example in section 2.2 can be written as:

```

IF (X.GT.100) THEN BEGIN
  Y = X + 10
  CALL ERROR (X)
  WRITE (6,10) Y
  X = 0
END

```

This block structure is different from BEGIN-END in PL/I. In PL/I it is often used for new declarations of identifiers or when the meanings of names are to be restricted within a program segment.

2.3 DECISION STRUCTURES

These are structured statements which control the execution of their scope on the basis of a logical expression. These are second type of program components.

2.3.1 THE IF STATEMENT

The format of this construct is:
IF (condition) THEN statement-1
ELSE statement-2

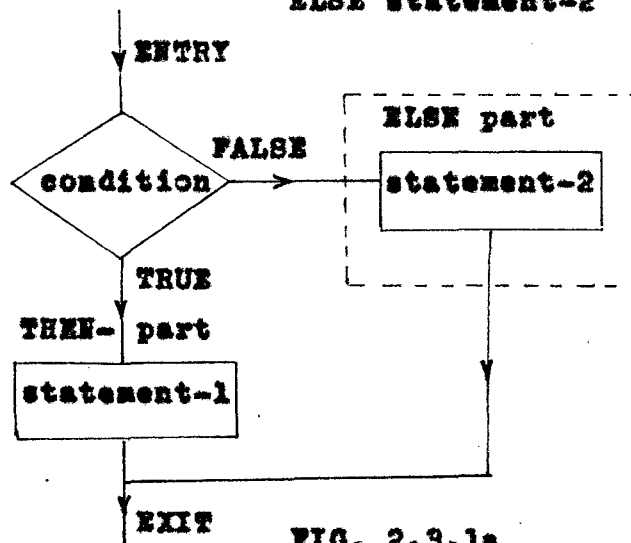


FIG. 2.3.1a

The condition part of the statement is a valid FORTRAN logical expression with value FALSE or TRUE. Statement-1 and statement-2 are any of the EXTFOR statements with some restrictions on using a mixed combination of FORTRAN statements and extended features like FORTRAN DO used as statement in THEN or ELSE unit.

The execution of IF statement first evaluates the condition, if the condition is TRUE then the statement following THEN is executed and ELSE part (if present) is skipped; if the condition results as FALSE the statement following THEN is skipped with execution of the ELSE

part (see FIG. 2.3.1a). So one and only one of the two statements (statement-1 or statement-2) is executed when the IF-THEN-ELSE is encountered. The ELSE part is optional, if not used, the statement is equivalent to FORTRAN logical IF statement. For example:

```
IF (LASTC .GT. MAXL) THEN LASTC = 1
```

is same as

```
IF (LASTC .GT. MAXL) LASTC = 1
```

Both these statements are permitted in EXTPOR.

The recursive use of logical IF statement for example:

```
IF (C.NE.EOF) IF (LINE.LT.60) PAGE=PAGE+1
```

is not allowed in FORTRAN. However, in EXTPOR this nesting of IF statements is permitted. This construct is useful in selecting the next statement execution on subjected to a number of conditions specified. For example:

```
IF (condition-1)
  THEN IF (condition-2)
    THEN statement-1
    ELSE statement-2
  ELSE IF (condition-3)
    THEN statement-3
    ELSE statement-4
```

Execution of this nested construction can be shown graphically as follows:

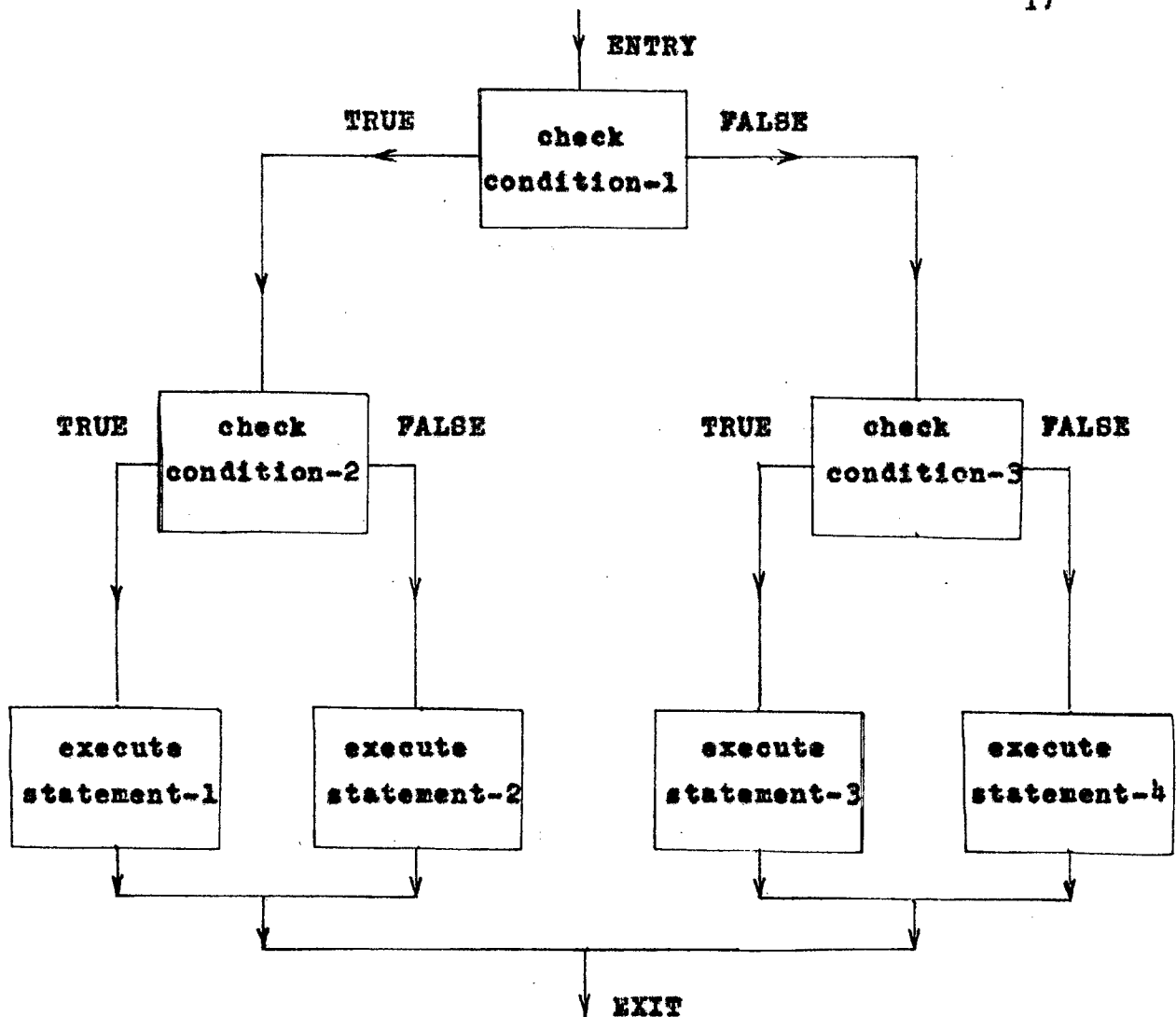


FIG. 2.3.1b

As in most of the languages, the above construction is ambiguous - the ELSE could be associated with either IF. BEGIN-END can be used to disambiguate this as desired. In absence of BEGIN-END, each ELSE goes with the previous un-ELSEd IF, just as in PL/I.

EXTFOR does not provide a CASE statement, since it may be readily simulated with a series of ELSE IF statements as:

```

IF (condition) THEN
  ELSE IF (condition) THEN
    ---
  ELSE IF (condition) THEN
    ---
  ELSE
    ---
  
```

2.4 LOOP STRUCTURES

These control statements are used to form loops; for both conditional and unconditional iterations. The flow of control in these control structures is pictured in the following figure.

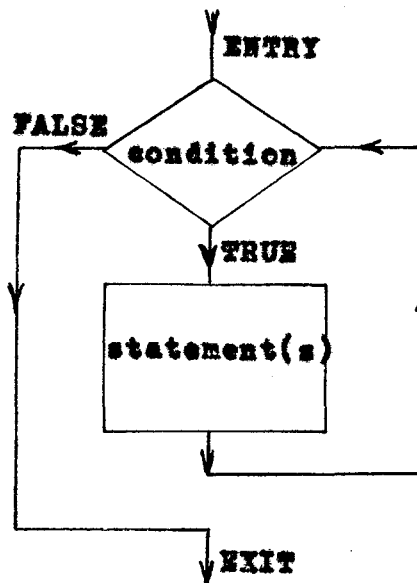


FIG. 2.4

The diagram has a decision symbol following the entry. Depending on the outcome of the test (value TRUE or FALSE), the flow continues either to the processing statement(s), the flow loops back to the entry connection; or exit from the loop. Looping is the objective of this control structure.

The number of times the processing statement(s) called the scope of loop, executed is determined by the specified conditions.

Four loop structures are selected for EXTFOR: DO-statement, WHILE-statement, FOR and REPEAT-UNTIL statements.

2.4.1 THE DO STATEMENT

The format is:

```
DO  index=i1,i2,step
    statements
END
```

The function of this DO is identical to the FORTRAN DO loop. The difference is that it does not use any statement number and the loop scope is limited to the matching END statement. It helps to avoid use of statement label, which has to be remembered to terminate the loop scope.

The scope of the loop is executed with index = initial value (i1); then index is incremented by the step value and loop is executed till the index value exceeds the final value i2, with increments in index after each execution of loop. The diagram shows the flow of control:

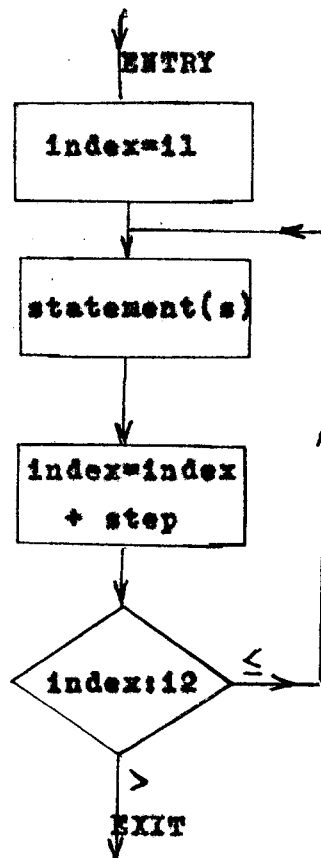


FIG. 2.4.1

Example:	equivalent FORTRAN DO
DO I=1,N	DO 10 I=1,N
X(I)=0	X(I)=0
Y(I)=0	Y(I)=0
END	10 CONTINUE

Nested DOs are permitted in EXTFOR provided the structures are ended with matching END statements.

Example:	equivalent FORTRAN DO
DO I=1,10	DO 30 I=1,10
DO J=1,10	DO 20 J=1,10
C(I,J)=0	C(I,J)=0
DO K=1,10	DO 10 K=1,10
C(I,J)=C(I,J)+A(J,K)	C(I,J)=C(I,J)+A(J,K)
*B(K,J)	*B(K,J)
END	10 CONTINUE
END	20 CONTINUE
END	30 CONTINUE

Like FORTRAN DO, EXTFOR DO is highly restrictive: the index must typically run from positive, non-zero integer value up to some positive integer limit, and in many versions of FORTRAN the loop is always obeyed at least once regardless of the limits, because the first test is done at the bottom. Jumps into a DO loop are forbidden. The jump out of the loop is by reaching, the statement following the loop after the loop index has executed its maximum value. However, jumps out of the DO loops to other parts of the program are of course permitted, of course this practice is discouraged.

2.4.2 THE WHILE STATEMENT

To overcome the difficulties with DO statement; EXTFOR provides the WHILE statement. The WHILE loop causes its scope to be repeatedly executed while a specified condition is true. As the following diagram shows the condition is checked prior to the first execution of the scope, thus if the condition is initially false the scope will not be executed at all. The format is:

```

WHILE (condition)
    statement-1
    statement-2
    .
    .
END
  
```

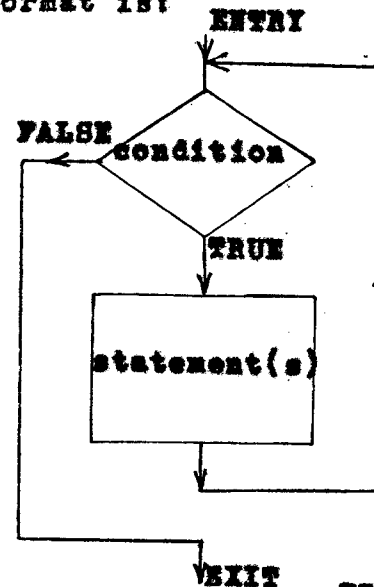


FIG. 2.4.2

The condition part is same as a valid FORTRAN condition. The scope of the loop is terminated by a matching end.

The nested WHILE statements are allowed provided the structures are ended with matching END statements.

Example:

```

WHILE (ABS (TERM) > 8)
    TERM = - TERM * X
    SN    = SN + TERM
    WHILE (SN > 0)
        SN = - SN * 2
        TERM2 = TERM2 + SN
    END
    TERM = TERM2 ** 2
END
  
```



The innermost END ends the innermost WHILE statement. TH-1222

2.4.3 THE FOR STATEMENT

This statement compactly summarizes all the loop control code in one line, in other words it separates the scope of the loop from the reason-for-looping. The FOR statement allows explicit initialization and increment step as part of the statement. The test part of loop is placed at the top of loop instead of the bottom which eliminates a potential boundary error. The format is:

```
FOR (initialize; (condition); reinitialize)
  statement
```

"initialize" is any single FORTRAN statement, which is executed once before the loop begins. "reinitialize" is any single FORTRAN statement, which is executed at the end of each pass through the loop. "condition" is any logical expression.

The diagram illustrates the flow of control.

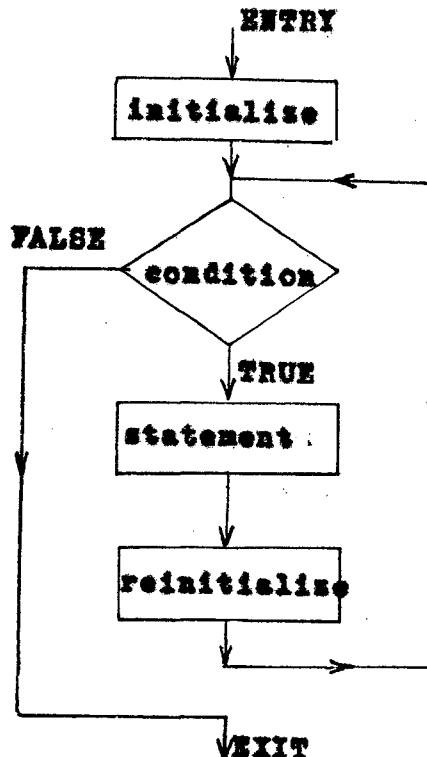


FIG. 2.4.3

The initialize statement is executed only once. Then the condition is checked, if it results a FALSE the statement following the FOR structure is given control. For TRUE value the scope of loop (statement) is executed followed by the execution of reinitialize. The loop continues from the evaluation of condition and further action depending upon the FALSE or TRUE value as discussed above.

The FOR statement is useful for backward loops, chaining along lists, loops that might be performed zero times, and things which are hard to express with a DO statement. For example, here is a 'backward DO loop':

```
FOR (I=80; (I.GE.50); I=I-1)
```

```
    IF (STR (I).EQ.BLANK) THEN NBLK=NBLK +1
```

The above code scans the character string STR(80) from last character to 50th character in the right to left direction.

Example:

```
TERM=1.0
```

```
FOR (I=2.35; (I.GT.0); I=FUN(I))
```

```
    TERM=TERM+VALUE (I)
```

From this example it is clear that the increment in FOR statement need not be an arithmetic progression. FUN and VALUE are functions, FUN generate random values for I.

In FOR any of initialize, condition and reinitialize may be omitted, although the semicolons must remain. For example

- (a) FOR (; (condition); reinitialize)
- (b) FOR (initialize; ;reinitialize)

```
(c)   FOR (initialize; (condition); )
(d)   FOR ( ; (condition); )
(e)   FOR ( ; ; )
```

Examples (b) and (e), where the condition is omitted represent indefinite repetition of loop scope execution. Examples (a) and (d) are similar to WHILE loop.

The nested FOR loops are allowed in EXTFOR, for example:

```
FOR (I=2; (BUF(I).NE.BLANK); I=I+1)
  FOR (J=I-1; (BUF(J).NE.EOS); )
    TOKEN (J)=BUF(I)
```

2.4.4 THE REPEAT-UNTIL STATEMENT

The REPEAT-UNTIL is a loop that is repeated one or more times until the condition test at the trailing end of loop results a TRUE. The format is:

```
REPEAT
  statement-1
  statement-2
  .
  .
  statement-n
UNTIL (condition)
```

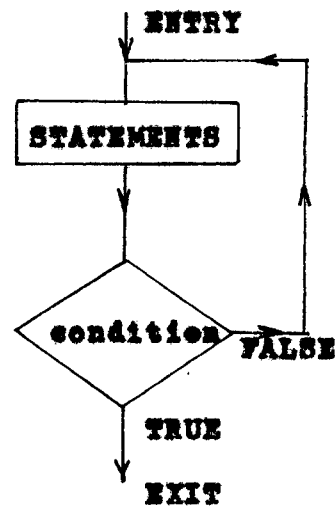


FIG. 2.4.4

The figure illustrates the flow of control within the structure. The scope of the loop (statement(s)) is executed first. The condition is checked and the FALSE value results in a repetition of the loop scope execution, which continues until the condition check results a TRUE value. This loop structure differs from others, because of the reverse condition for looping i.e. for exit out of the loop the condition must become a TRUE value. The loop is repeated atleast once, similar to the DO statement and the condition is examined only at the end of the loop. We shall remember, that the scope of the loop should modify the values of variables involved in condition checked at the end of loop such that after a finite number of repetitions the control comes out of the loop structure. For example the following statements result an indefinite loop.

```

NUMB=1
SUM=100
REPEAT
    SUM=SUM+NUMB
    NUMB=NUMB+1
UNTIL (SUM.EQ.100)

```

The value of SUM keeps on increasing by successive executions of loop scope and it is never equal to 100.

The condition part of UNTIL is optional; if it is omitted, the result is an infinite loop, unless the loop is broken by some conditional or unconditional branch from the loop scope. Example:

```

I=1
REPEAT
    STR1(I)=BUF(LASTC)
    IF (STR1(I).EQ.STR2(I)) THEN I=I+1
    IF (STR1(I).EQ.EOS) THEN BREAK
    LASTC=LASTC+1
UNTIL
NTOK=NTOK+1

```

The above loop is executed infinitely unless the underlined IF statement is there. When the condition of this IF results a TRUE, the loop is broken by BREAK statement transferring control to the statement NTOK=NTOK+1 which follows the loop.

2.5 THE BREAK STATEMENT:

As a policy we don't want to allow frequent use of GOTO's in EXTFOR programs since it causes a lot of confusion in program reading by creating more interactions among the program units. Therefore some statements are required to be included in our language to provide implied branches. The BREAK statement provides one way to branch out of a loop structure. Whenever this statement is encountered during the program control flow, the control is transferred to the statement, next to the loop structure. The format is:

BREAK level

where level is a positive integer indicates the depth of the loop nest to be broken. For example

```

WHILE (NGETC (C).NE.EOF)
  CALL PUTBAK (C)
  REPEAT
    IF((C.NE.BLANK) .AND. (C.NE.TAB)) THEN BREAK
    CALL PUTBAK (C)
    UNTIL
    I=1
    WHILE (I.LT.TOKS12)
      GETTOK = TYPE (NGETC (LEXSTR (1)))
      IF (GETTOK.NE.LETTER) THEN BREAK 2
    END
    GETTOK = ALPHA
  END

```

The execution of the first BREAK when the condition results a TRUE causes control to be given to the statement I=1 which just follows the REPEAT loop. Here the level is omitted and only one loop is terminated by default even if the BREAK is contained inside several nested loops. The second BREAK statement terminates two loops i.e. the control comes to the statement next to the outermost END.

In case the level specified is greater than the number of nested loops enclosing it, the execution of BREAK transfers control to the statement next to the outermost loop structure. This statement also provides one way to jump out of the BEGIN-END block. For example

```

-----
BEGIN
  SUBST = INJECT (NEW)
  IF (SUBST.EQ.ERR) THEN BREAK
  SUBST = OK
  LASTM = M
END
CALL DELET (LINE, LINE, STATUS)

```


When SUBST has value equal to ERR the BREAK statement causes an immediate exit from the block and control is given to statement CALL DELET.

The BREAK statement should not be used in an un-disciplined way, we should always try to express the exit condition, that is an IF controlling a BREAK. The BREAK statement without an enclosing loop is treated as a comment anyway an error message is flagged.

2.6 THE NEXT STATEMENT

This statement causes whatever loop it is contained in to branch to the bottom of the loop; skipping the rest of the loop body, so it causes the next iteration to begin. The format is:

NEXT level

In a WHILE, REPEAT or DO, it goes immediately to the condition part; in a FOR, it goes to the reinitialize step. "level" is a positive integer specifies the level of loop in a nest of several loops for which the next iteration to begin. The other loops inside the loop at "level" are terminated consequently.

Example:

```

DO I = 1, N
    IF (X (I) .LT.0.0) THEN NEXT
    SUM = SUM + X(I)
END

```

This loop skips over negative values to be added in SUM because whenever the value of X (I) is equal to 0.0 the NEXT causes the SUM = SUM + X (I) statement to be skipped and next iteration begins with next value of I. If I = N and X (I) 0.0 then the execution of NEXT terminates the loop. The "level" can be omitted which has default value as one.

Example:

```

FOR (I=1; (I.LE.10); I=I+1) BEGIN
  A (I) = B(I) **2 + 4*A(I)*C(I)
  DO J = 1, 10
    N = A (I) + B (J)
    CALL FIGR (N)
    IF (FIGR.GT.100) THEN NEXT 2
    STERM = STERM + A (5) * N
  END
  WRITE (6, 100) STERM
END

```

In this example when FIGR is greater than 100, the execution of NEXT causes, first to terminate DO loop and then the control is transferred to the initialize statement I=I+1 by skipping WRITE statement. In the next iteration DO loop is started with value J=1 as a fresh.

2.7 THE EXTFOR GRAMMAR

The syntax (grammar) is a set of rules by which a legal program in the language is written or recognized. The formal representation of a language by grammar has several advantages rather than describing the language with an informal description in words. The language specifications can be made fairly precise this way.

The language grammar is written as a finite non-empty set of rules, usually written as

$$U ::= X$$

where U is a symbol called non-terminal or syntactic entity. X as the right part of the rule is a non-empty finite string of symbols. The symbols appearing on right part may be a mixture of non-terminals and other symbols called terminals.

There are different notations used to describe the language grammar formally, the one used here is the Backus-Naur Form (BNF). In BNF non-terminals are written as symbols enclosed in corner brackets " \langle and \rangle " and sign " $::=$ " reads "is replaced by". Alternate ways of rewriting a given non-terminal are separated by a vertical bar " $|$ " (read "or"). Symbol $\left[A \right]_a^b$ represents the repeated appearances of A with minimum a times and maximum b times.

The grammar for the specifications of the Extended Fortran language is following one:

```

<program> ::= <statement>
           | <program> <statement>
<statement> ::= IF (condition) THEN <statement>
              | IF (condition) THEN <statement>
              | ELSE <statement>
              | BEGIN
                <statement>
              | END
              | DO limits
                <statement>
              | END
              | WHILE (condition)
                <statement>
              | END
              | FOR (initialize;(condition);reinitialize)
                <statement>
              | REPEAT
                <statement>
              | UNTIL (condition)
              | BREAK <level>
              | NEXT <level>
              | <label> <statement>
              | * END
              | OTHER
<level> ::= {<digit>}02
<label> ::= {<digit1>}01 {<digit>}0k
<digit> ::= 0|1|2|3|4|5|6|7|8|9
<digit1> ::= 1|2|3|4|5|6|7|8

```

The first rule says that a program is a statement, or a program followed by a statement. In other words, a program consists of one or more statements. A statement is one of the constructs represented by the rules other

than first. The vertical bar "|" indicates a choice of alternatives. All the underlined symbols are keywords; keyword is a symbol which has special meaning when used at a specified place in the program. The user is restricted from using these keywords as variable names, particularly the first symbol of the statement. Here EXTFOR grammar differs from FORTRAN since there is no keyword in FORTRAN. A statement consists of, for example keyword WHILE followed by parenthesized condition and a group statement followed by END. The definition of a statement is recursive in EXTFOR which is not allowed in FORTRAN for example, a nest of IF statements.

The '*' END' statement is used to indicate the end of subprogram, '*' is put in first column of the source line. The grammatical type OTHER represents all the standard FORTRAN statements. This type is of important simplification, for it frees EXTFOR translator from having to know very much about FORTRAN. Any statement which is not discussed in the preceding sections is recognized as OTHER. Therefore, any statement which does not begin with one of the underlined keywords (leaving the label apart), it must be an OTHER, and no real processing is needed on it.

"Level" is a 2 digit positive integer specifies the loop level to be broken by BREAK/NEXT statement.

"Label" is a standard FORTRAN 5 digit statement number less than 90000. The EXTFOR translator uses label numbers greater than 89999 as generated labels (discussed in chapter 3). Although these labels are relatively rare in EXTFOR programs, the grammar does allow them.

The format ^{of} source is same as that of FORTRAN. A 'C' in column 1 indicates a comment line, 1 to 5 columns for statement label, column 6 is continuation column and columns 73 to 80 are for documentation purposes and ignored by EXTFOR.

EXTFOR PREPROCESSOR3.1 PREPROCESSOR AS A TRANSLATOR

A program written in a high level language, say FORTRAN can not be executed by computer directly. A translator is required to translate the source program into a set of executable instructions (machine language) for the computer as shown in the following figure.

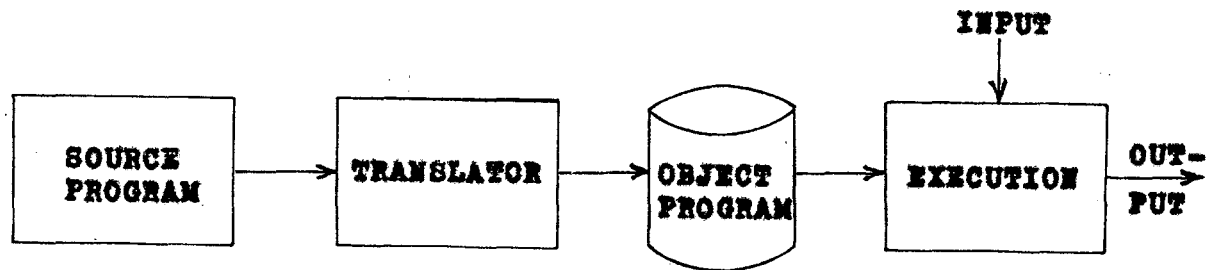


FIG. 3.1

Some translators accept a source program, translate it into some intermediate structure, then execute each operation given in the intermediate structure. Such a translator is called an interpreter. An interpreter is considerably less efficient in execution, but a program can be developed very fast with an interpreter.

A translator which accepts source program and translates it into machine executable code is called a compiler. The advantage of a compiler over an interpreter

is that it generates an efficient machine code. The high level language gives user the flexibility of writing programs independent of computer to be used. A FORTRAN program should yield the same results for a given input regardless of the to which it is translated. These results should be expected from the specification of the FORTRAN language and the algorithm as expressed in FORTRAN.

A program written in EXTFOR language has to be transformed into some code to make it executable or interpretable on a computer. One way to do it is, develop a compiler; which means too much of work involvement for a few extensions to FORTRAN language, which can be translated by a compiler. This approach will treat EXTFOR a new language. In situations when a user wants to avail some features which are not regular in a high level language like FORTRAN, same can be enhanced with a preprocessor. A preprocessor works on the extended features only and translates into the high level language. Thus the second approach is justified here for the translation of EXTFOR programs.

3.2 AN OVERVIEW OF EXTFOR PREPROCESSOR

The EXTFOR preprocessor which will be referred as "EXTFOR" is a program written in FORTRAN, translates programs of EXTFOR language defined by grammar in section 2.7, to programs expressed in FORTRAN as object language without changing the logical flow of the program. The preprocessor functions as an interface between the source of EXTFOR program and the FORTRAN compiler as shown in the figure.

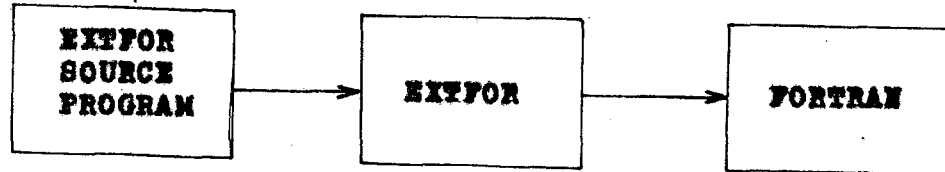


FIG. 3.2a

To translate a source program, the preprocessor must analyze it thoroughly and synthesize an equivalent object program.

During preprocessing that is EXTFOR to FORTRAN translation the process of source program analysis yields a variety of information about source program. This information is preserved by the preprocessor for it to produce an equivalent Fortran program. Various data structures, such as tables, lists, stacks etc. are employed by the preprocessor to preserve this information. The construction of an equivalent Fortran program is directed by the information preserved in these data structures.

The source program to preprocessor may be on any media such as card deck, disk, magnetic tape etc. The outputted object FORTRAN program is written by preprocessor on disk so as to make fast reading by FORTRAN compiler for the compilation into machine language. The complete process of preprocessing involves two steps, source program analysis and code generation.

The following figure (FIG.3.2) illustrates the logical parts of the EXTFOR preprocessor.

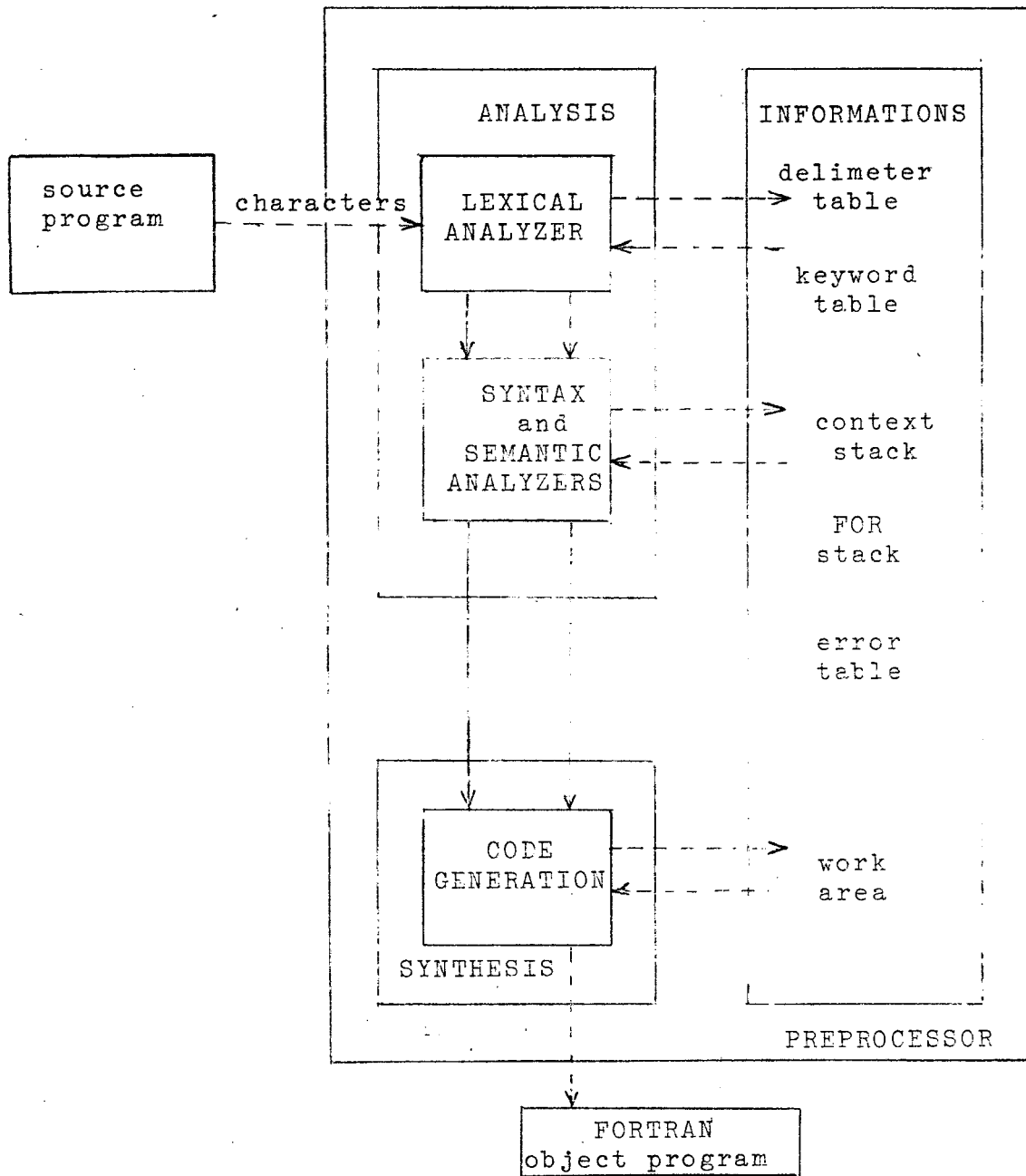


FIG. 3.2 Logical Parts of Preprocessor

3.2.1 SOURCE PROGRAM ANALYSIS

As input to the EXFOR, the source program is only a string of characters. From this linear representation of the program, the process of analysis should detect the structure and meaning of the program. This is very similar to the actions involved in finding constituent phrases in a sentence in English, in which case we do so by making use of the English grammar. From this analogy, the grammar defined in section 2.7 is used to recognize and formulate the constituents of the source program. The structural analysis of source program proceeds in two logical phases.

FIRST PHASE - LEXICAL ANALYSIS

The first phase consists of the analysis of the string of characters in the source program so as to form meaningful primitives (analogous to words and punctuation marks in an English sentence). Examples of such meaningful primitives are identifiers, operators like '+', '-', '*' etc. and separators or delimiters like ",", blank ' ', etc. The task of first phase is usually termed lexical analysis. Subsequent to this phase of analysis, a source program may be viewed as a sequence of meaningful primitive items referred to as lexemes or tokens.

SECOND PHASE - SYNTAX & SEMANTIC ANALYSIS

Following lexical analysis is the task of recognizing grammatical phrases in a source program. This is a more complicated task; sequences of tokens are grouped together to form simple structures of the source language;

these simple structures are used to form more complex ones and, ultimately, source programs. All these actions are the substance of syntax analysis also called parsing. The end result of this analysis is that we have discerned the grammatical structure of the source program. The routine called parser, analyses the grammatical structure of input, recognizes the statement of EXTFOR by seeing the first token of the statement. For instance, when an IF is seen the parser calls a routine which handles IF statements, which does the further semantic checking of the statement. When the parser recognizes a source language construct it calls a so called semantic procedure which takes the construct, checks it for semantic correctness that is the meaning part of the construct, and stores necessary information about it into context stack. This stack contains information about the scope, recursion and nesting of structures. The semantic procedure of the preprocessor is very simple since the most of the semantic checks are performed by the FORTRAN compiler at compilation level. After the syntax and semantic checks the next step is the code generation.

3.2.2 CODE GENERATION

This is the actual translation of the EXTFOR source program into FORTRAN language. In this section we will discuss the FORTRAN codes for the structures selected for EXTFOR. The preprocessor while generating codes generates statement labels, therefore a label slot 90000 to 99999 is reserved and the users are restricted to use them. In this section anything enclosed in square brackets [and] is optional.

BEGIN statement

If a BEGIN is encountered with a label then we replace BEGIN by

Label CONTINUE

No code is generated for unlabelled BEGIN. A label L is reserved and when an END terminates the BEGIN block, it is replaced by

L CONTINUE

If a BREAK or NEXT statement is enclosed by a BEGIN and the level of loop/structure corresponds to BEGIN then the statement BREAK or NEXT is replaced by

GO TO L

IF statement

The translation of

IF (condition) THEN statement

is something like

IF (condition is not true) go around statement

Thus when an IF is encountered, we

isolate the condition part

generate and save some unique label L

output "IF (.NOT.(condition)) GO TO L"

The .NOT. inverts the value of the condition. The statement label with IF (if any) is put as it is. When we get to the end of the statement that follows THEN, there are two possibilities:

If there is no ELSE following, we need output only

L CONTINUE

If an ELSE follows, however, we must generate another label and output

L1 GO TO L2

L CONTINUE

to branch around the ELSE part, and then, after whatever

statement follows the ELSE, we output

```
L2      CONTINUE
```

to terminate the IF-THEN-ELSE construction, the extra label L1 is required because there could be a statement like BREAK, NEXT or GO TO as the THEN part, then FORTRAN compiler flags an obvious error since any statement following unconditional branch has to be labelled. In summary, the code generation for

```
[label]IF (condition) THEN statement
```

is

```
[label] IF (.NOT.(condition)) GO TO L
      statement
```

```
L      CONTINUE
```

and for

```
[label] IF (condition) THEN statement1
      ELSE statement2
```

is

```
[label] IF (.NOT.(condition)) GO TO L
      statement1
```

```
L1     GO TO L2
```

```
L      CONTINUE
```

```
      statement2
```

```
L2     CONTINUE
```

EXTFOR generates three consecutive labels when an IF is seen: L1 is L+1 and L2 is L+2. If one of the labels is not used because there is no ELSE, it doesn't cost anything. As we know that the labels are always L, L+1 and L+2, only one of them need be remembered; the others are deduced by adding 1 and 2 respectively.

DO statement

The EXTFOR DO is a FORTRAN DO without a label. When a DO is encountered, we

isolate the "limits"
generate a label L
output "DO L limits"

No check is performed on "limits" by EXTFOR and left for FORTRAN. Then at the end of DO scope indicated by an END, we output

L CONTINUE

Thus the EXTFOR

```
[label] DO "limits"
      statements
      END
```

is translated into

```
[label] DO L "limits"
      statements
L      CONTINUE
L+1    CONTINUE
```

The second CONTINUE has to be produced in case the loop contains a BREAK statement. In this case we generate the second CONTINUE regardless of whether or not there is a BREAK; it is simply too much effort to check. To make out code generation task easier, we take advantage of the fact that FORTRAN compilers are usually quite clever about dealing with unreferenced CONTINUES.

The enclosed BREAK statement is replaced by

GO TO L+1

and the NEXT by

GO TO L

which causes next loop iteration to start.

WHILE statement

When a WHILE is encountered, we isolate the condition part generate and save 3 unique labels L, L+1 and L+2
 output "L+1 "IF (.NOT.(condition)) GO TO L+2"
 when we get an END terminating the WHILE, we output

```

L      GO TO L+1
L+2    CONTINUE

```

Thus the EXTFOR

```

[label] WHILE (condition)
      statements
      END

```

is translated into

```

[label] CONTINUE]
L+1    IF (.NOT.(condition)) GO TO L+2
      statements
L      GO TO L+1
L+2    CONTINUE

```

the CONTINUE before 'IF' is for the labelled WHILE as in:

```

10     WHILE (LASTC.LT.73)

```

...

The BREAK statement to break the WHILE loop is translated into

```

GO TO L+2

```

L+2 is the exit point of the loop. The NEXT statement terminating the current iteration and starting the next is translated as

```

GO TO L+1

```

FOR statement

When a FOR statement

```

[label] FOR (initialise; (condition); reinitialise)
      statement

```

is encountered, we

```

output [label] initialize
isolate the condition part
preserve the reinitialize part
output "L IF (.NOT.(condition)) GO TO L+2"

```

Then at the end of the statement associated with FOR, we output

```

L+1      reinitialize
          GO TO L
L+2      CONTINUE

```

Thus the code generation for FOR statement is:

```

[label] initialize
L        IF (.NOT.(condition)) GO TO L+2
          statement
L+1      reinitialize
          GO TO L
L+2      CONTINUE

```

Three labels L, L+1 and L+2 are reserved for a FOR statement. Here we note that reinitialize can not be a branch statement since GO TO L is unlabelled.

If "statement" in FOR is a BREAK which causes exit from FOR structure, it is translated into

```
GO TO L+2
```

similarly the NEXT statement which transfers control to the reinitialize part is translated into

```
GO TO L+1
```

REPEAT statement

When a REPEAT is encountered, we
output L CONTINUE

Then at the end of the loop which is identified by UNTIL (condition), we

```

isolate the condition part, if present
output "L+1 IF (condition) GO TO L
L+2 CONTINUE"

```


Thus the

```
[label] REPEAT
      statements
      UNTIL [(condition)]
```

is translated into

```
[label CONTINUE]
L      CONTINUE
      statements
L+1    [IF (condition)] GO TO L
L+2    CONTINUE
```

The first CONTINUE is generated only if REPEAT is labelled. In case the loop is infinite that is condition part is omitted "L+1 IF (condition) GO TO L" is replaced by "L+1 GO TO L". Three generated labels are reserved for a REPEAT statement. The BREAK statement encountered in the REPEAT loop to cause the exit from loop, is translated into

```
GO TO L+2
```

similarly NEXT statement causing control transfer to the UNTIL part, is translated into

```
GO TO L+1
```

Labels and Others

The statement label given by the user in EXTFOR source program is output at column 1 and followed by enough blanks that the next character will come in column 7. Input statement of lexical type OTHERS is copied from input to output without any change.

The error detecting abilities of EXTFOR are not as good as they might be with a more comprehensive grammar. This is not a serious drawback, however, since we are translating into FORTRAN, and FORTRAN compilers are perfectly capable of detecting any syntax errors that escape the pre-processor. The errors detected by EXTFOR are output as comments in the listing of program.

3.3 EXTFOR PREPROCESSOR ORGANISATION

The organisation of EXTFOR preprocessor is shown in FIG.3.3. The top level routine **PARSER** controls by analyzing the grammatical structure of EXTFOR source program as input it sees. It employs other routines of lower level to help in analyzing, storing informations and finally outputting the FORTRAN code.

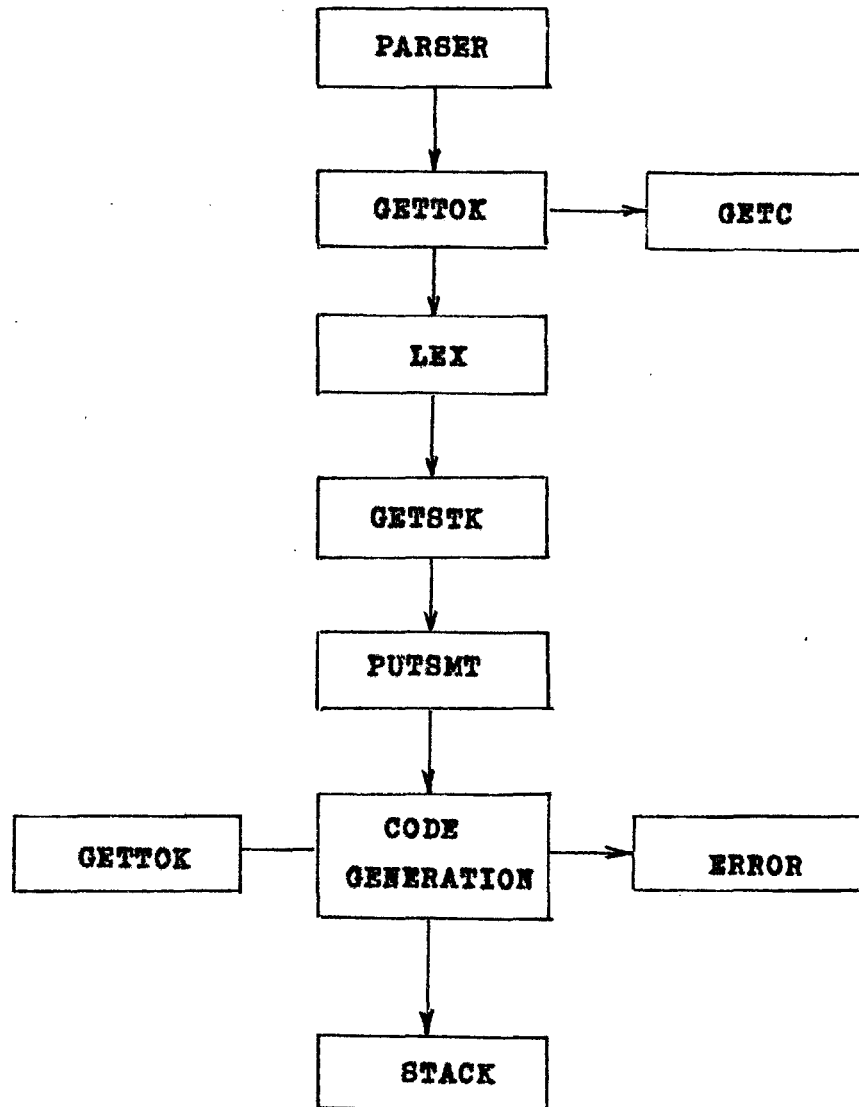


FIG. 3.3

The PARSER calls GETTOK to get the next token of the statement in the input buffer. GETC is the lowest level routine which gets a character from the input buffer to GETTOK. GETTOK constructs a token; LEX is called by PARSER to work on the first token of the statement and the token is identified as one of the keywords. This way PARSER determines the statement type. GETSTK routine tells about the context of the current statement which helps in analysis and proper code generation. For instance, when an IF is seen PARSER calls a routine which handles IF statements. That routine in turn isolates the condition part and makes an entry in the context stack with help of STACK. Necessary codes are generated. When the end of the "statement" part of THEN unit is reached, the correct terminating code for an IF can be produced by calling GETSTK to know about the IF and the generated labels reserved for it. This also include dealing with an ELSE if it is present.

PUTSMT outputs the text of the statement presently being processed as a comment for the user's reference. The FORTRAN codes are output just after the text of source, with "++" in columns 73 and 74 of the object line.

There are number of code generation routines called by PARSER; simple codes are generated by PARSER itself. Some code generation routines also use GETTOK to read further parts of the statement being processed. ERROR outputs the appropriate error message detected during the analysis of the current statement. The calling routine passes the error code to ERROR to output the error message from Error Table.

Now, we describe the various buffers, tables and stacks used to store the informations, lexical analysis, parsing, code generation and general purpose routines.

The following discussion is based on the source listing of EXTFOR preprocessor (Appendix A).

3.3.1 BUFFERS AND BLOCKS

A number of buffers are defined in the work area of EXTFOR, used to keep the parts of source program as well as object program for the current processing. The work area is a set of FORTRAN COMMON blocks.

INBLK Input buffer, COMMON block of 80 bytes. It holds one line of source statement at a time. When the line is processed, next line is read into this block.

INPTR A block which contains the pointer to the last character referred in INBLK.

TKN Holds the current token under processing. The token is constructed by GETTOK.

TKSIZE It contains the number of characters in the token TKN.

INLBL The common block for the statement label of source statement under processing.

CONBLK The condition buffer to contain the condition part of control structures like IF, WHILE, UNTIL and FOR. The condition part is separated out from the statement and put into this buffer till the code "IF (.NOT.(condition))" is generated. CONBLK can contain as long as 17 lines of condition.

CONPTR A pointer used to scan the condition CONBLK.

ERRNO The serial number of error in a statement, detected by EXTFOR.

NOERR It contains the total number of errors encountered in the source program at preprocessor level.

RTPLAB The next generated label available to EXTFOR for code generation is put in this block by function LABGEN which reserves next N labels by incrementing label in

RTFLAB by N and returning old value in LABGEN. It is referred as LABGEN(N).

3.3.2 KEYWORD TABLE

A table is an ordered set of elements of the same type. The ordering of elements in table is used to access the elements, by means of one or more subscripts. The following keywords are stored in form of a table. Each keyword is associated with a unique number which is returned by function LEX and called the lexical type. The lexical types are given negative values to avoid any confusion for LEX, with a printable character.

<u>Keyword</u>	<u>Lexical Type</u>
IF	-30
DO	-31
WHILE	-32
BEGIN	-33
FOR	-34
REPEAT	-35
ELSE	-36
END	-37
NEXT	-38
BREAK	-39
UNTIL	-40
OTHERS	-41

The OTHERS is not a keyword, it corresponds to a token which is not identified as any of the preceding types. This category actually encompasses most of the FORTRAN statements. The function LEX returns -30 as lexical type for IF when referred as LEX(IF).

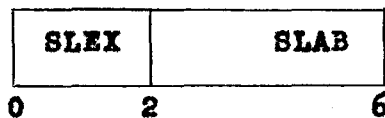
3.3.3 ERROR TABLE

This table contains the text of all the error messages with their codes. The routine ERROR uses this table to output

the error message .

3.3.4 CONTEXT STACK

A stack is a data structure in which elements are arranged in an order having only one end open for addition and deletion of data elements. Stacks are widely used in systems programming for housekeeping functions needed for recursive programs or structures, block structures, and bracketed expressions, etc. The context stack is used for these purposes except the expression processing since this work we leave for FORTRAN compiler. Each entry in context stack is of 6 bytes.



SLEX is the lexical type of the keyword and SLAB is the generated label preserved in the context stack.

Whenever either of the statements IF, ELSE, BEGIN, DO, WHILE, FOR and REPEAT is encountered, the lexical type and the generated label reserved for the structure are pushed on the stack. Whenever a statement ends the context stack is checked for the structure to generate codes for the exit point. For example, when a BEGIN is encountered the lexical type -33 for BEGIN and the generated label L is pushed down in the stack. Later on seeing the matching END statement the context stack is checked; if the stack does not contain BEGIN or anything, an error is flagged for unexpected appearance of END in the source unless it is the last statement of a subprogram. For BEGIN the entry from stack is removed and code "L CONTINUE" as the end point of the block structure is output.

A pointer called stack pointer (sp) is used to point

at the top entry on the stack. To delete an entry the pointer is decreased by 1 and to add *sp* is incremented by 1 with proper contents on the stack entry. The common block *STKBLK* defines the area for context stack with a maximum of 100 entries and *STKPTR* is the stack pointer(*sp*). The routine *STACK* pushes down an entry in the stack. The lexical type and the generated label are passed as parameters to *STACK*. *GETSTK* gets the top most entry and puts the contents in the arguments *SLEX* and *SLAB*.

3.3.5 FOR STACK

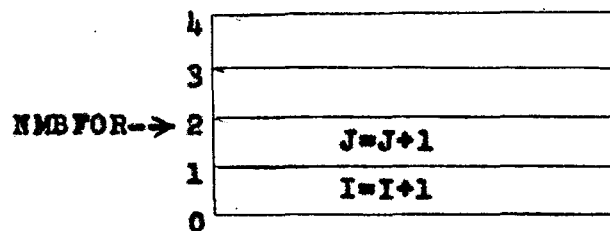
This stack is defined as the *COMMON* block *RIBLK* and the stack pointer *NMBFOR*. Whenever a *FOR* statement is encountered the reinitialize part is preserved in this stack, which is used to generate code for end of *FOR* statement.

The reinitialize part is limited to a length of 66 character i.e. one line of source. The size of this stack is four entries; thus limits the maximum number of nested *FOR* to 4.

Example:

```
FOR (I=1; (I.LT.100); I=I+1)
  FOR (J=1; (J.LT.10); J=J+1) ....
  ‡
```

when the input buffer pointer points at ‡ the *FOR* stack looks like



3.3.6 LEXICAL ANALYSIS

A token in EXTFOR is a character string constituted by the elements of FORTRAN character set except the (,), comma and blank. These four characters are significant delimiters for EXTFOR and separate the tokens in the source line. To construct a token GETTOK calls GETC to read the source line; GETC supplies characters one by one to GETTOK. GETTOK checks the type of the characters returned by GETC, if it is a delimiter then the character string formed by GETTOK in TKN is the token. Otherwise, the GETTOK keeps on adding the characters to TKN unless a delimiter or the end of source is encountered. At the end of source, EOF with value -1 is returned by GETC as character to GETTOK; also a flag EOSRC is set, used by PARSER to stop the processing.

Since GETTOK uses blanks to separate tokens, blanks are significant in EXTFOR which is not there in FORTRAN. The keywords like WHILE must not contain blanks, otherwise they won't be recognized.

GETC while reading the source preserves the statement label which is any thing appears before column 6 of the source line, unless it is a continuation line, in COMMON block INLBL. This label is output later on in the translated code. No check is made on the validity of statement label and left for FORTRAN compiler to do the check.

GETTOK also finds out whether the token is the first token of a statement. The function LEX takes this token and determines the lexical type of token by comparing it with all the keywords. This lexical type is used by PARSER in calling the appropriate routine to generate the code.

3.3.7 PARSING AND CODE GENERATION

When the beginning of a statement is encountered and recognized as one of the IF, ELSE, BEGIN, DO, WHILE, FOR and REPEAT, the corresponding lexical type and generated label are pushed on the context stack and code generation routine for that type is called after outputting the text of the source statement as a comment. At the end of the statement indicated by type END or BREAK, NEXT, or OTHERS, an appropriate code generation routine is called. The PARSER also checks the stack entry at the top and may be able to pop one or more things of the stack depending upon the lexical type on the stack. For example in

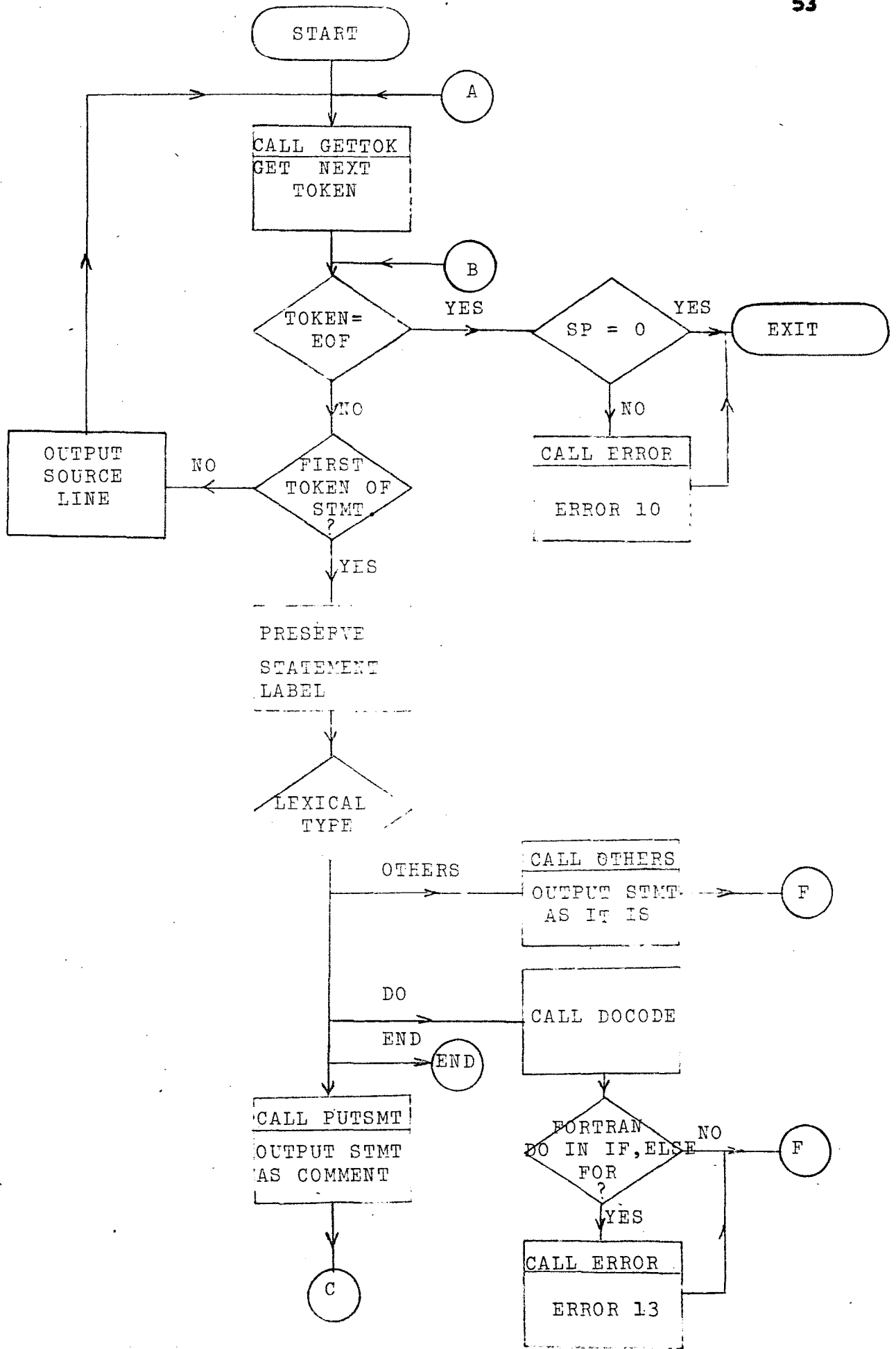
```

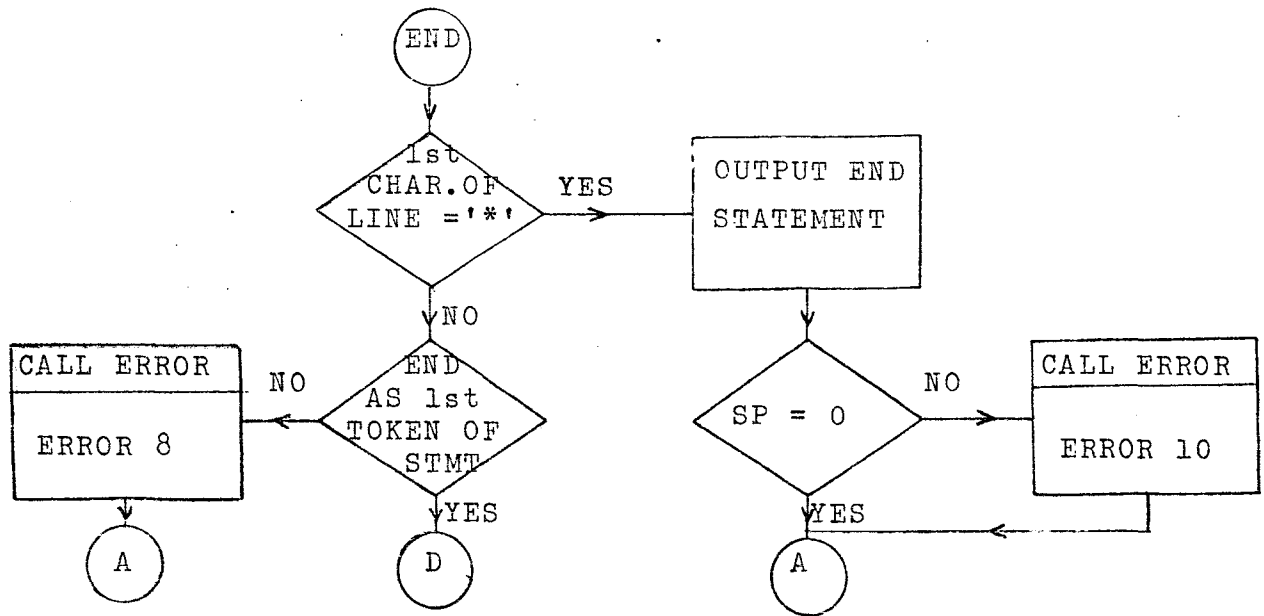
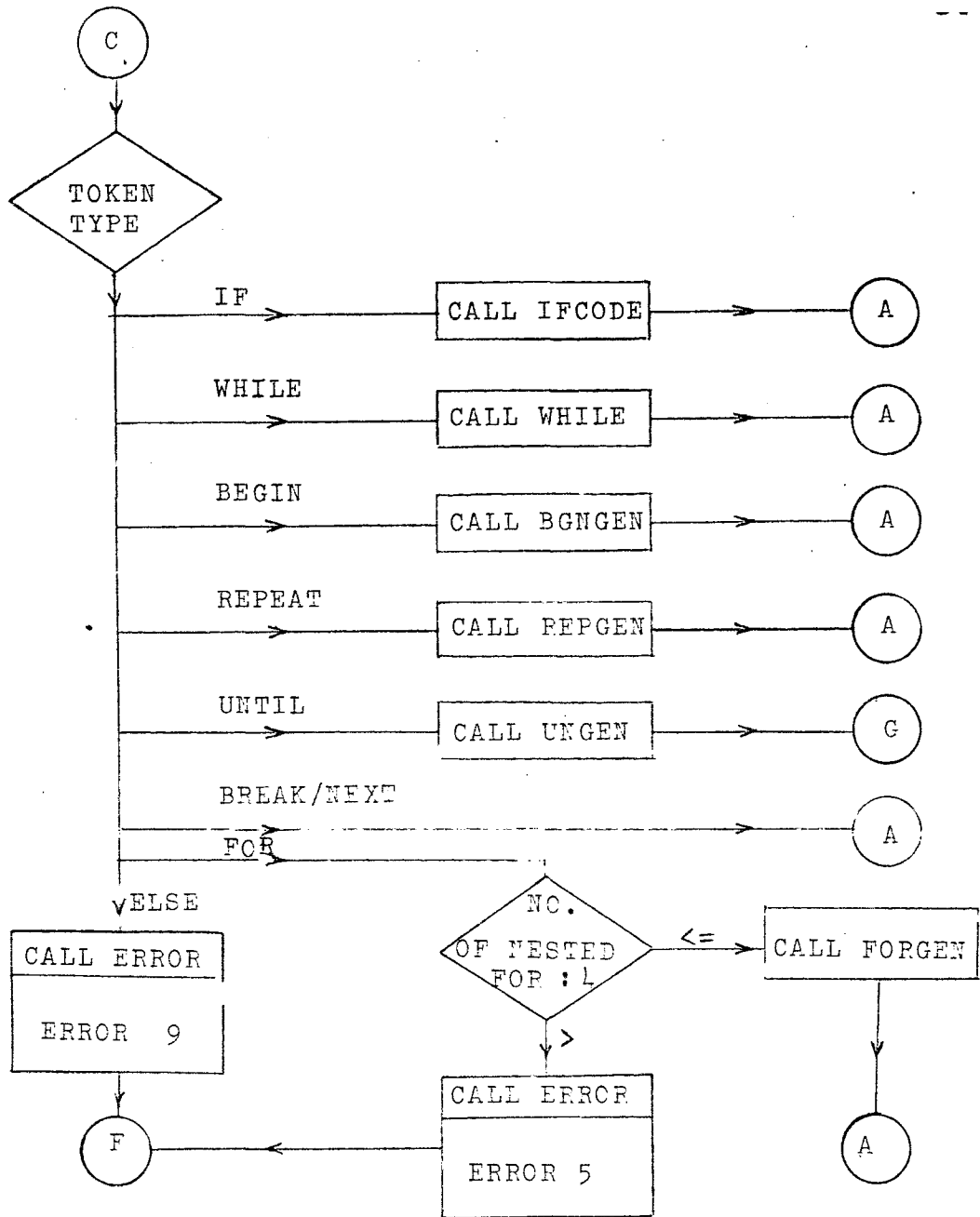
IF (cond-1) THEN
    IF (cond-2) THEN M=M+1
    ELSE M=M+2
LASTC=LASTC+1

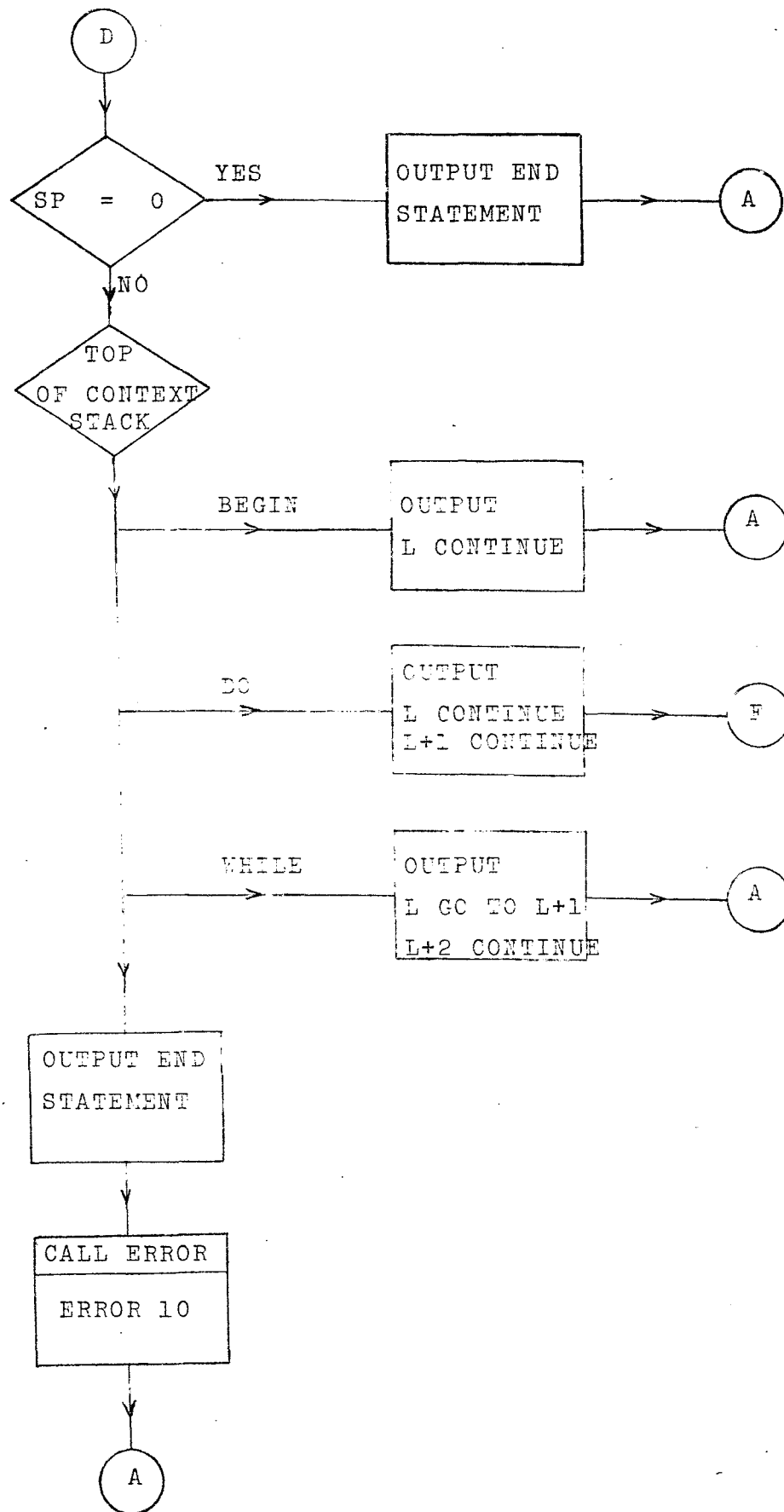
```

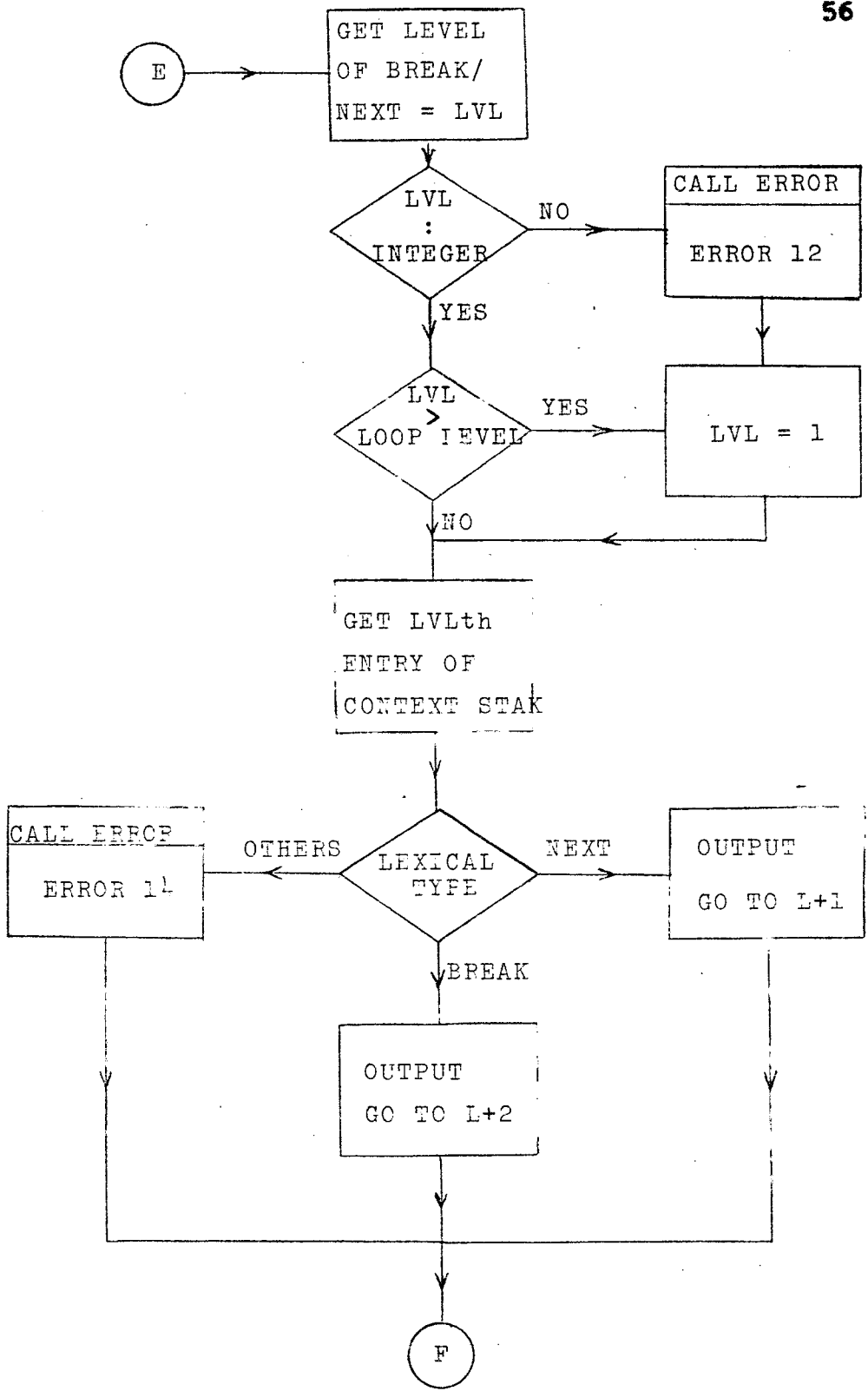
when LASTC=LASTC+1 is seen (lexical type OTHERS), there are entries for ELSE and both IFs on the stack. PARSER pops ELSE and both IFs; because there is no following ELSE matching with first IF, both IF statements are finished. The following ELSE is checked by looking ahead one token.

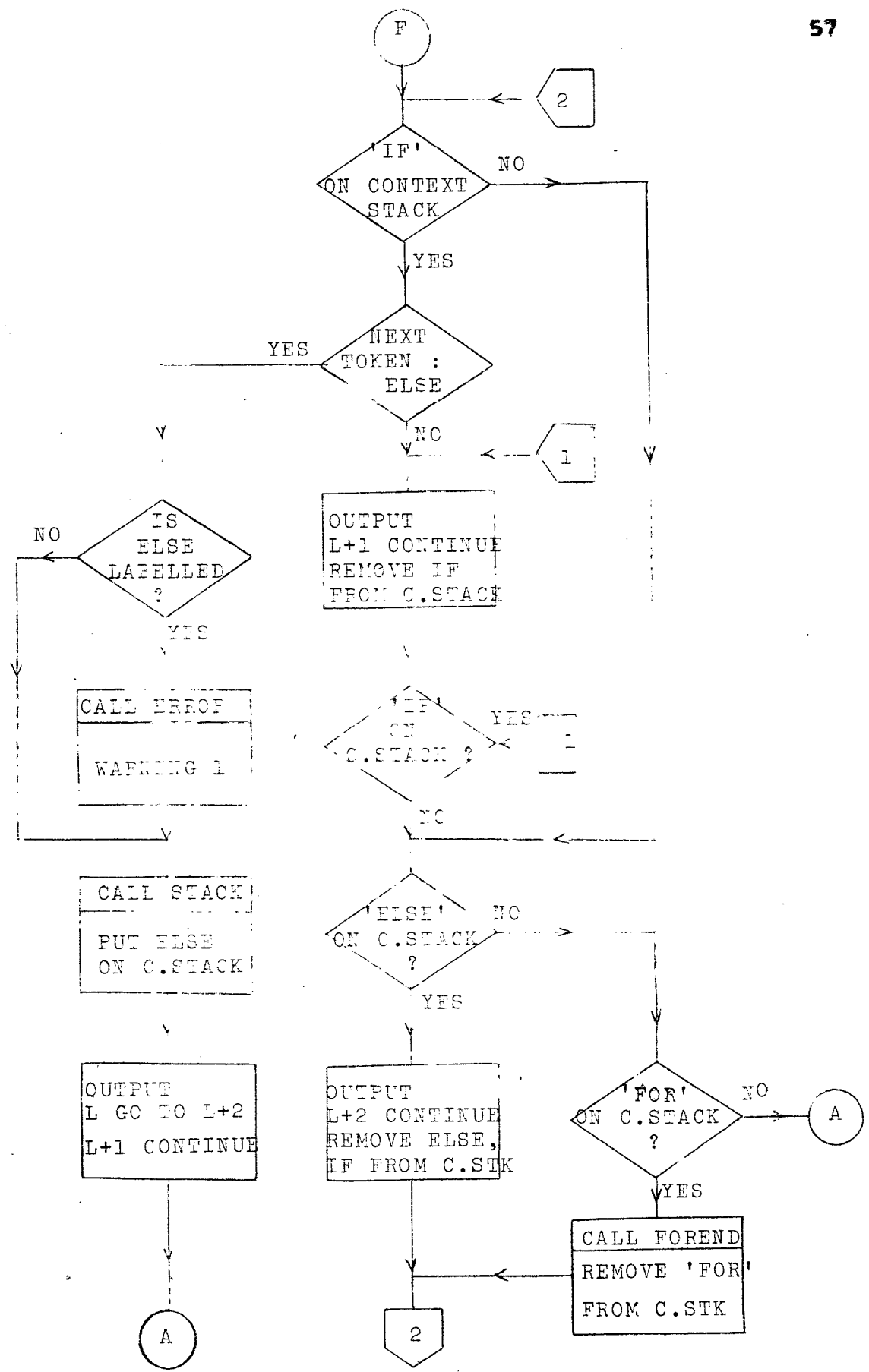
The following flow-chart of PARSER and descriptions of the routines give a close idea how EXTFOR works.











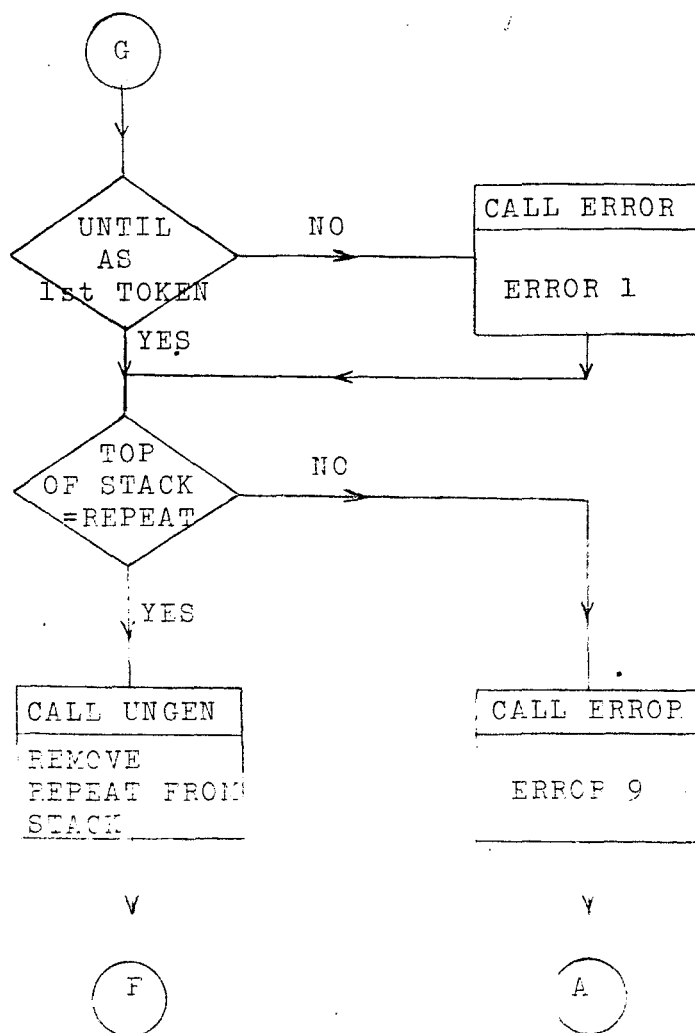


FIG:3.3.6 Flow Chart of PARSEP

ROUTINE DESCRIPTIONS

ERROR This routine outputs the error messages (Appendix D) when the error code is supplied as a parameter.

OTHERS This routine outputs the statement which is not any of the extended features.

DOCODE This routine processes DO statements. The token following 'DO' is checked, if it is an integer number then the DO is a FORTRAN DO and it is output as such. Otherwise, the 'limits' which follows DO is collected, two labels L and L+1 are generated and DOCODE outputs

DO L 'limits'

BGNGEN This routine generates one label L for the BEGIN. If BEGIN statement is labelled "Label CONTINUE" is output. The lexical type of BEGIN is put on the context stack.

BALPAR This routine collects the condition part of the control structure, which is a string enclosed in balanced parentheses. If this is spreaded over several lines BALPAR handles continuations upto 17 lines and puts in COMMON block CONBLK with the number of characters in condition, in COMMON block CONPTR.

PUTIF This routine is referred as CALL PUTIF (TYPE, LAB). In case value of TYPE is passed as 1 PUTIF outputs

```
IF (.NOT.(condition)) GO TO LAB
```

when TYPE=0, PUTIF outputs

```
IF (condition) GO TO LAB
```

IFGO The routine processes the condition part of IF, WHILE, UNTIL and FOR statements. If IFGO is called with option CNTRL=1, it checks for FORTRAN logical and arithmetic IF statements and outputs without any change. With option CNTRL=0, it calls BALPAR and PUTCHD to put the condition part in condition buffer CONBLK and, then calls PUTIF to generate the code

```
[label] IF (.NOT.(condition)) GO TO L
```

IFCODE Routine is called to process the IF statement. Three generated labels are reserved and IFGO is called. If the statement is found to be anyone of the FORTRAN IF statements control is returned to PARSER after outputting the statement as it is. For 'IF' statement the code

```
IF (.NOT. (condition)) GO TO L
```

is output and lexical type IF and the label L are put on the context stack.

WHILE Routine generates FORTRAN code for WHILE statement. Three generated labels are reserved for each appearance of WHILE. Routine outputs

label CONTINUE

if statement is labelled and calls IFGO to output

L+1 IF (.NOT.(condition)) GO TO L+2

The lexical type WHILE and label L are put on the context stack.

FORGEN This routine processes the statement:

FOR (initialize; (condition); reinitialize) statement

The procedure followed is:

Call PUTEND and BALPAR to collect source between outermost parentheses, in condition buffer CONBLK.

Scan CONBLK to find the positions of semicolons as PTR1 and PTR2; if number of semicolons is not equal to 2 then, call ERROR(6) and output the statement; return control to PARSER.

Output the initialize part which appears before first semicolon as [label] initialize

Push down the reinitialize part which appears after second semicolon, on FOR stack.

Shift the condition part which appears between two semicolons i.e. PTR1 and PTR2 positions, to the first character position of CONBLK and set pointer CONPTR=PTR2-PTR1.

Call PUTIF to output

L IF (.NOT.(condition)) GO TO L+2

Stack the lexical type FOR and generated label L.

Return control to PARSER.

WHILE

FOREND This routine generates FORTRAN code for end of FOR structure. The code output are

```

      L+1      reinitialize
              GO TO L
      L+2      CONTINUE

```

The reinitialize part is taken from FOR stack and stack pointer is decremented by 1. The FOR entry from context is also removed.

REPGEN The routine reserves three generated labels for each appearance of REPEAT and outputs

```

      Label   CONTINUE   (if REPEAT is labelled)
      L       CONTINUE

```

The lexical type REPEAT and generated label L are stacked.

UNGEN Processes UNTIL part of REPEAT structure. In case UNTIL is preceded by label ERROR(1) is flagged. If there is no token which follows UNTIL, the loop is indefinite and UNGEN outputs

```

      L+1      GO TO L
      L+2      CONTINUE

```

returns the control to PARSER.

For finite loop UNGEN calls BALPAR and PUTCND to collect the condition part of UNTIL. Finally PUTIF is called to output

```

              IF (condition) GO TO L+2
      L+1      GO TO L
      L+2      CONTINUE

```

The entry for REPEAT from context stack is removed and control is returned to PARSER.

There are some small routines called by PARSER, their functions and algorithms are quite simple, one can refer EXTFOR source listing (Appendix A).

EFFICIENCY AND OVERHEADS IN PREPROCESSOR

4.1 EFFICIENCY

The efficiency increases considerably while writing programs in EXFOR but the additional translation step from EXFOR to FORTRAN increases the cost of producing an executable version of the program.

One of the major factors affecting the efficiency of preprocessor is the features selected, to be translated by preprocessor. More sophisticated features and lot of improvements attribute to the inefficiency of preprocessor. The best way is to select a small set of improvements, depending on the user's requirements. A better design and implementation also make the preprocessor an efficient tool.

The first FORTRAN preprocessor RATFOR developed by Dr Kernighan is found to be extremely inefficient. The cost ratios of a RATFOR and FORTRAN programs in terms of CPU times, are from 10:1 for small programs to 19:1 for RATFOR compiling itself. During the final testing of EXFOR and the test programs attached in Appendix C, it is found that same ratio for EXFOR is 1.1:1 to 2:1; a substantial improvement in efficiency of preprocessor.

The efficiency can be improved further by identifying some of the routines as 'high spot' in preprocessor and record them to improve efficiency. Dr Kernighan measured that about 60% time is spent doing input and output at the lowest level. A major part of remaining 40% time is shared by GETTOK. We might improve the efficiency by concentrating on the better design of these

'high spot' routines. User also has to pay penalty in terms of run-time efficiency of the code produced by the preprocessor and the storage requirements.

The straightforward translation process from EXTFOR to FORTRAN makes little attempt to produce optimal FORTRAN code. Sometimes unnecessary GOTOs and CONTINUEs are produced and left for FORTRAN compiler to optimize and avoid the redundancy in FORTRAN code. In the early stages of the idea of FORTRAN preprocessor, it was thought that preprocessed and hand coded versions of a program would not display significantly different time and space characteristics if both were compiled using a good optimizing compiler even if the program involved a moderate amount of input/output. Later on, a number of experiments were done on several preprocessors, compilers, and machines and following facts were concluded.

Two factors would affect overall time and space characteristics - the quality of optimization performed by the FORTRAN compiler and the amount of input/output performed.

4.2 EXECUTION TIME OVERHEAD

With moderate input/output activities, the run-time overhead incurred by using a preprocessor is insignificant. With output removed, the results change considerably. The execution time overhead ranges from 5 to 30 percent. When a compiler's optimization is turned on interesting things were observed. For programs with a moderate amount of output, there is an insignificant change in execution time overhead with improved optimization. Surprisingly, in several cases the overhead

is increased by optimizing the code. This behaviour is found to be different for different compilers. With output removed the run-time overhead increases with greater optimization! Except for a few preprocessors and compilers.

4.3 STORAGE OVERHEAD

The storage penalty incurred by using a preprocessor seems highly dependent on the style of the generated FORTRAN code and on the sophistication of the preprocessor features. A preprocessor with more sophisticated language features has more storage overheads than one with simple features. The presence of large arrays in a program would decrease the storage overhead. With input/output the storage overhead are less than the same for programs without input/output. The compiler's optimization has mixed affects on storage overhead.

The overhead incurred seems highly dependent on both the FORTRAN compilers and the particular characteristics of the generated code. Anyway the use of preprocessors is not discouraged. They offer many advantages in coding over working with FORTRAN code directly. However, it is important for a user to recognize all of the costs of using the preprocessor, as well as all of the benefits received.

EXPERIENCE AND CONCLUSION

The preprocessor is a good software tool for enhancing the facilities in a programming language. It is quite obvious that many things can be done with a preprocessor once it is available. For example, the data types of FORTRAN can be improved by adding the CHARACTER type, modifications in EXTFOR to translate the operators =, <, >, |, >=, <=, ^=, and & as .EQ., .LT., .GT., .NOT., .OR., .GE., .LE., .NE., and .AND..

When designing a language or enhancing it, it is important to concentrate on the essential requirement of providing the user with the best language possible for a given effort. One must avoid throwing features - things which the user may trivially construct within the existing framework. For example, the REPEAT statement which is a loop with its test at the bottom. This statement encourages programs which fail at their boundaries. In the few cases where it is needed, it can be easily simulated with an infinite loop and a test and break at the bottom.

The main problem faced while using preprocessor is the fact that errors detected by the compiler are reported with respect to the preprocessor generated program, not the program written by the user. However, the readability of the program, easier maintenance and modification, and the reduction in programming time by structured programming are not disputable with difficulty faced because of bad error detection. While designing EXTFOR we concentrated on this aspect.

The PARSER is made better to recover from invalid syntax like ELSEs without IFs, UNTILs without REPEAT, etc. The situation where a matching END is not encountered to terminate a BEGIN, WHILE or DO loop, becomes embarrassing because the END which is also used as the last statement of a subprogram, is assumed to be associated with the loop or block and the next subprogram becomes the part of the first. This unrecovered error keeps on communicating; to avoid this EXTFOR provides the user to specify '*' in the first column of the line with END statement, used as the last statement of subprogram; seeing it, EXTFOR terminates the scope of current subprogram. If the context stack is not empty it is made so and an error is reported; processing of next subprogram starts as a fresh with no interaction with errors in previous subprogram.

The error detection capabilities of EXTFOR^{are} improved by tagging each output line of object code with the source line of extended feature, that created the code. The source lines of an extended feature are listed as comment followed by equivalent FORTRAN code generated by EXTFOR with '++' in columns 73 and 74 of each object line. This way user can distinguish between the regular FORTRAN features in his program and the EXTFOR generated code. The errors which come from the FORTRAN compiler in generated codes can be readily associated with their source of extended features. The errors reported by preprocessor are input with the source of EXTFOR appearing as comments in the listing.

The choice of selecting a programming language to write the preprocessor is opened to any language but,

one with better character handling facilities is the ideal choice; preferably PL/I or PASCAL could be selected to improve the efficiency of the preprocessor. Our choice for EXTFOR writing is the FORTRAN because we already have FORTRAN compiler to compile the EXTFOR generated code. The EXTFOR is portable to all computer installations where FORTRAN users want to avail the EXTFOR features.

EXTFOR is written using the modular approach. The program is divided into small subroutines and functions. Subroutines and functions rarely spread over more than two pages; most are much shorter. As a result the code is readable. It is easy to convince that EXTFOR module is probably correct because it is broken up into pieces that one can grasp one at a time. It is tried to make the routines easy to modify by keeping the routines unknown about information that not needed by them. The informations, commonly used by more than two routines are defined in COMMON blocks referred as EXTFOR work area. The less frequently used information is communicated between two routines through parameters.

It is well known that no program comes out to a perfect work of art on its first draft, regardless of the technique one uses to write it. Most of the routines were written, modified and tested several times, yet we still would not claim that any one is flawless. The EXTFOR module is tested thoroughly particularly at boundaries.

EXTFOR SOURCE LISTING

```

C*****
C
C           A FORTRAN PREPROCESSOR
C-----
C
C   THIS PROGRAM TRANSLATES 'FORTRAN' WITH EXTENDED FEATURES,
C   INTO 'FORTRAN', THE LANGUAGE IS SPECIFIED BY FOLLOWING GRAMMAR
C
C   PROGRAM : STATEMENT
C             PROGRAM STATEMENT .
C   STATEMENT I IF (CONDITION) THEN STATEMENT
C             I IF (CONDITION) THEN STATEMENT
C             ELSE STATEMENT
C             I BEGIN
C             STATEMENT
C             END
C             I DO LIMITS
C             STATEMENT
C             END
C             I WHILE (CONDITION) STATEMENT
C             END
C             I FOR (INITIALIZE;(CONDITION);REINITIALIZE)
C             STATEMENT
C             I REPEAT STATEMENT
C             UNTIL < (CONDITION) >
C             I LABEL STATEMENT
C             I BREAK <LEVEL>
C             I NEXT <LEVEL>
C             I <PROGRAM>
C*****

```

*** PARSER ***

```

C           I OTHER
C
C           THE VERTICAL BAR I INDICATES A CHOICE OF ALTERNATIVES
C           < > INDICATES TO BE OPTIONAL, 'OTHER' IS ANYTHING THAT
C           WAS NOT RECOGNIZED AS ANY OF THE PRECEDING TYPES.
C           TO TRANSLATE THE EXTENDED CONTROL STRUCTURE, AT THE
C           BEGINNING OF EACH STATEMENT THIS PROGRAM CALLS GETTOK TO
C           GET THE FIRST TOKEN OF THE STATEMENT WHICH DETERMINES
C           THE STATEMENT TYPE. AT A TIME 80 CHARACTERS OF A LINE
C           ARE READ INTO IN-BUFFER WITH A SOURCE POINTER LASTC USED
C           TO SCAN THE SOURCE LINE. WHEN THE STATEMENT IS CLASSIFIED,
C           THIS ROUTINE CALLS THE APPROPRIATE CODE GENERATION ROUTINE
C           AND CORRESPONDING TYPE WITH A UNIQUE GENERATED LABEL ARE
C           PUT ON THE CONTEXT STACK IN CASE OF 'IF', 'ELSE', 'BEGIN',
C           'DO', 'WHILE', 'REPEAT', AND 'FOR'. WHEN THE END OF STATEMENT
C           OR MATCHING 'END' FOR 'DO', 'BEGIN', AND 'WHILE' IS ENCOUN-
C           TERED APPROPRIATE CODES ARE GENERATED ALSO, ENTRY FROM
C           CONTEXT STACK IS REMOVED. WHEN A LINE IS PROCESSED, NEXT
C           LINE IS READ FOR THE CONTINUATION OF THE STATEMENT OR
C           THE BEGINNING OF NEXT STATEMENT.
C           THE EXTENDED CONTROL STATEMENT ARE PRINT AS SUCH STARTING
C           WITH 'C+' AND GENERATE CODES ENDING WITH '++' IN 72-73
C           COLUMN, APPROPRIATE ERROR MESSAGES ARE PRINTED WHEN AN
C           ERROR IS DETECTED, THE CONDITION PART OF SOME OF THE
C           STATEMENT IS LEFT TO FORTRAN COMPILER TO CHECK IT.
C           THE FIRST GENERATED LABEL IS 900000.
C
C
C
C
C*****

```

*** PARSER ***

```

C*****
C
C      PARSER
C
C.....
C      INTEGER ERRNOB,J,LAB,LABS,LABST,LASTC,LEVEL          PARS0001
C      INTEGER NLASTC,NFOR,NERR,OUTP,PTR,SP                PARS0002
C      INTEGER*2 ZERO/0/                                   PARS0003
C      INTEGER*2 STAR/'*'/,BLANK/' '/,ZERO/'0'/,EOS/-2/,EOF/-1/ PARS0004
C      INTEGER*2 LEXIF/-30/,LEXDO/-31/,LEXWHL/-32/,LEXBEG/-33/ PARS0005
C      INTEGER*2 LEXFOR/-34/,LEXREP/-35/,LEXELS/-36/,LEXEND/-37/ PARS0006
C      INTEGER*2 LEXNXT/-38/,LEXBRK/-39/,LEXUNT/-40/,LEXOTR/-41/ PARS0007
C      INTEGER*2 LEX,LEXS,LEXST,LEXTKN,NSTTOK,RRNFLG,TKNSWH,IFLAG,I PARS0008
C      INTEGER*2 BUF(80),CLVEL(2),INLAB(5),TOKEN(80)      PARS0009
C
C      COMMON BLOCKS
C      COMMON /COMPTR/ PTR                                  PARS0010
C      COMMON /ERRNO / ERRNOB                              PARS0011
C      COMMON /EOSRC / TEOSRC                              PARS0012
C      COMMON /INPTR / LASTC                              PARS0013
C      COMMON /INBLK / BUF                                 PARS0014
C      COMMON /INLAB / INLAB                              PARS0015
C      COMMON /NSTFLG/ NSTTOK                             PARS0016
C      COMMON /NMBFOR/ NFOR                                PARS0017
C      COMMON /NDERR / NERR                                PARS0018
C      COMMON /NLST / NLASTC                              PARS0019
C      COMMON /OUTPTR/ OUTP                                PARS0020
C      COMMON /RNFLG / RRNFLG                             PARS0021
C      COMMON /RTFLAB/ LAB                                 PARS0022
C      COMMON /STKPTR/ SP                                  PARS0023

```

*** PARSER ***

	COMMON /SYSIN / TSYSTEM	PARS0024
	COMMON /SYSOUT/ TSYOUT	PARS0025
	COMMON /TKN / TOKEN	PARS0026

C	NERR IS NUMBER OF DIAGNOSTICS GENERATED.	*****
C	ERRNOB IS NUMBER OF ERROR IN A STATEMENT.	*****
C	IEDSRC IS END OF SOURCE FLAG.	*****
C	SP IS STACK POINTER.	*****
C	LASTC IS SOURCE BUFFER POINTER.	*****
C	PTR IS CONDITION BUFFER POINTER.	*****
C	NFOR IS NO. OF NESTED FOR.	*****
C	LAB IS GENERATED LABEL.	*****
C	RRNFLG IS RERUN FLAG.	*****
C	TKNSWH IS TOKEN SWITCH.	*****
C	NSTTOK IS NEW STATEMENT TOKEN.	*****

	TSYSIN=1	PARS0027
	TSYOUT=3	PARS0028
	NERR=0	PARS0029
	ERRNOB=0	PARS0030
	IEDSRC=0	PARS0031
	SP=0	PARS0032
	LASTC=72	PARS0033
	OUTP=6	PARS0034
	PTR=14	PARS0035
	NFOR=0	PARS0036
	LAB=90000	PARS0037
	RRNFLG=0	PARS0038
	TKNSWH=0	PARS0039
	GO TO 50	PARS0040

*** PARSER ***

5	NLASTC=LASTC+1	PARS0041
C	GET THE NEXT TOKEN AND CHECK FOR A NEW STATEMENT	*****
	CALL GETTOK	PARS0042
	IF (TOKEN(1).EQ.EOF) GO TO 3000	PARS0043
	IF (NSTTOK.EQ.0) GO TO 158	PARS0044
	LASTC=6	PARS0045
	GO TO 2001	PARS0046
10	RRMFLG=0	PARS0047
20	IF (LASTC.EQ.6) GO TO 50	PARS0048
40	LASTC=72	PARS0049
50	NSTTOK=0	PARS0050
	NLASTC=1	PARS0051
100	CALL GETTOK	PARS0052
	NLASTC=1	PARS0053
	IF (TOKEN(1).EQ.EOF) GO TO 3000	PARS0054
C	CHECK FOR THE CONTINUATION OF THE STATEMENT AND	*****
C	PRINT IF SO.	*****
	IF (NSTTOK.NE.0) GO TO 150	PARS0055
	CALL PUTLIN(BUF)	PARS0056
	GO TO 40	PARS0057
C	STATEMENT LABEL.	*****
150	DO 155 I=1,5	PARS0058
	INLAR(I)=BUF(I)	PARS0059
155	CONTINUE	PARS0060
158	BUF(6)=BLANK	PARS0061
	CALL GFTSTK (LEXS,LARS)	PARS0062
160	LEXTKN=LEX(TOKEN)	PARS0063
170	IF (TOKEN(1).EQ.EOF) GO TO 3000	PARS0064
	IF (LEXTKN.EQ.LEXOTR) GO TO 1200	PARS0065
	IF (LEXTKN.EQ.LEXDO) GO TO 300	PARS0066

*** PARSER ***

C	IF (LEXTKN.EQ.LEXEND) GO TO 900	PARS0067
	PRINT OUT THE STATEMENT AS A COMMENT.	*****
	CALL PUTSMT(NLASTC)	PARS0068
	I=- (LEXTKN)-29	PARS0069
C	BRANCH TO PROCESS THE STATEMENT DEPENDING	*****
C	UPON THE THE FIRST TOKEN OF THE STATEMENT.	*****
	GO TO(200,300,400,500,600,700,800,900,1000,1000,1100,1200),I	PARS0070
C		*****
C*	IF STATEMENT	*****
200	CALL IFCODE	PARS0071
	GO TO 2000	PARS0072
C		*****
C*	DO STATEMENT	*****
300	CALL DDCODE (I,NLASTC)	PARS0073
	IF ((I.EQ.7ER0).OR.((LEXS.NE.LEXIF)	PARS0074
	* .AND.(LEXS.NE.LEXELS).AND.	PARS0075
	* (LFXS.NE.LEXFOR))) GO TO 10	PARS0076
C	ERROR--'FORTRAN DO STATEMENT ILLEGAL HERE'	*****
	CALL ERROR (13)	PARS0077
	GO TO 2001	PARS0078
C		*****
C*	WHILE STATEMENT	*****
400	CONTINUE	PARS0079
	CALL WHILE	PARS0080
	RRNFLG=1	PARS0081
	GO TO 5	PARS0082
C		*****
C*	REGIN STATEMENT	*****
500	CONTINUE	PARS0083
	CALL CHKBLK	PARS0084

*** PARSER ***

	CALL BGENEN	PARS0085
	GO TO 10	PARS0086
C		*****
C*	FOR STATEMENT	*****
600	CONTINUE	PARS0087
	IF(NFOR.LT.4) GO TO 650	PARS0088
C	' TOO MANY NESTED 'FOR' STATEMENTS'	*****
	CALL ERROR (5)	PARS0089
	GO TO 2001	PARS0090
650	CONTINUE	PARS0091
	CALL FORGEN	PARS0092
	GO TO 2000	PARS0093
C		*****
C*	REPEAT STATEMENT	*****
700	CONTINUE	PARS0094
	CALL REPGEN	PARS0095
	RRNFLG=1	PARS0096
	GO TO 5	PARS0097
C		*****
C*	ELSE STATEMENT	*****
800	CONTINUE	PARS0098
C	ERROR= 'UNM	
	CALL ERROR (9)	PARS0099
	GO TO 2001	PARS0100
C		*****
C*	END STATEMENT	*****
900	CONTINUE	PARS0101
	IF (RRNFLG.NE.7ER0) GO TO 910	PARS0102
	IF (BUF(1).NE.STAR) GO TO 950	PARS0103
901	WRITE (ISYOUT,902) (BUF(I),I=2,80)	PARS0104

*** PARSER ***

902	FORMAT (1X,79A1)	PARS0105
	IF(SP.EQ.0) GO TO 10	PARS0106
C	ERROR--'UNEXPECTED END OF SUBPROGRAM'	*****
	CALL ERROR (10)	PARS0107
	SP=0	PARS0108
	GO TO 10	PARS0109
905	IF ((LEXS.NE.LEXIF).AND.(LEXS.NE.LEXREP)	PARS0110
1	.AND.(LEXS.NE.LEXFOR)) GO TO 950	PARS0111
910	CALL PUTSMT (NLASTC)	PARS0112
C	ERROR--'CONDITIONAL END STATEMENT'	*****
915	CALL ERROR (8)	PARS0113
	GO TO 2001	PARS0114
950	IF (SP.NE.0) GO TO 952	PARS0115
	CALL PUTLIN(BUF)	PARS0116
	GO TO 40	PARS0117
952	CALL PUTSMT(NLASTC)	PARS0118
	CALL CHKBLK	PARS0119
	IF ((LEXS.EQ.LEXFOR).OR.	PARS0120
1	(LEXS.EQ.LEXREP)) GO TO 915	PARS0121
C	REMOVE ENTRY FROM STACK	*****
	SP=SP-1	PARS0122
	I=- (LEXS)-30	PARS0123
	GO TO (960,970,965),I	PARS0124
C		*****
C*	END OF BEGIN/DO UNIT	*****
C	GENERATE 'LABS+1 CONTINUE'	*****
C	'LABS CONTINUE'	*****
960	CALL OUTCON (LABS+1)	PARS0125
965	CALL OUTCON(LABS+2)	PARS0126
	GO TO 2001	PARS0127

*** PARSER ***

C		*****
C*	END OF WHILE UNIT	*****
C	GENERATE 'LABS GO TO LABS+1'	*****
C	'LABS+2 CONTINUE'	*****
970	CALL OUTGON(LABS,LABS+1)	PARS0128
	CALL OUTCON(LABS+2)	PARS0129
	GO TO 2001	PARS0130
995	CONTINUE	PARS0131
C		*****
C*	BREAK/NEXT STATEMENT	*****
C	CHECK THE LEVEL OF THE LOOP TO BE BROKEN	*****
C	IF LEVEL IS NOT SPECIFIED '1' IS DEFAULT	*****
1000	CONTINUE	PARS0132
	CALL GFTTOK	PARS0133
	IF(NSITOK.NE.1) GO TO 1008	PARS0134
	TKNSWH=1	PARS0135
1005	LEVEL=1	PARS0136
	GO TO 1030	PARS0137
1008	CONTINUE	PARS0138
	DO 1010 I=1,80	PARS0139
	IF(TOKEN(I).EQ.EOS) GO TO 1015	PARS0140
1010	CONTINUE	PARS0141
1015	IF(I.EQ.1) GO TO 1005	PARS0142
	CLVEL(1)=ZERO	PARS0143
	IF(I.EQ.2) GO TO 1020	PARS0144
	CLVEL(1)=TOKEN(1)	PARS0145
	CLVEL(2)=TOKEN(2)	PARS0146
	GO TO 1023	PARS0147
1020	CLVEL(2)=TOKEN(1)	PARS0148
1023	CALL CHKBLK	PARS0149

*** PARSER ***

1025	LEVEL=10*((CLVFL(1)-ZERO)/256)	PARS0150
*	+((CLVFL(2)-ZERO)/256)	PARS0151
1030	IF(LEVEL.EQ.0) GO TO 2001	PARS0152
	IF(LEVEL.LT.0) CALL ERROR (12)	PARS0153
C	ERROR-- 'INVALID LEVEL SPECIFIED'	*****
	J=SP	PARS0154
	LEXST=LEXS	PARS0155
	LARST=LABS	PARS0156
	IFLAG=ZERO	PARS0157
	DO 1050 I=1,LEVEL	PARS0158
	IF ((LEXS.LE. LEXDO).AND.(LEXS.GE.LEXFOR))	PARS0159
1	GO TO 1035	PARS0160
	I=I-1	PARS0161
	IF(SP.GT.0) GO TO 1040	PARS0162
	IF (IFLAG.NE.ZERO) GO TO 1060	PARS0163
	CALL ERROR (14)	PARS0164
C	ERROR-- 'BREAK/NEXT ILLFGAL HERE'	*****
	SP=J	PARS0165
	GO TO 2001	PARS0166
1035	IFLAG=1	PARS0167
	LEXST=LEXS	PARS0168
	LARST=LABS	PARS0169
1040	SP=SP-1	PARS0170
	CALL GFTSTK(LEXS,LABS)	PARS0171
1050	CONTINUE	PARS0172
1060	CALL BRKNXT(LEXTKN,LEXST,LABST)	PARS0173
	SP=J	PARS0174
	GO TO 2001	PARS0175
C		*****
C*	UNTIL UNIT	*****

*** PARSER ***

1100	CONTINUE	PARS0176
	IF (RRNFLG.EQ.ZERO) GO TO 1120	PARS0177
C	ERROR--'CONDITIONAL UNTIL'	*****
	CALL ERROR (8)	PARS0178
	GO TO 2001	PARS0179
1120	IF (LEYS.EQ.LEXRFP) GO TO 1130	PARS0180
C	ERROR--'UNMATCHING UNTIL'	*****
	CALL ERROR (9)	PARS0181
	GO TO 40	PARS0182
1130	CONTINUE	PARS0183
	CALL UNGEN (LABS)	PARS0184
	SP=SP-1	PARS0185
	GO TO 2001	PARS0186
C		*****
C*	FORTTRAN STATEMENTS	*****
C	IF THE STATEMENT IS NOT RECOGNIZED AS AN EXTENDED	*****
C	STATEMENT, THE TRANSLATER ASSUMES IT MUST BE A	*****
C	FORTTRAN STATEMENT AND PASSES IT UNALTERED.	*****
1200	CALL OTHERS(NLASTC)	PARS0187
	GO TO 2001	PARS0188
C	NON-ZERO RRRFLG INDICATES THE STATEMENT IS IF, WHILE,	*****
C	FOR STATEMENTS. IN CASE THERE IS AN ENTRY FOR 'IF' ON	*****
C	THE TOP OF STACK, NEXT STATEMENT IS CHECKED WHETHER	*****
C	IT IS 'ELSE', OTHERWISE CODE IS GENERATED FOR THE END	*****
C	OF IF STATEMENT, FOR ELSE AN ENTRY IS MADE ON STACK	*****
C	AND CODES ARE GENERATED FOR BEGIN OF ELSE-UNIT.	*****
2000	IF (RRNFLG.NE.ZERO) GO TO 5	PARS0189
	ERRNOB=0	PARS0190
2001	CALL GETSTK (LEXS,LABS)	PARS0191
	RRNFLG=ZERO	PARS0192

*** PARSER ***

2002	IF (LEYS.NE.LEXIF) GO TO 2020	PARS0193
	IF (TKNSWH.EQ.1) GO TO 2006	PARS0194
	IF (LASTC.EQ.6) GO TO 2005	PARS0195
2004	LASTC=72	PARS0196
2005	NSTTOK=ZERO	PARS0197
	CALL GETTOK	PARS0198
	TKNSWH=1	PARS0199
2006	NLASTC=1	PARS0200
	IF (NSTTOK.NE.ZERO) GO TO 2008	PARS0201
	CALL PUTLIN(BUF)	PARS0202
	GO TO 2004	PARS0203
2008	LEXTKN=LEX(TOKEN)	PARS0204
	DO 2009 I=1,5	PARS0205
	INLAR(I)=BUF(I)	PARS0206
2009	CONTINUE	PARS0207
	IF(LEXTKN.NE.LEXELS) GO TO 2015	PARS0208
C		*****
C*	ELSE STATEMENT	*****
2010	CALL PUTSMT (1)	PARS0209
	DO 2011 I=1,5	PARS0210
	IF(BUF(I).NE.BLANK) GO TO 2012	PARS0211
2011	CONTINUE	PARS0212
	GO TO 2013	PARS0213
2012	CALL ERROR (1)	PARS0214
C	LABEL ILLEGAL HERE-WARNING	*****
2013	CONTINUE	PARS0215
C	STACK 'ELSE'	*****
	CALL STACK(LEXFLS,IARS)	PARS0216
	K=LARS+2	PARS0217
	L=LARS+1	PARS0218

*** PARSEP ***

C	GENERATE 'LABS GO TO LABS+2'	*****
C	'LABS+1 CONTINUE'	*****
	CALL OUTCON (LABS,K)	PARS0219
	CALL OUTCON (L)	PARS0220
	RRNFLG=1	PARS0221
	TKNSWH=ZERO	PARS0222
	GO TO 5	PARS0223
C*	CODE FOR END OF IF UNIT	*****
C	'LABS+1 CONTINUE'	*****
2015	CALL OUTCON (LABS+1)	PARS0224
C	POP OFF 'IF' FROM CONTEXT STACK	*****
	SP=SP-1	PARS0225
	CALL GETSTK(LEXS,LABS)	PARS0226
	IF(LFXS.EQ.LEXIF) GO TO 2015	PARS0227
2020	IF(LFXS.NE.LEXELS) GO TO 2100	PARS0228
C*	CODE FOR END OF ELSE UNIT	*****
C	'LABS+2 CONTINUE'	*****
	CALL OUTCON (LABS+2)	PARS0229
C	POP OFF 'ELSE' AND 'IF' FROM CONTEXT STACK	*****
	SP=SP-2	PARS0230
	GO TO 2001	PARS0231
C	THE TOP MOST ENTRY AT STACK IS CHECKED,	*****
C	IF 'FOR' THEN CODE FOR END OF FOR UNIT	*****
C	ARE GENERATED AND ENTRY IS REMOVED.	*****
2100	IF (LEYS.NE.LEXFOR) GO TO 2200	PARS0232
	CALL FOREND (LABS)	PARS0233
	SP=SP-1	PARS0234
	GO TO 2001	PARS0235
2200	CONTINUE	PARS0236
	IF(TKNSWH.NE.1) GO TO 10	PARS0237

*** PARSER ***

	TKNSWH=ZERO	PARS0238
	GO TO 170	PARS0239
C	WHEN END OF SOURCE IS ENCOUNTERED THE STACK IS CHECKED,	*****
C	WHICH MUST BE EMPTY FOR A COMPLETE PROGRAM.	*****
3000	CONTINUE	PARS0240
	IF (SP.NE.0) CALL ERROR (10)	PARS0241
C	ERROR--'UNEXPECTED END OF SUBPROGRAM'	*****
	WRITE (6,3010) NERR	PARS0242
3010	FORMAT (1H1,1X/1H ,10 X,'** STATISTICS AT PREPROCESSOR LEVEL	*PARS0243
	1*'/1H ,1X/1H ,10X,I3,' DIAGNOSTICS GENERATED')	PARS0244
	STOP	PARS0245
	END	PARS0246

*** GFTTOK ***

```

SUBROUTINE GETTOK
C*****
C
C   THIS ROUTINE GETS TOKEN AND PUT IN COMMON BLOCK TKN,
C   FLAG NSTTOK IS SET TO 1,IF TOKEN IS THE FIRST TOKEN OF THE
C   NEW STATEMENT, TO BE USED TO CLASSFY THE STATEMENT.
C   THE END OF TOKEN IS INDICATED BY EOS.
C*****
C
C   INTEGER   LASTC,LENTH1,TKLNTH
C   INTEGER*2 C,INBUF(80),LEXSTR(80),NSTTOK,STMCON,TYPE
C   INTEGER*2 BLANK/' '/,DLIMIT/-10/,EOF/-1/,EOS/-2/,ZERO/'0'/
C
C   COMMON BLOCKS
C   COMMON /INBLK / INBUF
C   COMMON /INPTR / LASTC
C   COMMON /NSTFLG/ NSTTOK
C   COMMON /PLENTH/ LENH1
C   COMMON /TKN   / LEXSTR
C   COMMON /TKSIZE/ TKLNTH
C
C
C   TKLNTH IS THE NO. OF CHARS IN THE TOKEN.
C   IF TOKEN IS CONTINUED ON THE NEXT LINE THEN LENH1
C   INDICATES THE NO. OF CHARS OF TOKEN ON THE RIRST LINE.
C
C   NSTTOK=0
C   TKLNTH=0
C   I=1
C   LENH1=0
10  CALL GETC(C)
C   IF (LASTC.NE.7) GO TO 25

```

```

GETT0000
*****
*****
*****
*****
*****
*****
GETT0001
GETT0002
GETT0003
*****
GETT0004
GETT0005
GETT0006
GETT0007
GETT0008
GETT0009
*****
*****
*****
*****
GETT0010
GETT0011
GETT0012
GETT0013
GETT0014
GETT0015

```

*** GETTOK ***

C	STMCON IS STATEMENT CONTINUATION CHARACTER IN THE 6TH COL.	*****
	STMCON=INBUF(6)	GETT0016
	IF ((STMCON.EQ.BLANK).OR.(STMCON.EQ.ZERO)) NSTTOK=1	GETT0017
20	IF(C.EQ.EOF) GO TO 60	GETT0018
25	IF(C.NE.BLANK) GO TO 30	GETT0019
	CALL GETC(C)	GETT0020
	IF (LASTC.NE.7) GO TO 25	GETT0021
	STMCON=INBUF(6)	GETT0022
	IF ((STMCON.EQ.BLANK).OR.(STMCON.EQ.ZERO)) GO TO 200	GETT0023
	GO TO 20	GETT0024
C	THE SIGNIFICANT DELIMITER FOR PREPROCESSOR ARE ')','(','!','	*****
C	AND BLANK	*****
30	IF(TYPE(C).EQ.DLIMIT) GO TO 60	GETT0025
	LEXSTR(I)=C	GETT0026
	I=I+1	GETT0027
40	CALL GETC(C)	GETT0028
	IF(LASTC.NE.7) GO TO 30	GETT0029
	STMCON=INBUF(6)	GETT0030
50	IF((STMCON.EQ.BLANK).OR.(STMCON.EQ.ZERO)) GO TO 200	GETT0031
	LENTH1=I-1	GETT0032
	GO TO 30	GETT0033
60	IF(I.GT.1) GO TO 100	GETT0034
C	THE TOKEN IS A NON-BLANK DELIMITER	*****
70	LEXSTR(1)=C	GETT0035
	LEXSTR(2)=EOS	GETT0036
	RETURN	GETT0037
100	LASTC=LASTC-1	GETT0038
	LEXSTR(I)=EOS	GETT0039
	TKLNTH=I-1	GETT0040
	RETURN	GETT0041

*** GETTOK ***

200 LASTC=6
250 LEXSTR (I)=EOS
TKLNTH=I-1
RETURN
END

GETT0042
GETT0043
GETT0044
GETT0045
GETT0046

*** GETC ***

5	READ (TSYSIN,10,END=200) (BUF(I),I=1,MAXCRD)	GETC0015
10	FORMAT(80A1)	GETC0016
C	COMMENT STATEMENT IS PRINTED AS IT IS.	*****
	IF (BUF(1).NE.COMT) GO TO 50	GETC0017
	WRITE (TSYOUT,10) (BUF(I),I=1,MAXCRD)	GETC0018
C	GO TO READ THE NEXT LINE	*****
	GOTO 5	GETC0019
50	LASTC=7	GETC0020
C	STORE THE STATEMENT LABEL	*****
	DO 60 I=1,5	GETC0021
	LABEL(I)=BUF(I)	GETC0022
60	CONTINUE	GETC0023
100	C=BUF(LASTC)	GETC0024
	RETURN	GETC0025
200	IEOSRC=1	GETC0026
C	SET END OF SOURCE FLAG	*****
250	C=EOF	GETC0027
	LASTC=7	GETC0028
	BUF(6)=BLANK	GETC0029
	BUF(7)=BLANK	GETC0030
	RETURN	GETC0031
	END	GETC0032

*** CHKBLK ***

```

SUBROUTINE CHKBLK                                     CHKB0000
C*****
C
C      THIS ROUTINE CHECKS FOR A NON-BLANK CHARACTER IN THE REMAINING
C      PART OF THE BUFFER AND ISSUES AN ERROR IF ENCOUNTERED.
C*****
C*****
C
C      INTEGER*2 BUF(80),BLANK/' '/
C      INTEGER LASTC
C      COMMON/INBLK/BUF
C      COMMON/INPTR/LASTC
C
C      J=LASTC+1
C      IF(J.GE.73) RETURN
C      DO 10 I=J,72
C      IF(BUF(I).NE.BLANK) GO TO 100
C      CONTINUE
C      RETURN
C      ERROR--'SYNTAX ERROR'
100    CALL ERROR (15)
C      RETURN
C      END
CHKB0001
CHKB0002
CHKB0003
CHKB0004
*****
CHKB0005
CHKB0006
CHKB0007
CHKB0008
CHKB0009
CHKB0010
*****
CHKB0011
CHKB0012
CHKB0013

```


*** STACK ***

	SUBROUTINE STACK (SLEX,SLAB)	STAC0000
C	*****	*****
C		*****
C	THIS ROUTINE POPS IN THE CONTEXT STACK WITH SLEX=SLAB	*****
C	MAKING ONE ENTRY.	*****
C		*****
C	*****	*****
C		*****
	DIMENSION STKLEX(100),STKLAB(100)	STAC0001
	INTEGER STKLAB,SLAB,SP	STAC0002
	INTEGER*2 STKLEX,SLEX	STAC0003
	COMMON/STKPTR/SP	STAC0004
	COMMON/STKBLK/STKLEX,STKLAB	STAC0005
C		*****
C	INCREMENT THE STACK POINTER.	*****
	SP=SP+1	STAC0006
	IF(SP.GT.100) GO TO 10	STAC0007
	STKLEX(SP)=SLEX	STAC0008
	STKLAB(SP)=SLAB	STAC0009
	RETURN	STAC0010
10	CONTINUE	STAC0011
	RETURN	STAC0012
	END	STAC0013

*** ERROR ***

```

SUBROUTINE ERROR (FRCODE)
C*****
C
C PRINTS ERROR MESSAGE
C*****
C*****
C
INTEGER FRCODE, LASTC, NERR, ERRNOB, ERCD?
INTEGER*2 ISNTX(66)/66* ' ', DOLLER/'$'/, BLANK/' '/
INTEGER*2 ERCD, ZFRO/'0'/, ONE/'1'/
INTEGER ERRORM(8,14) /'WARN', 'ING-', '-LAB', 'EL I', 'LLEG', 'AL H',
* 'ERE ', '
* ' ('-'', '-MIS', 'STNG', '
* 'UNMA', 'TCHI', 'NG ', ' (''S
* 'UNRE', 'COGN', 'IZAR', 'LE S', 'TATE', 'MENT', '
* 'TOO ', 'MANY', ' NES', 'TED ', ' 'FOR', ' ' ST', 'ATEM', 'ENTS',
* 'SYNT', 'AX E', 'RROR', ' IN ', ' 'FOR', ' ' ST', 'ATEM', 'ENT ',
* 'TOO ', 'LARG', 'E A ', 'COND', 'ITIO', 'N/ST', 'ATEM', 'ENT ',
* 'COND', 'ITIO', 'NAL ', 'END/', 'ELSE', '/UNT', 'IL ',
* 'UNMA', 'TCHI', 'NG E', 'LSE/', 'UNTI', 'L
* 'UNEX', 'PECT', 'ED E', 'ND O', 'F SU', 'BPRO', 'GRAM',
* 'NO L', 'OOP ', 'EXIS', 'TS
* 'ILLE', 'GAL ', 'LEVE', 'L SP', 'ECIF', 'IED
* 'FORT', 'RAN ', 'DO S', 'TMT ', 'ILLE', 'GAL ', 'HERE',
* 'BREA', 'K/NE', 'XT I', 'LLEG', 'AL H', 'ERE ', ' '
C
COMMON BLOCKS
COMMON /ERRNO / ERRNOB
COMMON /INPTR / LASTC
COMMON /NOERR / NERR
COMMON /SYSOUT / ISYOUT

```

ERR00000

ERR00001
ERR00002
ERR00003
ERR00004
ERR00005
ERR00006
ERR00007
ERR00008
ERR00009
ERR00010
ERR00011
ERR00012
ERR00013
ERR00014
ERR00015
ERR00016
ERR00017
ERR00018

ERR00019
ERR00020
ERR00021
ERR00022

*** ERROR ***

C		*****
C	NERR IS TOTAL NUMBER OF ERROR S IN A PROGRAM	*****
	NERR=NERR+1	ERR00023
C	ERRNOB IS SERIAL NO. OF ERROR IN A STATEMENT	*****
	ERRNOB=ERRNOB+1	ERR00024
	IF(ERCODE,GE.10) GO TO 20	ERR00025
	ERCODE2=ERCODE	ERR00026
	ERCD=ZERO	ERR00027
	GO TO 30	ERR00028
20	ERCD=ONE	ERR00029
	ERCODE2=ERCODE-10	ERR00030
	IF (ERCODE.GT.13) GO TO 100	ERR00031
30	WRITE (ISYOUT,50) ERRNOB,ERCD,ERCODE2,(ERRORM(I,ERCODE),I=1,8)	ERR00032
50	FORMAT ('C***0',I1,') JFP0',A1,I1,'I ',8A4,29X)	ERR00033
C	THE FORMAT OF ERRORS MESSAGE IS	*****
C	C***01) JFP002I TEXT OF MESSAGE	*****
	RETURN	ERR00034
100	T=LASTC-6	ERR00035
	ISNTX(I)=DOLLER	ERR00036
	WRITE (ISYOUT,150) ISNTX,ERRNOB,ERCD,ERCODE2	ERR00037
150	FORMAT ('C ',6A1,8X/	ERR00038
1	'C***0',I1,') JFP0',A1,I1,'I SYNTAX',55X)	ERR00039
	ISNTX(I)=BLANK	ERR00040
	RETURN	ERR00041
	END	ERR00042

*** BNGEN ***

```

SUBROUTINE BNGEN
C*****
C
C   THIS ROUTINE GENERATE CODE FOR BEGIN STATEMENT,
C   NO CODE IS GENERATED IF STATEMENT IS UNLABELED,
C   AN ENTRY IS MADE ON CONTEXT STACK
C
C*****
C
C   INTEGER LAB,LABGEN
C   INTEGER*2 LABEL(5),LEXRGN/-33/,BLANK/' '/,I
C   COMMON/INLRL/LABEL
C   COMMON /SYSOUT/ISYOUT
C
C   LAB=LARGEN (1)
C   1 GENERATED LABEL IS RESERVED
C   DO 10 I=1,5
C   IF (LABEL(I).NE.BLANK) GO TO 20
10  CONTINUE
C   GO TO 50
C   FOR LABELED 'BEGIN' CODE 'LABEL CONTINUE' IS GENERATED
20  WRITE (ISYOUT,30) LABEL
30  FORMAT(5A1,' CONTINUE',58X,'++',6X)
C   POP IN LEXRGN AND GENERATED LABEL
50  CALL STACK(LEXRGN,LAB-2)
C   RETURN
C   END
BNGG0000
*****
*****
*****
*****
*****
*****
BNGG0001
BNGG0002
BNGG0003
BNGG0004
*****
BNGG0005
*****
BNGG0006
BNGG0007
BNGG0008
BNGG0009
*****
BNGG0010
BNGG0011
*****
BNGG0012
BNGG0013
BNGG0014

```

*** BALPAR ***

```

SURROUTINE BALPAR BALP0000
C*****
C
C BALPAR COLLECTS AND PUT THE CONDITION PART IN CONBUF,
C WHICH IS A STRING ENCLOSED IN BALANCED PARENTHESES.
C*****
C*****
C
C INTEGER*2 ZERO/1H0/,BLANK/1H /,EOF/-1/,LBRK/1H(/,RBRK/1H)/
C INTEGER*2 BUF(80),C,CONT,NL BALP0001
C INTEGER LASTC BALP0002
C COMMON/INBLK/ BUF BALP0003
C COMMON/INPTR/LASTC BALP0004
C BALP0005
C*****
C NL IS NO. OF '(' ENCOUNTERED
C NL=1
C GO TO 17 BALP0006
C
C 10 NL=NL+1 BALP0007
C 15 CALL PUTCND(C) BALP0008
C 17 CALL GETC(C) BALP0009
C IF(LASTC.NE.7) GO TO 20 BALP0010
C CONT=BUF(6) BALP0011
C CONTINUATION OF CONDITION ON NEXT LINE
C IF ((CONT.EQ.ZERO).OR.(CONT.EQ.BLANK)) GO TO 50 BALP0012
C 20 IF(C.EQ.LBRK) GO TO 10 BALP0013
C IF(C.NE.RBRK) GO TO 15 BALP0014
C NL=NL-1 BALP0015
C IF(NL.EQ.0) GOTO 100 BALP0016
C GO TO 15 BALP0017
C 50 LASTC=6 BALP0018
C BALP0019

```

*** BALPAR ***

	IF(NL.EQ.0) RETURN	BALP0020
C	ERROR--'UNMATCHING ')''	*****
	CALL ERROR (3)	BALP0021
	RETURN	BALP0022
100	CALL PUTEND(C)	BALP0023
	RETURN	BALP0024
	END	BALP0025

*** IFCODE ***

```

SUBROUTINE IFCODE
C*****
C
C      PROCESS IF STATEMENT, FORTRAN IF STATEMENTS ARE BYPASSED
C*****
C*****
C      INTEGER LAB,LARGEN,KEY
C      INTEGER*2 LEXIF/-30/,RRNFLG
C      COMMON/RRNFLG/RRNFLG
C
C      RRNFLG=0
C      3 GENERATED LABELS ARE RESERVED FOR IF STATEMENT
C      LAB=LARGEN(3)
C      CALL IFGO(1,LAB+1,KEY)
C      KEY VALUES 1 WHEN STATEMENT IS FORTRAN IF STATEMENT
C      TF(KEY.EQ.1) RETURN
C      AN ENTRY FOR 'IF' ON CONTEXT STACK.
C      CALL STACK (LEXIF,LAB)
C      RRNFLG=1
C      RETURN
C      END
IFC00000
*****
*****
*****
*****
IFC00001
IFC00002
IFC00003
*****
IFC00004
*****
IFC00005
IFC00006
*****
IFC00007
*****
IFC00008
IFC00009
IFC00010
IFC00011

```

*** IFGO ***

```

SUBROUTINE IFGO(CNTRL,LAR,TYPE)                                IFGO0000
C*****
C
C      THIS ROUTINE IS CALLED TO PROCESS THE CODITION PART OF   IFGO0001
C      IF, WHILE, UNTIL, AND FOR STATEMENT. IF CALLED WITH      IFGO0002
C      CNTRL=1 IT CHECKS FOR FORTRAN IF STATEMENTS,IF THE STMT  IFGO0003
C      IS NOT FORTRAN STMT KEY IS SET TO 0 OTHERWISE 1.        IFGO0004
C
C*****
C      INTEGER    CNTRL,LAR,LASTC,TYPE                          IFGO0005
C      INTEGER*2  THEN(5) /'I','H','E','N',-2/,LBRK/'(' /,EOS/-2/ IFGO0006
C      INTEGER*2  C,COMPAR,NSTTOK,TFLG,TOKEN(80)                IFGO0007
C
C      COMMON BLOCKS                                           IFGO0008
C      COMMON /INPTR /LASTC                                     IFGO0009
C      COMMON /TKNFLG/TFLG                                      IFGO0010
C      COMMON /TKN    /TOKEN                                    IFGO0011
C      COMMON /NSTFLG/NSTTOK                                    IFGO0012
C
C      TFLG=0                                                  IFGO0013
C      CHECK FOR THE CODITION PART                             IFGO0014
C      CALL GETTOK                                             IFGO0015
C      IF (NSTTOK.EQ.1) GO TO 300
C      IF (TOKEN(1).EQ.LBRK) GO TO 20
C      ERROR--'(' MISSING
C      CALL ERROR (2)
C      THE CODITION PART IS STORED IN CODITION BUFFER
20  DO 30 I=1,80
C=C=TKEN(I)
IF (C.EQ.EOS) GO TO 35

```

*** IFGO ***

	CALL PUTCND (C)	IFG00016
30	CONTINUE	IFG00017
35	CONTINUE	IFG00018
C	BALPAR TO BALANCE THE PARENTHESIS	*****
	CALL BALPAR	IFG00019
C	WHEN CNTRL IS 1 STATEMENT IS IF-STATEMENT	*****
	IF (CNTRL.EQ.0) GO TO 100	IFG00020
	IF (LASTC.EQ.6) GO TO 50	IFG00021
	CALL GETTOK	IFG00022
	IF (NSTTOK.NE.1) GO TO 40	IFG00023
	LASTC=6	IFG00024
	GO TO 50	IFG00025
40	IF (COMPAR(TOKEN,THEN).EQ.1) GO TO 100	IFG00026
	TFLG=1	IFG00027
50	CALL PUTIF(0,LAB)	IFG00028
C	THE FORTRAN IF STATEMENTS ARE PRINTED AS SUCH AND	*****
C	TYPE IS SET TO 1	*****
	TYPE=1	IFG00029
	RETURN	IFG00030
100	CALL PUTIF(1,LAB)	IFG00031
C	THE CODE IF(.NOT.(CONDITTON))----- IS GENERATED	*****
	TYPE=0	IFG00032
	RETURN	IFG00033
300	LASTC=6	IFG00034
	TYPE=0	IFG00035
	RETURN	IFG00036
	END	IFG00037

*** PUTIF ***

```

SUBROUTINE PUTIF (TYPE,LAB)                                PUTI0000
C*****                                                    *****
C
C   THIS ROUTINE GENERATES CODE 'IF(.NOT.(CONDITION)) ...'  *****
C   OR 'IF (CONDITION) ....' WHEN TYPE=1 OR 0; THE CONDITION *****
C   IS TAKEN FROM THE CONDITION BUFFER.                    *****
C*****                                                    *****
C
C   DIMENSION IFNOT(8)                                     PUTI0001
C   INTEGER*2 CONBUF(1254),TOKEN(80),INBUF(80),INLAB(5),C,CNT,TFLG PUTI0002
C   INTEGER*2 IF(2)/'I','F',IFNOT/'I','F',/'(','.',',','N','O','T','.',',',PUTI0003
C   *          BLANK/' ',ZERO/'0',EOF/'-1',EOS/'-2',
C   *          GOTO(8)/' ','G','O',' ','T','O',' ','
C   INTEGER I,I1,I2,J,LASTC,PTR,TYPE                       PUTI0006
C   COMMON BLOCKS                                         *****
C   COMMON /CONBLK/ CONBUF                                PUTI0007
C   COMMON /CONPTR/ PTR                                  PUTI0008
C   COMMON /INRLK / INBUF                                PUTI0009
C   COMMON /INPTR / LASTC                                PUTI0010
C   COMMON /INLBL / INLAB                                PUTI0011
C   COMMON /SYSOUT/ ISYOUT                                PUTI0012
C   COMMON /TKNFLG/ TFLG                                  PUTI0013
C   COMMON /TKN / TOKEN                                  PUTI0014
C*****                                                    *****
C   DO 10 I=6,14                                         PUTI0015
C   CONBUF(I)=BLANK                                       PUTI0016
C   CONTINUE                                             PUTI0017
C   MOVE STATEMENT LABEL TO CONBUF                       *****
C   DO 20 I=1,5                                           PUTI0018

```

*** PUTIF ***

	CONBUF(I)=INLAB(I)	PUTI0019
20	CONTINUE	PUTI0020
C	CLEAR INLAB	*****
	CALL CLRFLD (INLAB,1,5)	PUTI0021
25	CONTINUE	PUTI0022
30	IF(TYPE.EQ.1) GO TO 1000	PUTI0023
C	WHEN TYPE=0 'IF' IS MOVED TO 14TH COLUMN OF CONBUF	*****
	CONBUF(13)=IF(1)	PUTI0024
	CONBUF(14)=IF(2)	PUTI0025
	I1=13	PUTI0026
	IF(TFLG.EQ.0) GO TO 100	PUTI0027
C	UNPROCESSED TOKEN IS PUT IN CONBUF	*****
	DO 50 I=1,80	PUTI0028
	C=TOKEN(I)	PUTI0029
	IF (C.EQ.EOS) GO TO 60	PUTI0030
	CALL PUTCND(C)	PUTI0031
50	CONTINUE	PUTI0032
60	TFLG=0	PUTI0033
C	THE UNCOVERED PART OF THE STATEMENT IS PUT IN CONBUF	*****
100	CALL GETC(C)	PUTI0034
	IF(LASTC.EQ.7) GO TO 200	PUTI0035
150	CALL PUTCND(C)	PUTI0036
	GO TO 100	PUTI0037
200	CNT=INBUF(6)	PUTI0038
	IF (.NOT. ((CNT.EQ.BLANK).OR.(CNT.EQ.ZERO))) GO TO 150	PUTI0039
	LASTC=6	PUTI0040
	I2=78	PUTI0041
	GO TO 2000	PUTI0042
C	THE FOLLOWING PART GENERATES THE CODE	*****
C	'IF (.NOT.(CONDITION)) GO TO LAB'	*****

*** PUTIF ***

1000	DO 1100 I=1,8	PUTI0043
	J=I+6	PUTI0044
C	MOVE 'IF (.NOT.' IN CONBUF LEFT TO CONDITION	*****
	CONBUF(J)=IFNOT(I)	PUTI0045
1100	CONTINUE	PUTI0046
C	MOVE 'GO TO' IN CONBUF RIGHT TO THE CONDITION	*****
	DO 1200 I=1,8	PUTI0047
	C=GOTO(I)	PUTI0048
	CALL PUTCND(C)	PUTI0049
1200	CONTINUE	PUTI0050
	I=LAB	PUTI0051
C	CONVERT LAB INTO CHARACTER FORM AND MOVE TO CONBUF	*****
	CALL DGTCHR(I,INLAB)	PUTI0052
	DO 1300 I=1,5	PUTI0053
	C=INLAB(I)	PUTI0054
	CALL PUTCND(C)	PUTI0055
1300	CONTINUE	PUTI0056
	CALL CLRFLD (INLAB,1,5)	PUTI0057
	T1=7	PUTI0058
	T2=72	PUTI0059
C	FOLLOWING PART PRINTS THE CONTENTS OF CONBUF	*****
2000	IF (PTR.EQ.I2) GO TO 2400	PUTI0060
	J=PTR+1	PUTI0061
	IF (PTR.GT.I2) GO TO 2100	PUTI0062
	K=I2	PUTI0063
	GO TO 2200	PUTI0064
2100	IQ=(PTR-I2)/66	PUTI0065
	T=(IQ*66)+I2	PUTI0066
	IF(I.NE.PTR) GO TO 2150	PUTI0067
	K=PTR	PUTI0068

*** PUTTF ***

	GO TO 2400	PUTI0069
2150	K=((IQ+1)*66)+I2	PUTI0070
2200	DO 2300 I=J,K	PUTI0071
	CONBUF(I)=BLANK	PUTI0072
2300	CONTINUE	PUTI0073
2400	WRITE (ISYOUT,2500) (CONBUF(I),I=1,5),(CONBUF(I),I=I1,I2)	PUTI0074
2500	FORMAT (5A1,' ',66A1,'+',6X)	PUTI0075
	IF(PTR.LE.I2) GO TO 3000	PUTI0076
	I2=I2+1	PUTI0077
	WRITE (ISYOUT,2600) (CONBUF(I),I=I2,K)	PUTI0078
2600	FORMAT (' *',66A1,'+',6X)	PUTI0079
3000	PTR=14	PUTI0080
	RETURN	PUTI0081
	END	PUTI0082

*** DDCODE ***

```

SUBROUTINE DDCODE (ID,NLASTC)                                DDC00000
C*****                                                    *****
C                                                                 *****
C   THIS ROUTINE PROCESS 'DO' STATEMENT, IN CASE THE STATEMENT *****
C   IS FORTRAN 'DO' VALUE OF (ID) IS RETURNED AS 1, OTHERWISE 0 *****
C   A GENERATED LABEL IS INSERTED BETWEEN 'DO' AND INDEX, WITH *****
C   AN ENTRY ON CONTEXT STACK *****
C                                                                 *****
C*****                                                    *****
C   INTEGER LASTC,NLASTC,LAB,LARGEN                          DDC00001
C   INTEGER*2 DIGITS(10)/'1','2','3','4','5','6','7','8','9','0'/, DDC00002
C   * DO(2)/'0','0'/,BUF(80),TOKEN(80),INLAB(5),OBUF(80)/80*' '/,ID DDC00003
C   INTEGER*2 BLANK/' '/,LEXDO/-31/,FOS/-2/,STAR/'*'/        DDC00004
C   COMMON BLOCKS                                           *****
C   COMMON/TKN/TOKEN                                         DDC00005
C   COMMON /INBLK / BUF                                       DDC00006
C   COMMON /INPTR / LASTC                                     DDC00007
C   COMMON /INLBL / INLAB                                     DDC00008
C   COMMON /SYSOUT/ ISYOUT                                    DDC00009
C                                                                 *****
C   IF(NLASTC.EQ.1) NLASTC=7                                DDC00010
C   N=LASTC                                                  DDC00011
C   CALL GETTOK                                              DDC00012
C   WHEN THE TOKEN AFTER 'DO' IS AN INTEGER (DIGITS), THE *****
C   STATEMENT IS RECOGNIZED AS FORTRAN 'DO' *****
C   DO 10 I=1,10                                            DDC00013
C   IF (TOKEN(1).EQ.DIGITS(I)) GO TO 500                    DDC00014
10  CONTINUE                                                DDC00015
   LAB=LARGEN(2)                                           DDC00016

```

*** DDCODE ***

	CALL STACK (LEXDD,LAB-1)	DDC00017
	DO 30 I=1,5	DDC00018
	OBUF(I)=INLAB(I)	DDC00019
30	CONTINUE	DDC00020
	CALL CLRFLD (OBUF,6,80)	DDC00021
	OBUF(7)=DO(1)	DDC00022
	OBUF(8)=DO(2)	DDC00023
C	CONVERT LAB INTO CHARACTER FORM	*****
	CALL DGTCHR (LAB,INLAB)	DDC00024
	J=0	DDC00025
	DO 40 I=10,14	DDC00026
	J=J+1	DDC00027
	OBUF(I)=INLAB(J)	DDC00028
40	CONTINUE	DDC00029
	I=15	DDC00030
	DO 50 J=1,80	DDC00031
	I=I+1	DDC00032
	IF (TOKEN(J).EQ.FOS) GO TO 60	DDC00033
	OBUF(I)=TOKEN(J)	DDC00034
50	CONTINUE	DDC00035
60	I=I-1	DDC00036
65	IF (LASTC.EQ.72) GO TO 200	DDC00037
	LASTC=LASTC+1	DDC00038
	IF(I.LT.72) GO TO 80	DDC00039
	WRITE (ISYOUT,70) (OBUF(IJ),IJ=1,72)	DDC00040
70	FORMAT (72A1,'+',6X)	DDC00041
	OBUF(6)=STAR	DDC00042
	I=6	DDC00043
80	IF (OBUF(LASTC).EQ.BLANK) GO TO 100	DDC00044
	I=I+1	DDC00045

*** DOCODE ***

	OBUF(I)=BUF(LASTC)	DOC00046
100	GO TO 65	DOC00047
200	WRITE (ISYOUT,70) (OBUF(IJ),IJ=1,72)	DOC00048
	TD=0	DOC00049
	RETURN	DOC00050
500	IF(LASTC.LT.N) LASTC=6	DOC00051
	WRITE (ISYOUT,600) INLAB,(BUF(I),I=NLASTC,72)	DOC00052
600	FORMAT (5A1,' ',66A1,8X)	DOC00053
	TD=1	DOC00054
	RETURN	DOC00055
	END	DOC00056

*** WHILE ***

```

SUBROUTINE WHILE
C*****
C
C ROUTINE TO PROCESS 'WHILE' STATEMENT
C*****
C*****
C
C INTEGER LAB,LARGEN,I
C INTEGER*2 LEXWHL/-32/,BLANK/' '/,INLAB(5)
C COMMON/INLRL/INLAB
C
C 3 GENERATED LABEL RESERVED
C LAB=LARGEN(3)
C DO 10 I=1,5
C IF(INLAB(I).NE.BLANK) GO TO 100
10 CONTINUE
C GO TO 200
C FOR LABELED STATEMENT THE CODE 'LABEL CONTINUE' IS GENERATED
100 CALL LABCON(INLAB)
200 I=LAB+1
C CALL DGTCHR(J,INLAB)
C CODE 'IF (.NOT.(CONDITION)) GO TO LAB+2' IS GENERATED AND
C ENTRY IS MADE ON CONTEXT STACK.
C CALL IFGO(0,LAB+2,T)
C CALL STACK (LEXWHL,LAB)
C RETURN
C END
WHILO000
*****
*****
*****
*****
WHILO001
WHILO002
WHILO003
*****
*****
WHILO004
WHILO005
WHILO006
WHILO007
WHILO008
*****
WHILO009
WHILO010
WHILO011
*****
WHILO012
WHILO013
WHILO014
WHILO015

```

*** REPGEN ***

```

      SURROUTINE REPGEN
C*****
C
C      PROCESSES THE 'REPEAT' STATEMENT AND MAKES AN ENTRY IN
C      CONTEXT STACK.
C*****
C*****
C      INTEGER*2 INLAB(5),BLANK/' '/,LEXREP/-35/
C      INTEGER LAB,LARGEN
C      COMMON/INLRL/INLAB
C
C      LAB=LARGEN(3)
C      3 GENERATED LABELS ARE RESERVED
C      STACK 'REPEAT'
C      CALL STACK (LEXREP,LAB)
C      DO 10 J=1,5
C      IF(INLAB(J).NE.BLANK) GO TO 50
10    CONTINUE
C      GO TO 100
C      CODES GENERATED ' INLAB  CONTINUE' FOR LABELED STATEMENT
C      ' LAB    CONTINUE'
50    CALL LABCON(INLAB)
100   CALL OUTCON(LAB)
      RETURN
      END
      REPG0000
      *****
      *****
      *****
      *****
      REPG0001
      REPG0002
      REPG0003
      *****
      REPG0004
      *****
      REPG0005
      REPG0006
      REPG0007
      REPG0008
      REPG0009
      *****
      *****
      REPG0010
      REPG0011
      REPG0012
      REPG0013

```

*** UNGEN ***

```

SUBROUTINE UNGEN(LRL)
C*****
C
C      GENERATE CODE FOR 'UNTIL' UNIT
C*****
C*****
C      INTEGER LBL
C      INTEGER*2 TOKEN(80),INLAB(5),BLANK/' '/,TFLG,EOS/-2/
C      COMMON BLOCKS
C      COMMON /INLBL / INLAB
C      COMMON /TKN / TOKEN
C      COMMON /TKNFLG/ TFLG
C      COMMON /SYSOUT/ ISYOUT
C
C      CALL GETTOK
C      TFLG=0
C      LABEL IS NOT ALLOWED WITH UNTIL
C      DO 10 I=1,5
C      IF (INLAB(I).NE.BLANK) GO TO 50
10    CONTINUE
C      GO TO 60
C      ERROR--LABEL ILLFGAL HERE
50    CALL ERROR (1)
60    CONTINUE
C      CHECK FOR INFINITE REPEAT LOOP (UNCONDITIONAL REPEAT)
C      IF (TOKEN(1).EQ.EOS) GO TO 200
C      FOLLOWING PART GENERATES CODE FOR CONDITIONAL REPEAT
C      I=LBL+1
C      CALL DGTCHR(I,INLAB)

```

UNGE0000

 UNGE0001
 UNGE0002

 UNGE0003
 UNGE0004
 UNGE0005
 UNGE0006

 UNGE0007
 UNGE0008

 UNGE0009
 UNGE0010
 UNGE0011
 UNGE0012

 UNGE0013
 UNGE0014

 UNGE0015

 UNGE0016
 UNGE0017

*** UNGEN ***

	DO 100 I=1,80	UNGE0018
	IF (TOKEN(I).EQ.FDS) GO TO 150	UNGE0019
	CALL PUTCND (TOKEN(I))	UNGE0020
C	CONDITION IS BEING PUT IN CONBUF	*****
100	CONTINUE	UNGE0021
150	CALL BALPAR	UNGE0022
C	CODE 'IF (CONDITION) GO TO LBL' IS PRINTED	*****
	CALL PUTIF(0,LBL)	UNGE0023
	WRITE (ISYOUT,160) LBL	UNGE0024
160	FORMAT (' *GO TO ',I5,55X,'++',6X)	UNGE0025
	GO TO 300	UNGE0026
C	CODE ' LBL+1 GO TO LBL '	*****
200	CALL OUTGON(LBL+1,LBL)	UNGE0027
C	CODE ' LBL+2 CONTINUE '	*****
300	CALL OUTCON(LBL+2)	UNGE0028
	RETURN.	UNGE0029
	END	UNGE0030

*** FORGEN ***

```

SUBROUTINE FORGEN
C*****
C
C      PROCESS THE 'FOR' STATEMENT.
C      INITIALIZATION, CONDITION AND REINITIALIZATION PARTS ARE
C      SEPARATED BY ';', RE-INITIALIZATION PART IS PUT IN AREA
C      RETEXP.
C*****
C*****
C      INTEGER    T,J,K,LAB,LABGEN,NFOR,PTR,PTR1,PTR2
C      INTEGER*2  CONBUF(1254),TOKEN(80),INLAB(5),RETEXP(4,66)
C      INTEGER*2  LEXFOR/-34/,LBRK/'('/,RBRK/')'/,SCOLN/';'/,EOS/-2/,
C      * BLANK/' '/,NSMCLN,NSTTOK,RRNFLG
C      COMMON BLOCKS
C      COMMON /CONBLK/ CONBUF
C      COMMON /CONPTR/ PTR
C      COMMON /INLBL / INLAB
C      COMMON /NMBFOR/ NFOR
C      COMMON /NSTFLG/ NSTTOK
C      COMMON /RINBLK/ RETEXP
C      COMMON /RNFLG / RRNFLG
C      COMMON /SYSOUT/ ISYOUT
C      COMMON /TKN   / TOKEN
C
C      RRNFLG=0
C      CALL GETTOK
C      IF(NSTTOK.EQ.1) GO TO 1100
C      IF(TOKEN(1).EQ.LBRK) GO TO 100
C      '(' MISSING

```

```

FORG0000
*****
*****
*****
*****
*****
*****
*****
*****
FORG0001
FORG0002
FORG0003
FORG0004
*****
FORG0005
FORG0006
FORG0007
FORG0008
FORG0009
FORG0010
FORG0011
FORG0012
FORG0013
*****
FORG0014
FORG0015
FORG0016
FORG0017
*****

```

*** FORGEN ***

	CALL ERROR (2)	FORG0018
	PTR=14	FORG0019
C	LBRK MISSING	*****
	DO 10 I=1,80	FORG0020
	IF(TOKEN(I).EQ.EOS) GO TO 100	FORG0021
	CALL PUTCND (TOKEN(I))	FORG0022
10	CONTINUE	FORG0023
100	CALL BALPAR	FORG0024
	NSMCLN=0	FORG0025
	DO 200 I=15,PTR	FORG0026
	IF(CONRUF(I).NE.SCOLN) GO TO 200	FORG0027
	NSMCLN=NSMCLN+1	FORG0028
	IF(NSMCLN.EQ.1) PTR1=I	FORG0029
	IF(NSMCLN.EQ.2) PTR2=I	FORG0030
200	CONTINUE	FORG0031
	IF (NSMCLN.EQ.2) GO TO 300	FORG0032
C	ERROR--'SYNTAX ERROR IN 'FOR' STATEMENT'	*****
	CALL ERROR (6)	FORG0033
	PTR=14	FORG0034
	RETURN	FORG0035
C	3 GENERATED LABELS	*****
300	LAB=LAPGEN(3)	FORG0036
	IF(CONRUF(PTR).EQ.PBRK) PTR=PTR-1	FORG0037
C	FOLLOWING PART PUTS THE RE-INITIALISTING PART IN AREA REIEXP	*****
	NFOR=NFOR+1	FORG0038
	DO 350 I=1,66	FORG0039
	REIEXP(NFOR,I)=@BLANK	FORG0040
350	CONTINUE	FORG0041
	J=PTR2+1	FORG0042
	K=1	FORG0043

*** FORGEN ***

	DO 400 I=J,PTR	FORG0044
	RETEXP(NFOR,K)=CONBUF(I)	FORG0045
	IF(K.GF.66) GO TO 500	FORG0046
	K=K+1	FORG0047
400	CONTINUE	FORG0048
500	I=PTR1-14	FORG0049
	K=15	FORG0050
	DO 600 J=1,I	FORG0051
	TOKEN(J)=CONBUF(K)	FORG0052
	IF(J.EQ.66) GO TO 700	FORG0053
	K=K+1	FORG0054
600	CONTINUE	FORG0055
	CALL CLRFLD(TOKEN,J,66)	FORG0056
700	WRITE (ISYOUT,750) INLAR,(TOKEN(I),I=1,66)	FORG0057
750	FORMAT (SA1,' ',66A1,'+',6X)	FORG0058
C	ABOVE IS TO PRINT INITIALISING STATEMENT	*****
	I=LAB	FORG0059
	CALL DGTCHR(I,INLAR)	FORG0060
	J=PTR2-PTR1-1	FORG0061
	K=15	FORG0062
	PTR1=PTR1+1	FORG0063
	DO 1000 I=1,J	FORG0064
	CONBUF(K)=CONBUF(PTR1)	FORG0065
	K=K+1	FORG0066
	PTR1=PTR1+1	FORG0067
1000	CONTINUE	FORG0068
	PTR=J+14	FORG0069
C	CODE 'IF (NOT.(CONDITION)) GO TO LAB+2 '	*****
	CALL PUTIF(1,LAB+2)	FORG0070
C	SET RE-RUN FLAG ON AND PUT THE 'FOR' ON STACK	*****

*** FORGEN ***

	RRNFLG=1	FORG0071
	CALL STACK (LEXFOR,LAB)	FORG0072
	RETURN	FORG0073
1100	CALL ERROR (6)	FORG0074
	RETURN	FORG0075
	END	FORG0076

*** OTHERS ***

```

SUBROUTINE OTHERS (NLASTC)
C*****
C
C      THIS ROUTINE OUTPUTS A STATEMENT WHICH IS NOT CLASSIFIED
C      AS ANY OF THE EXTENDED STATEMENTS, WITHOUT ANY CHANGE.
C*****
C*****
C      INTEGER  NLASTC,LASTC,LNTH,TKLNTH,ISYOUT
C      INTEGER*2 BUF(80),TOKEN(80),RRNFLG,C
C      INTEGER*2 BLANK/' '/,EOS/-2/,STAR/'*'/,ZERO/'0'/
C
C      COMMON BLOCKS
C      COMMON /INRLK / BUF
C      COMMON /INPTR / LASTC
C      COMMON /PLENTH/ LNTH
C      COMMON /RNFLG / RRNFLG
C      COMMON /SYSOUT/ ISYOUT
C      COMMON /TKN  / TOKEN
C      COMMON /TKSIZE/ TKLNTH
C
C      IF (TOKEN(1).EQ.EOS) RETURN
C      IF (RRNFLG.EQ.0) GO TO 8
C      IF (LNTH.EQ.0) GO TO 6
C      WRITE (ISYOUT,10) (TOKEN(I),I=1,LNTH)
C      BUF(6)=STAR
C      GO TO 25
6      IF (LASTC.NE.6) GO TO 7
C      WRITE (ISYOUT,10) (TOKEN(I),I=1,TKLNTH)
C      GO TO 22
7      J=LASTC+1

```

OTHE0000

 OTHE0001
 OTHE0002
 OTHE0003

 OTHE0004
 OTHE0005
 OTHE0006
 OTHE0007
 OTHE0008
 OTHE0009
 OTHE0010

 OTHE0011
 OTHE0012
 OTHE0013
 OTHE0014
 OTHE0015
 OTHE0016
 OTHE0017
 OTHE0018
 OTHE0019
 OTHE0020

*** OTHERS ***

	WRITE (ISYOUT,110) (TOKEN(I),I=1,TKLNTH),(BUF(I),I=J,72)	OTHE0021
	GO TO 20	OTHE0022
8	IF (NLASTC.EQ .1) GO TO 25	OTHE0023
	WRITE (ISYOUT,10) (BUF(I),I=NLASTC,72)	OTHE0024
10	FORMAT (6X,74A1)	OTHE0025
20	LASTC=72	OTHE0026
	CALL GFIC(C)	OTHE0027
22	C=BUF(6)	OTHE0028
	IF((C.EQ.BLANK).OR.(C.EQ.ZERO)) GO TO 100	OTHE0029
25	CALL PUTLIN(BUF)	OTHE0030
	GO TO 20	OTHE0031
100	LASTC=6	OTHE0032
	RETURN	OTHE0033
110	FORMAT (6X,66A1,8X/5X,'*',66A1,8X)	OTHE0034
	END	OTHE0035

*** PUTSMT ***

```

SUBROUTINE PUTSMT(N)
C*****
C
C   OUTPUTS THE PART OF IN-BUFFER RIGHT TO THE POINTER 'N'
C   AS A COMMENT.
C*****
C*****
C
C   INTEGER*2 BUF(80)
COMMON/INBLK/BUF
COMMON /SYSOUT/ ISYOUT
C
C   IF (N.LE.6) N=6
WRITE (ISYOUT,10) (BUF(I),I=N,80)
10  FORMAT ('C+ ',75A1)
RETURN
END
PUTS0000
*****
*****
*****
*****
*****
*****
PUTS0001
PUTS0002
PUTS0003
*****
PUTS0004
PUTS0005
PUTS0006
PUTS0007
PUTS0008

```

*** PUTLIN ***

	SUBROUTINE PUTLIN (STRING)	PUTL0000
C	*****	*****
C	THIS ROUTINE PRINTS STRING AS SUCH.	*****
C	*****	*****
C	*****	*****
C	INTEGER*2 BLANK /' '/, STRING(80)	PUTL0001
	COMMON /SYSOUT/ ISYOUT	PUTL0002
C	*****	*****
	WRITE (ISYOUT,10) STRING	PUTL0003
10	FORMAT (80A1)	PUTL0004
	DO 20 I=1,80	PUTL0005
	STRING(I)=BLANK	PUTL0006
20	CONTINUE	PUTL0007
	RETURN	PUTL0008
	END	PUTL0009

*** LABCON ***

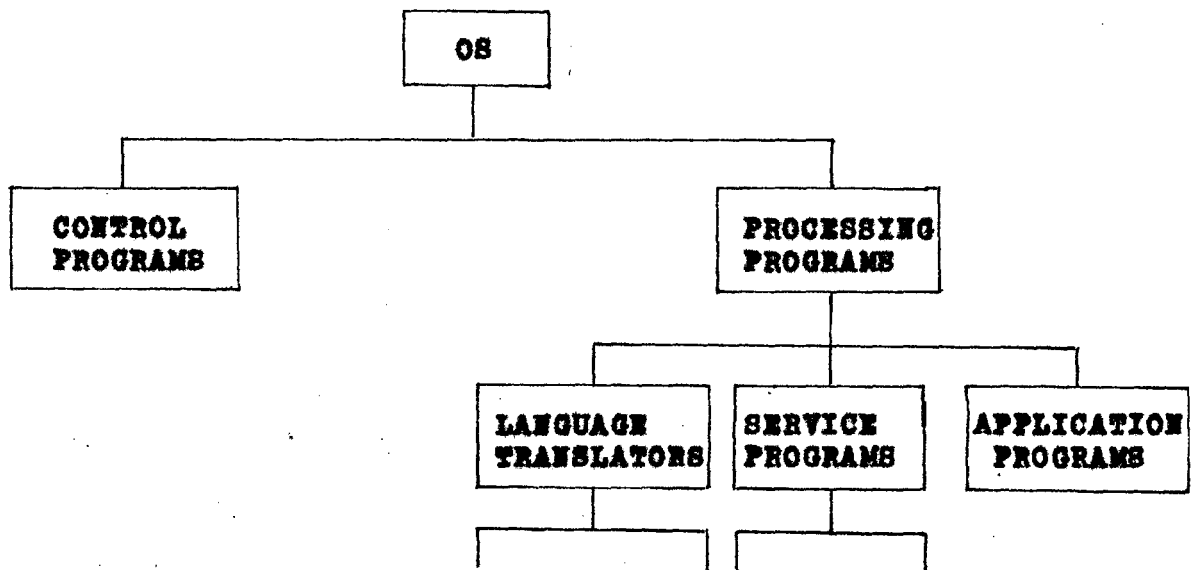
```

      SUBROUTINE LABCON (LABL)
C*****
C
C      GENERATE CODE 'LABL CONTINUE'
C*****
C*****
C      INTEGER*2 LABL(5)
      COMMON /SYSOUT/ ISYOUT
C
C      WRITE (ISYOUT,10) LABL
10    FORMAT (5A1,' CONTINUE',58X,'++',6X)
      RETURN
      END
LABC0000
*****
*****
*****
LABC0001
LABC0002
*****
LABC0003
LABC0004
LABC0005
LABC0006
```


OPERATING SYSTEM AND JOB CONTROL LANGUAGEOPERATING SYSTEM (OS)

The operating system (OS) introduces programs to the computing system, initiates their execution, and provides all the resources and services they require.

The programs and routines that compose the operating system are classified as control programs and processing programs as shown in the following figure.



FORTRAN PL/I COBOL LINKAGE
EDITOR
LOADER

The three main functions of the control programs are to accept and schedule jobs in a continuous flow (job management), to supervise each unit of work to be done (task management), and to simplify retrieval of all data, regardless of the way it is organized and stored (data management). The processing programs consist of language

translators (such as FORTRAN, PL/I, COBOL compilers), service programs (such as the linkage editor, loader), and application programs (such as user programs).

The operating system with which EXFOR has been interfaced is the OS-MFT (Multiprogramming with a Fixed number of Tasks) variation of OS. Several jobs can run concurrently under MFT.

OS-JCL (JOB CONTROL LANGUAGE)

JCL serves as a means of communication between the user's program and the system. JCL statements tell the system how to execute a job, request needed system facilities and describe required input/output(I/O) devices.

The normal JCL statements (or cards) are:

1. The JOB statement. This is the first control card; it marks the beginning of a job.
2. The EXEC (Execute) statement. This card follows the JOB card and names the program or procedure to execute.
3. The DD (Data Definition) statement. This card describes each data set (a file on tape or direct access device, or a deck of cards), and requests the allocation of I/O devices.
4. The Delimiter (/*) statement.

This is the "end-of-file" card for marking the end of a card deck.

Other JCL statements less often used are:

5. The Null (//) statement. The null may be used to mark the end of a job.

6. The comment (//*) statement. Comments may be coded in rest of the columns as an aid in documenting JCL.
7. The PROC statement. This card assigns default values for symbolic parameters in cataloged procedures.
8. The PEND statement. This card terminates the definition of a procedure.
9. The Command Statement. This card may be used by operators to enter operator commands from the input stream.

All the JCL statements (except the /* card) begin with a // in columns 1 and 2, followed by a name field, an operation field, and an operand field. The name field begins immediately after the second slash, while the name, operation, and operand fields are separated from each other by one or more blank spaces.

//name operation operand.

The name field identifies the control card so that other cards or system control blocks can refer to it. It can range from 1 to 8 characters in length and can contain any alphanumeric or national (@\$#) characters. However, the first character of the name must be alphabetic and in column 3.

The operation field specifies the type of control card: JOB, EXEC, DD, PROC, PEND, or an operator command.

The operand field contains parameters separated by commas.

The JOB, EXEC, and DD cards enable a complete job to be executed. As an example, suppose that some FORTRAN source language statements are to be compiled. Assume the FORTRAN compiler is a system program named IEYF0RT which requires three data sets: an input data set consisting of source language statements to compile, an output data set to print the compilation listing, and an output data set to contain the object module. Such a program might be executed by the following JCLs.

```

//TEST      JOB
//FORT      EXEC  PGM=IEYF0RT,PARM=(SOURCE,NOMAP)
//SYSPRINT  DD    SYSOUT=A
//SYSLIN    DD    DSNNAME=&LOADST,DISP=(NEW,PASS),
//           UNIT=2311,SPACE=(80,(200,100),RLSE),
//           DCB=BLKSIZE=80,VOL=SER=LIB400
//SYSIN     DD   *
              (source language statements to compile)
/*
//

```

The job is named TEST; the FORTRAN compiler named as IEYF0RT is executed in a job step named FORT. PARM specifies the parameters SOURCE and NOMAP to be passed to compiler.

```

//SYSPRINT  DD    SYSOUT=A

```

SYSOUT=A defines a print data set. The SYSOUT keyword instructs the system to queue the output on direct excess storage. A is the class of output devices defined by an installation, traditionally, SYSOUT=A is a printer.

The SYSLIN DD card defines a data set on direct access device 2311 volume name as LIB400, DSNNAME specifies the name of the data set &LOADST with disposition NEW and

PASS which says that after the job step FORT the data set &LOADST to be kept for the use of next job step (say linkage). The data sets with first character as "s" are temporary and deleted by the system as soon as the job is over. The SPACE parameter tells about the space to be allocated to the data set.

SPACE=(80,(200,100)RLSE) requests 200 blocks of 80 bytes each with an increment of 100 blocks for the secondary allocation in case the primary allocated space is full, maximum 15 secondary allocations can be given to a data set. RLSE releases all unused space when the data set is closed.

The DCB(Data Control Block) gives parameter for the data control block which is a table of data in core that describes each data set used by the program. The DCB parameters are LRECL (logical record length), BLKSIZE (blocksize), BUFNO (Number of buffers) and, RECFM (record format).

```
//SYSIN DD *
```

The asterisk (*) is a special code telling the system that data cards immediately follow the DD card. The end of the card deck is indicated by /* card.

CATALOGED PROCEDURE

JCL statements direct the operating system on the processing to be done on a job, and describe all I/O units required. Since these JCL statements are numerous and complex, the cards for frequently used procedures are kept on direct storage as a member of procedure library

named SYS1.PROCLIB. The user invokes these cataloged procedures by giving the system the name of the cataloged procedure rather than submitting the JCL cards. For example, a FORTRAN compilation linkage and execution job contains as many as 20 JCL cards, rather these statements are cataloged with a name FORTCLG. With use of cataloged procedure the JCL cards reduce to

```
//TEST JOB
//      EXEC  FORTCLG
      (the FORTRAN compile, link edit, go procedure
      is invoked)
//PORT.SYSIN DD *
[ source deck ]
/*
//GO.SYSIN   DD *
[ data cards ]
/*
//
```

When a procedure name is found in JCL statements the system searches the procedure library and the JCL statements are substituted in place of the procedure name, taking the parameters as specified with procedure name.

The procedure is defined by enclosing a set of JCLs between PROC and PEND statements. For example:

```
//RUN  PROC  NAME=LIB1
//STEP1 EXEC  PGM=ONE
//A    DD    DSN=&NAME, ---
//STEP2 EXEC  PGM=TWO
//B    DD    DSN=&NAME, ---
//      PEND
```

This procedure is cataloged in system procedure

library by adding it using the utility IEBUPDTE. Whenever the procedure is referred as

```
// EXEC RUN
```

in the JCLs, the procedure is substituted with each appearance of the symbolic parameter &NAME assigning a default value of LIB1. One might override this default for a run as

```
// EXEC RUN,PARM=USELIB
```

Each appearance of &NAME in the procedure is replaced by USELIB.

PROC FOR EXTFOR PREPROCESSOR

The complete process of translation from EXTFOR program to the execution may be divided into the following four job steps.

- Preprocessing
- Compilation
- Link Edit
- Execution

To perform these job steps, separate sets of JCL are defined. These job steps constitutes a procedure called FORTPCLG which is cataloged in procedure library SYS1.PROCLIB. This procedure can be used to perform all the 4 steps. In all 4 procedures are cataloged:

FORTP -For preprocessing an EXTFOR program.

FORTPC -For preprocessing and compilation of an EXTFOR program.

FORTPCL -For preprocessing, compilation and linkage edit of an EXTFOR program.

FORTPCLG-For preprocessing to execution of an EXTFOR program.

OS - EXTFOR INTERFACE AND OPERATING PROCEDURE

The FORTRAN source of EXTFOR preprocessor is compiled under the local OS environment. After linkage edit the load module which is now in executable form, is stored permanently as a member of processing program library of OS. The load module is preferably stored on the fastest device available. A temporary data set (file) is allocated each time the EXTFOR is loaded, to store the output of EXTFOR. The FORTRAN reads the input source (output of EXTFOR) from this temporary data set; rest of the procedure is same as followed for FORTRAN compilation.

The EXTFOR developed has been interfaced with OS-MFT on EC-1030 System at ORG Systems, Baroda. The EXTFOR load module (JFPREPRO) is stored as a member of library P.LOADLIB on direct access device 2311, volume serial no. LIB400 (system pack). The temporary data set &SRCSET on LIB400 is allocated to contain the output of EXTFOR. The FORTRAN compiler, which has IEYFORT as system name takes the further processing as desired by the user. The following procedure FORTPCLG is defined and cataloged to preprocess, compile, linkage edit, and finally execute the EXTFOR source program.

```

//FORTECLG PROC NAME='''NAME=MAIN''',SOURCE=SOURCE,LIST=HOLIST,      X
//          DECK=NODECK,MAP=MAP,LOAD=LOAD,LCNT=.,.,LINECNT=50;''',  X
//          LXREF=XREF,LLET=LET,LLIST=LIST
//**
//**          STEP TO PREPROCESS THE EXTFOR SOURCE PROGRAM          **
//PORTP     EXEC PGM=JYFPREPRO,COND=EVEN,TIME=2
//STEFLIB   DD   DSN=P.LOADLIB,UNIT=2311,DISP=SHR,VOL=SER=LIB400
//FT06FOO1  DD   SYSOUT=A
//SYSPRINT  DD   SYSOUT=A
//FT01FOO1  DD   DDNAME=SYSIN
//**          EXTFOR SOURCE ON CARDS          **
//FT03FOO1  DD   DSN=&SRCSET,DISP=(NEW,PASS),UNIT=2311,              X
//          SPACE=(CYL,(2,5,1),RLSE),                              X
//          DCB=(RECFM=FB,LRECL=80,BLKSIZE=1600),                  X
//          VOL=SER=LIB400
//**
//**          STEP TO COMPILE FORTRAN CODE          **
//FORT      EXEC PGM=IEYFORT,                                       X
//          PARM=(&NAME,&SOURCE,&LIST,&DECK,                          X
//          &MAP,&LOAD,&LCNT),COND=(4,LT,FORTP)
//SYSPRINT  DD   SYSOUT=A
//SYSLIN    DD   DSNNAME=&LOADSET,DISP=(NEW,PASS),UNIT=2311,        X
//          SPACE=(80,(200,100),RLSE),DCB=BLKSIZE=80,              X
//          VOL=SER=LIB400

```

```

//SYSIN DD DSNNAME=&SRCSRT,DISP=(OLD,DELETE)
// DD DNAME=SYSIN
//** SYSIN IS DATA SET NAME USER CAN CONCATENATE **
//** WITH PREPROCESSOR OUTPUT **
//**
//** LINKAGE EDIT STEP **
//LKED EXEC PGM=IEWL,PARM=(&LXREF,&LLET,&LLIST), X
// COND=((&L,LT,FORTP),(&L,LT,FORT))
//SYSPRINT DD SYSOUT=A
//SYSLIB DD DSNNAME=SYS1.FORTLIB,DISP=SHR
//SYSLMOD DD DSNNAME=&GOSRT(MAIN),DISP=(NEW,PASS),UNIT=2311, X
// SPACE=(1024,(20,10,1),RLSE),DCB=BLKSIZE=1024, X
// VOL=SER=LIB400
//SYSUT1 DD DSNNAME=&SYSUT1,UNIT=2311,SPACE=(1024,(20,10),RLSE), X
// DCB=BLKSIZE=1024,VOL=SER=LIB400
//SYSLIN DD DSNNAME=&LOADSET,DISP=(OLD,DELETE) X
// DD DNAME=SYSIN
//GO EXEC PGM=*.LKED.SYSLMOD, X
// COND=((&L,LT,FORTP),(&L,LT,FORT),(&L,LT,LKED))
//PT03FOO1 DD SYSOUT=A
//PT01FOO1 DD DNAME=SYSIN
//PT06FOO1 DD SYSOUT=A
// PENDING

```

To execute the EXTFOR program following set of JCLs is used:

```
//JOBNAM JOB
//          EXEC  FORTPCLG
//FORTP.SYSIN DD *
[EXTFOR source cards]
/*
//GO.SYSIN   DD *
[input card data]
/*
//
```

User can have EXTFOR source, the input and output data files (data sets) on any media. For source, the DD (Data Definition) card with name FORTP.SYSIN is replaced by DD of data set on which the source program exists. GO.SYSIN DD card defines the input file. The input/output files can be used in program with DD definition as:

```
//GO.FTnnFOO1 DD operands
```

The nn is the file reference number used in the program.

However, user can define his own set of procedures and data set definitions to fit his requirements and OS environment.

Following is the EXTFOR sample program:

IEF587 DIRECT SYSOUT=A.
//SAMPLE1 JOB
// EXEC FORTPCLG
//SYSIN DD *

** STATISTICS AT PREPROCESSOR LEVEL **

0 DIAGNOSTICS GENERATED

IEF373I STEP /FORTR / START 80363.1408
IEF374I STEP /FORTR / STOP 80363.1409 CPU 0MIN 13.58SEC MAIN 54K

```
C*****  
C   THIS PROGRAM SORTS A SEQUENCE OF N NUMBERS USING SHELL SORT METHOD *  
C   ILLUSTRATES USE OF FOR,WHILE,DO,REPEAT AND BREAK STATEMENT *  
C*****  
C
```

```
0001   INTEGER A(20)  
0002   READ (1,10) N,A  
0003   CALL SHLSRT (A,N)  
0004   WRITE (3,20) A  
0005   10  FORMAT (2I13)  
0006   20  FORMAT (1H ,20X,'THE SORTED SEQUENCE IS'/  
*       1H ,20X,10I5/1H ,20X,10I5)  
0007   STOP  
0008   END
```

SUBPROGRAMS CALLED

YMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOC
BCOM#	90	SHLSRT	94						

SCALAR MAP

YMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOC
	A0								

ARRAY MAP

YMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOC
	A4								

FORMAT STATEMENT MAP

YMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOC
10	F4	20	FA						

OPTIONS IN EFFECT IO,DKOI,SOURCE,NOLIST,NODECK,LOAD,MAP
 OPTIONS IN EFFECT NAME = MAIN , LINECNT = 50
 STATISTICS SOURCE STATEMENTS = 8,PROGRAM SIZE = 494
 STATISTICS NO DIAGNOSTICS GENERATED

```

0001      SUBROUTINE SHLSRT (SEQ,NUM)
0002      INTEGER SEQ(20),NUM,DIST,TEMP
      C+   FOR (DIST=4;(DIST.LT.NUM);DIST=DIST+DIST)
0003      DIST=4      ++
0004      90000 IF(.NOT.(DIST.LT.NUM)) GO TO 90002      ++
0005      90001 DIST=DIST+DIST      ++
0006      GO TO 90000      ++
0007      90002 CONTINUE      ++
0008      DIST=DIST-1
      C+   WHILE (DIST.GT.1) DIST=DIST/2
0009      90004 IF(.NOT.(DIST.GT.1)) GO TO 90005      ++
0010      DIST=DIST/2
0011      ND=NUM-DIST
0012      DO 90006,K=1,ND      ++
0013      I=K
      C+   REPEAT
0014      90008 CONTINUE      ++
0015      J=I+DIST
      C+   IF (SEQ(I).LE.SEQ(J)) THEN BREAK
0016      IF(.NOT.(SEQ(I).LE.SEQ(J))) GO TO 90012      ++
      C+   BREAK
0017      GO TO 90010      ++
0018      90012 CONTINUE      ++
0019      TEMP=SEQ(I)
0020      SEQ(I)=SEQ(J)
0021      SEQ(J)=TEMP
0022      I=I-DIST
      C+   UNTIL (I.GE.1)
0023      90009 IF(I.GE.1)      ++
      *GO TO 90008      ++
0024      90010 CONTINUE      ++
      C+   END
0025      90006 CONTINUE      ++
0026      90007 CONTINUE      ++
      C+   END
0027      90003 GO TO 90004      ++
0028      90005 CONTINUE      ++
0029      RETURN
0030      END

```

SCALAR MAP

SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION
DIST	DC	AREA	E0	NO	E4	K	E8	I	
J	F0	TEMP	F4						

ARRAY MAP

SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION	SYMBOL	LOCATION
SEQ	F8								

OPTIONS IN EFFECT ID,DKOI,SOURCE,NOLIST,NODECK,LOAD,MAP

OPTIONS IN EFFECT NAME = SHLSRT , LINECNT = 50

STATISTICS SOURCE STATEMENTS = 30,PROGRAM SIZE = 724

STATISTICS NO DIAGNOSTICS GENERATED

STATISTICS NO DIAGNOSTICS THIS STEP

IEF373I STEP /FORT / START 80363.1409
IEF374I STEP /FORT / STOP 80363.1409 CPU 0MIN 12.84SEC MAIN 106K

IEF373T STEP /LKFD / START 90363.1409
IEF374I STEP /LKFD / STOP 90363.1410 CPU 0MIN 07.683EC MAIN 68K
//GO.SYSIN DD *

THE SORTED SEQUENCE IS

2	3	4	7	8	8	21	32	35	59
62	83	95	123	236	365	532	568	586	986

IEF373I STEP /GO / START 80363.1410
IEF374I STEP /GO / STOP 80363.1411 CPU 0MIN 03.76SEC MAIN 48K
IEF375I JOB /SAMPLE1 / START 80363.1408
IEF376I JOB /SAMPLE1 / STOP 80363.1411 CPU 0MIN 37.86SEC

ERROR MESSAGES

The EXTFOR preprocessor gives error messages in the following format:

```
C***nn) JFPmmI text of the message
nn - the serial number of error in the current state-
ment.
mmm - the error code number.
```

The text of the message may be any one of the following list.

Code Message Text

1 WARNING--LABEL ILLEGAL HERE

In some of the EXTFOR extended features the statement label is not allowed, for example labelled ELSE, UNTIL parts of control structures.

```
IF (AMONT.GT.1000) THEN BREAK
10 ELSE INFRST = (AMONT*RATE*YRS)/100
The statement number 10 with ELSE is not allowed.
```

2 '(' - MISSING

The condition part of IF, FOR, WHILE, UNTIL is not preceded by '('.

3 UNMATCHING '('S

The number of '(' and ')' are not same in condition part of the control structure. For example

```
IF ((LASTC.LT.72).AND.(C.NE.BLANK)$THEN
      TOKEN(I) = C
```

The \$ points where a ')' should have been to match the outermost '('.

4 UNRECOGNIZABLE STATEMENT

The current statement has the first symbol as a valid keyword but the rest of the statement results an unpredicable syntax.

5 TOO MANY NESTED 'FOR' STATEMENTS

Maximum of 4 nested 'FOR' are allowed.

6 SYNTAX ERROR IN 'FOR' STATEMENT

Self explanatory.

7 TOO LARGE A CONDITION/STATEMENT

The maximum length of a condition or a statement is 1254 characters.

8 CONDITIONAL 'END/ELSE/UNTIL'

These keywords should not be used in the conditional part of a control structure. For example

```
IF (X(I).LT.0) THEN END
                ELSE UNTIL
```

9 UNMATCHING ELSE/UNTIL

Self explanatory

Example:

```
IF (I.EQ.10) THEN BREAK
A = 2
DO I=1,5
B(I)=0
ELSE -- error
END
```

10 UNEXPECTED END OF SUBPROGRAM

The source text is over or a new subprogram has appeared in the program without the matching END for WHILE, DO or BEGIN structures in the previous subprogram.

12 ILLEGAL LEVEL SPECIFIED

The level in BREAK or NEXT statement is not positive integer. For example

```
BREAK - 5
NEXT A
```

13 FORTRAN 'DO' STATEMENT ILLEGAL HERE

The FORTRAN DO statement is found just after THEN or ELSE in the IF statement or in FOR statement which is not allowed in EXTFOR.

14 BREAK/NEXT ILLEGAL HERE

The statement BREAK or NEXT is encountered without an enclosing loop or block structure.

15 SYNTAX

The error message points to the symbol in the statement by a "\$" where the error is encountered.

For errors at FORTRAN compilation level, refer to the FORTRAN manual.

At the end the EXTFOR specifies the number of errors encountered during preprocessing the source program. The message is

```
** STATISTICS AT PREPROCESSOR LEVEL **
nnn DIAGNOSTICS GENERATED
```

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